

Pointing, asymptotics, and random generation in unlabelled classes

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Motivations

Automatic methods for

- Enumeration (exact/asymptotic)
- Random generation (cf [Flajolet,F,Pivoteau'07])

in the unlabelled setting.

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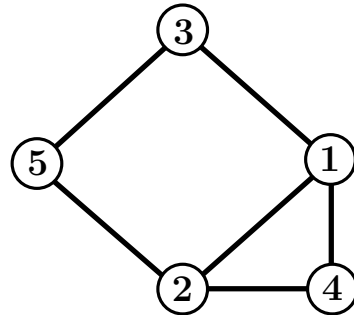
in the unlabelled setting.

References:

- Short version in SODA'07
- Long version almost finished, written in the framework of “Combinatorial species”, cf [Bergeron, Labelle, Leroux'98]

Labelled/Unlabelled structures

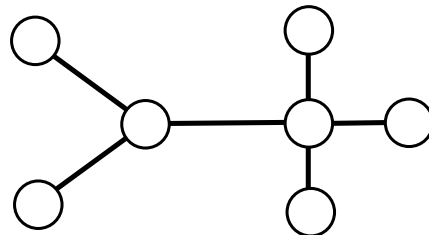
- **labelled** class $\mathcal{C} = \cup_n \mathcal{C}_n$



Labeled graph of size 5

$$\text{EGF : } C(x) = \sum_n \frac{1}{n!} c_n x^n, \text{ with } c_n = |\mathcal{C}_n|$$

- **Unlabelled** class $\tilde{\mathcal{C}} = \cup_n \tilde{\mathcal{C}}_n$



Unlabeled tree of size 7

$$\text{OGF : } C(x) = \sum_n \tilde{c}_n x^n, \text{ with } \tilde{c}_n = |\tilde{\mathcal{C}}_n|$$

Plan

- Decomposition strategy for labelled structures
 - Pointing + recursive decomp. + gen. functions
 - Examples: trees, planar graphs...

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 - Examples: trees, planar graphs...
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 - Difficulties due to **symmetries**
 - Solution: **unbiased pointing** + Pólya theory
- Application to **asymptotic enumeration**
- Application to **random generation**:
⇒ **Boltzmann samplers without rejection**

Decomposition strategy for labelled structures

Dictionary for EGF

- labelled class $\mathcal{C} = \cup_n \mathcal{C}_n$

$$\text{EGF : } C(x) = \sum_n \frac{1}{n!} c_n x^n, \text{ with } c_n = |\mathcal{C}_n|$$

- Simple **computation rule** for each **construction**:

Disjoint union	$\mathcal{C} = \mathcal{A} + \mathcal{B}$	$C(x) = A(x) + B(x)$
Cartesian product	$\mathcal{C} = \mathcal{A} \times \mathcal{B}$	$C(x) = A(x) \cdot B(x)$
Set	$\mathcal{C} = \text{Set}(\mathcal{A})$	$C(x) = \exp(A(x))$
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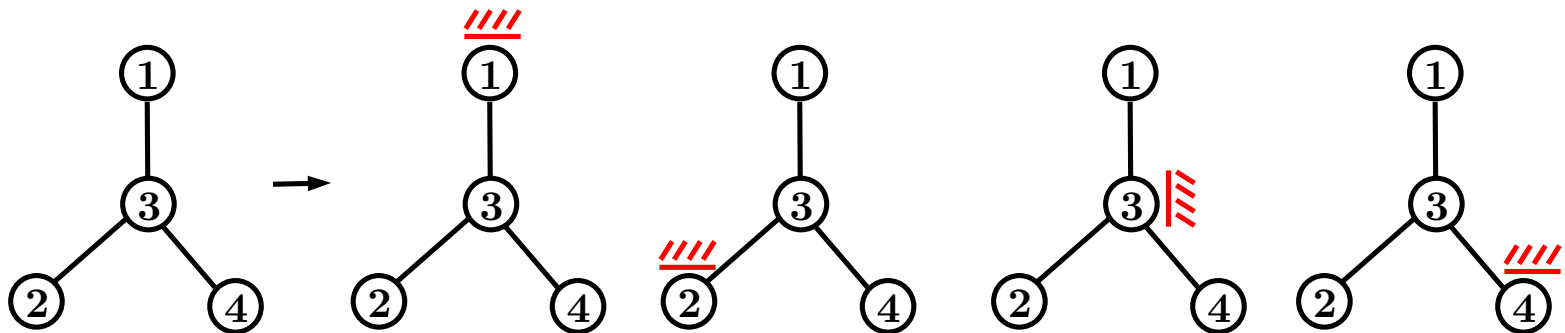
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- Remark.** **Substitution** rule implies **Set** rule since the EGF of the class Set is $\exp(z)$ (same for cycle, set, unoriented sequence, etc...)

Decomposition strategy for trees

- Goal: find t_n the number of (unrooted) trees of size n
- Important tool: pointing: $\mathcal{A} \mapsto \mathcal{A}^\bullet$
Let r_n be the number of rooted trees of size n

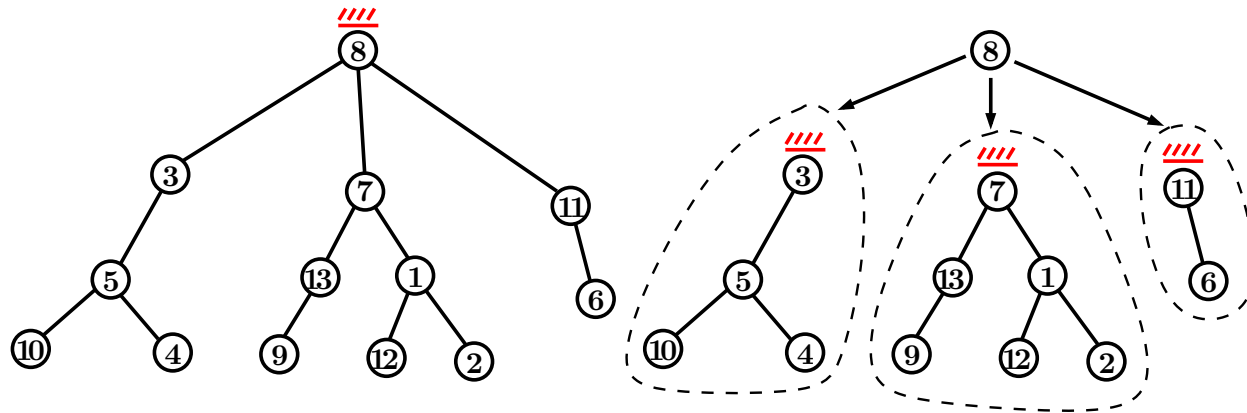


$$r_n = n \cdot t_n$$

\Rightarrow Counting trees reduces to counting rooted trees.

Rooted trees are decomposable

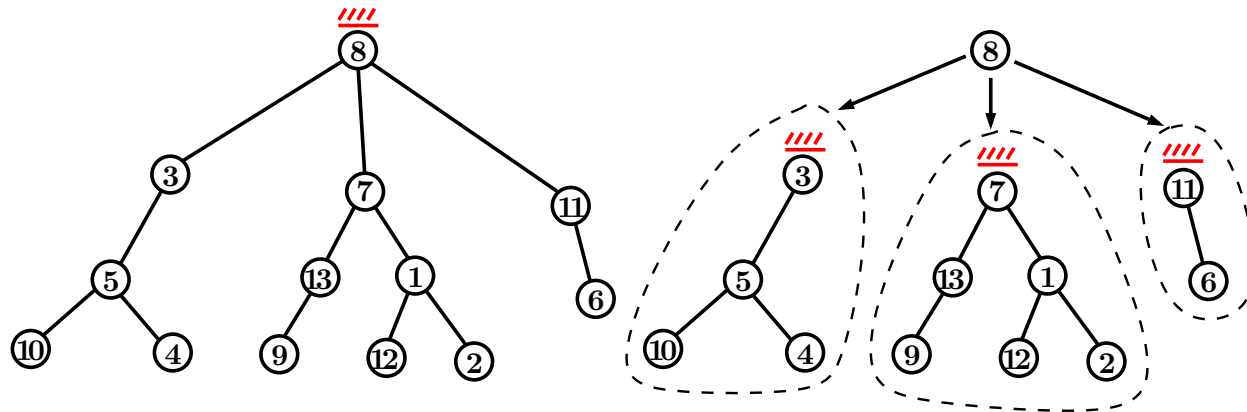
- The class \mathcal{R} of rooted trees satisfies the [decomposition](#)



$$\mathcal{R} = \mathcal{Z} \times \text{Set}(\mathcal{R}) \Rightarrow R(x) = x \exp(R(x))$$

Rooted trees are decomposable

- The class \mathcal{R} of rooted trees satisfies the **decomposition**



$$\mathcal{R} = \mathcal{Z} \times \text{Set}(\mathcal{R}) \Rightarrow R(x) = x \exp(R(x))$$

- Lagrange **inversion formula**

inverse of $R(x)$ is $R^{(-1)}(y) = y \exp(-y)$

\Rightarrow **Rooted trees:** $r_n = n^{n-1} \Rightarrow$ **Trees:** $c_n = n^{n-2}$

Counting labelled trees: summary

- **Decomposition** of rooted trees

$$R(x) = x \exp(R(x))$$

yields $r_n = n^{n-1}$ from Lagrange inversion formula

- **Pointing** relation: $t_n = r_n/n$:

$$R(x) = xT'(x)$$

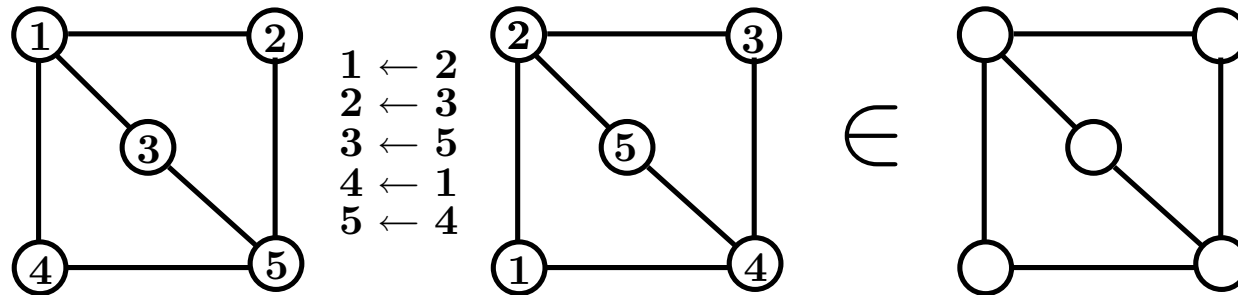
yields $t_n = n^{n-2}$

Same method applies for many classes (**planar graphs**)

Adaptation to the unlabelled setting

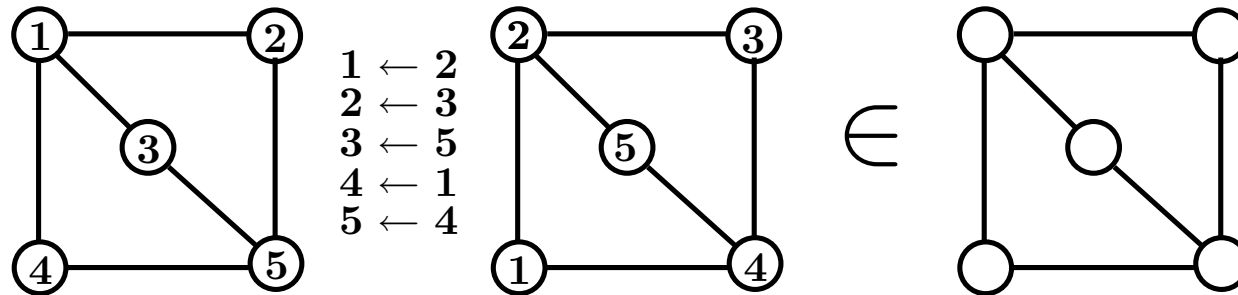
Unlabelled setting

- Unlabelled structures = labelled structures up to relabeling

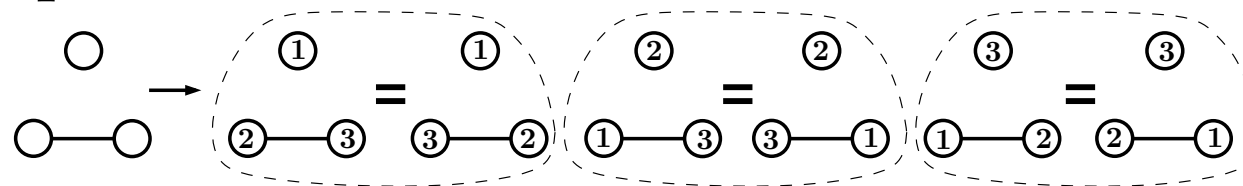


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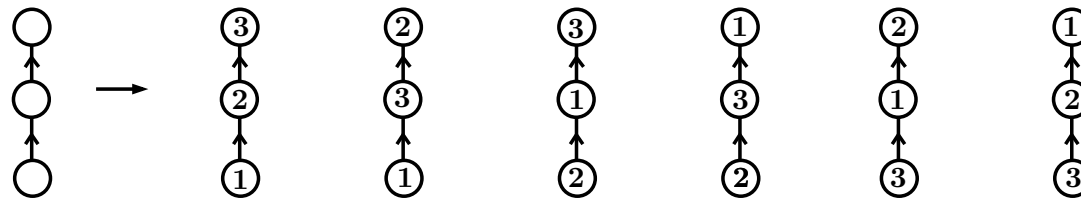
- **Unlabelled** structures = **labelled** structures **up to relabeling**



- **Examples:**



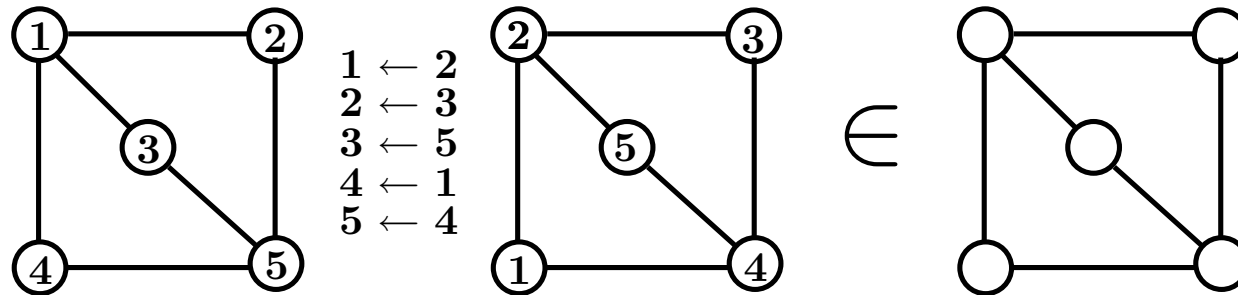
3 labeled objects (instead of $3! = 6$)



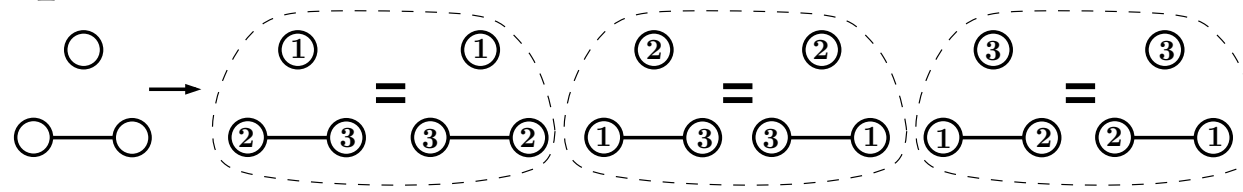
6 labeled objects (no symmetry)

Unlabelled setting

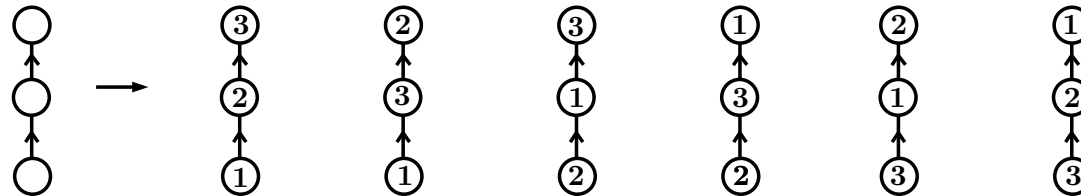
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- Examples:



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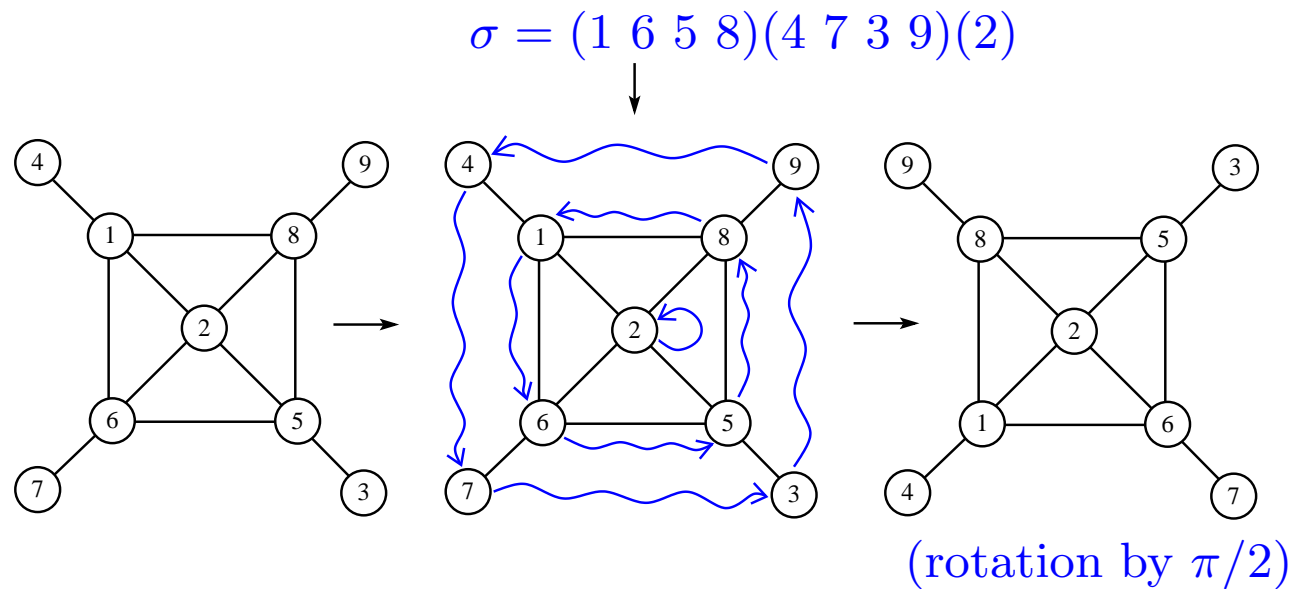
6 labeled objects (no symmetry)

- Unlabelled struct. size $n \rightarrow$ at most $n!$ labelled structures.
 $\Rightarrow \frac{1}{n!} a_n^{\text{label.}} \leq a_n^{\text{unlabel.}} \Rightarrow (\text{EGF}) A(x) \preceq \tilde{A}(x) (\text{OGF}).$

Symmetries

Let \mathcal{A} be a labelled class,

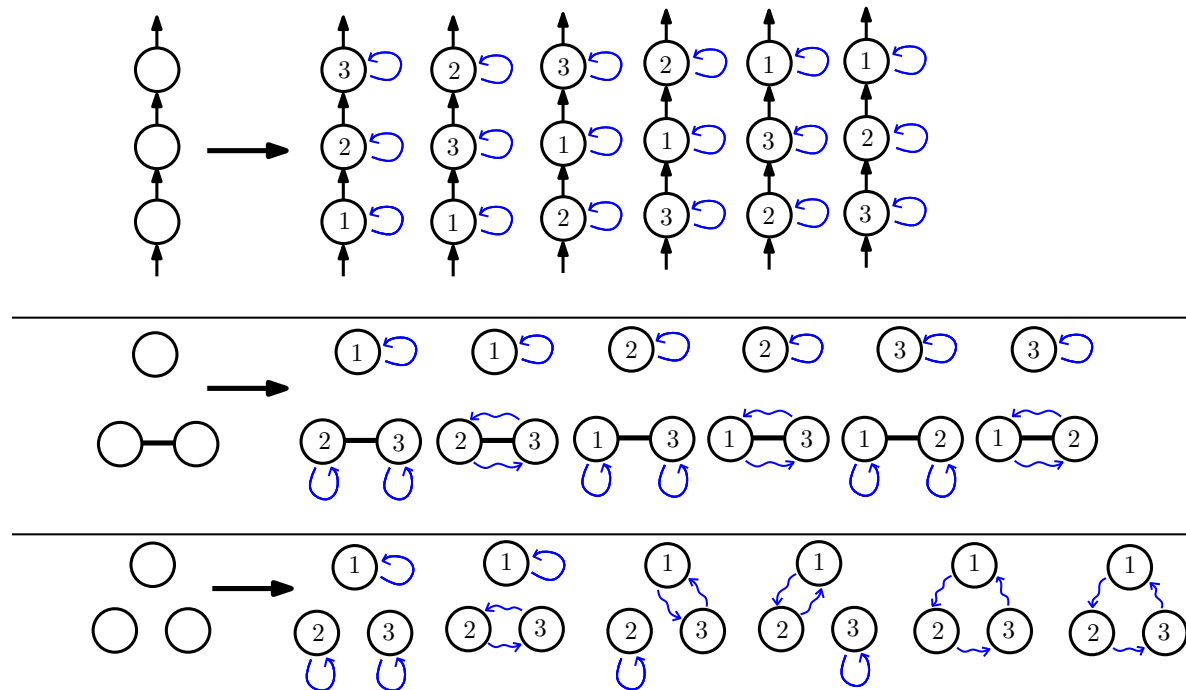
- A **symmetry** of size n on \mathcal{A} is a pair $(\sigma \in \mathfrak{S}_n, A \in \mathcal{A}_n)$ such that A is **fixed** by the action of σ .



Burnside's lemma

Given \mathcal{A} a labelled class (species of structures) let
 $\tilde{\mathcal{A}} = \mathcal{A}/\text{isomorphisms}$, $\text{Sym}(\mathcal{A}) = \{\text{Symmetries from } \mathcal{A}\}$

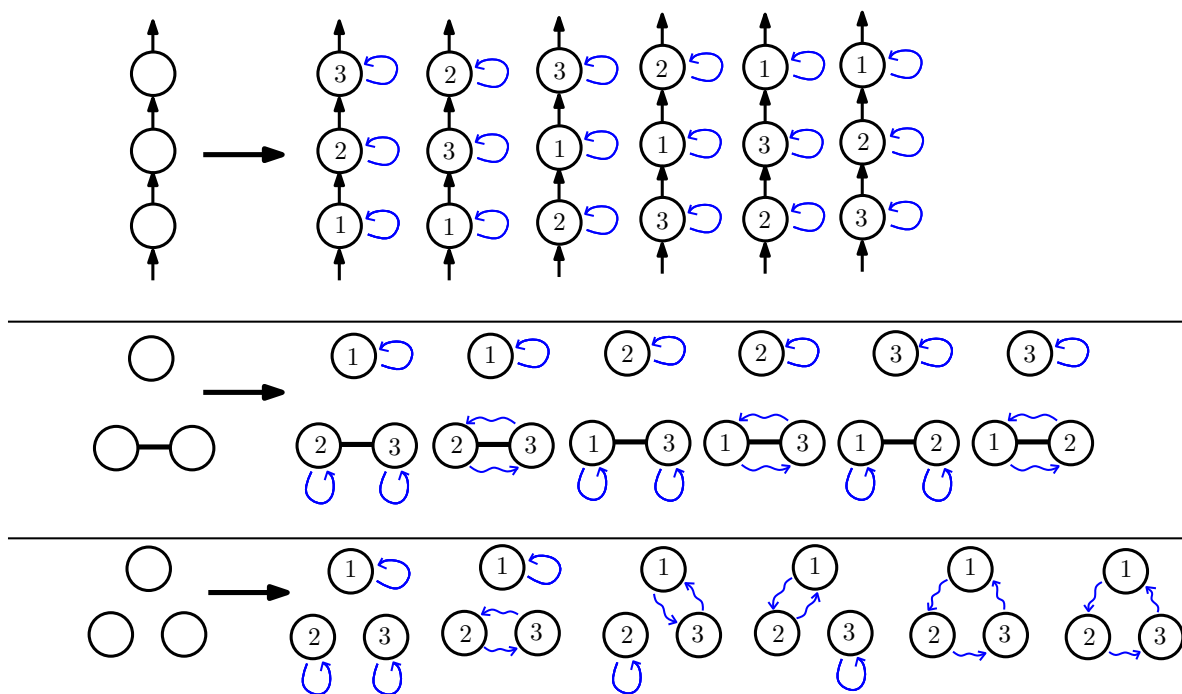
- **Burnside's lemma** $\Rightarrow \text{Sym}(\mathcal{A})_n \simeq n! \times \tilde{\mathcal{A}}_n$



Burnside's lemma

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- **Burnside's lemma** $\Rightarrow \text{Sym}(\mathcal{A})_n \simeq n! \times \tilde{\mathcal{A}}_n$



- Hence $\text{EGF of } \text{Sym}(\mathcal{A}) = \tilde{\mathcal{A}}(x) \text{ (OGF)}$

Cycle index sum

Let \mathcal{A} be a **labelled** class, $\text{Sym}(\mathcal{A})$ the **symmetry** class.

- **Refined weight** for $(\sigma, A) \in \text{Sym}(\mathcal{A})$

$$W(\sigma, A) := \frac{1}{n!} s_1^{c_1(\sigma)} s_2^{c_2(\sigma)} \dots s_n^{c_n(\sigma)}$$

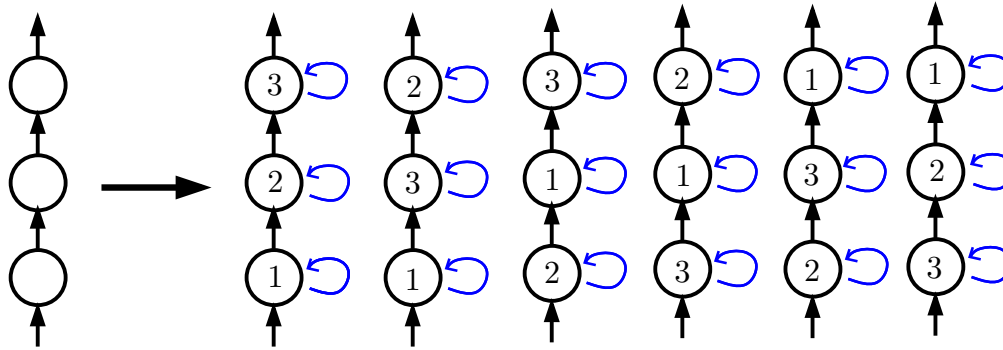
where $c_i(\sigma) = \#\{\text{cycles length } i \text{ in } \sigma\}$

- Cycle index sum of \mathcal{A} (cf Pólya) is the **multivariate** series

$$\begin{aligned} Z_{\mathcal{A}}(s_1, s_2, \dots) &= \sum_{\sigma \cdot A = A} W(\sigma, A) \\ &= \sum_{n \geq 1} \frac{1}{n!} \sum_{\sigma \in \mathfrak{S}_n} s_1^{c_1} \dots s_n^{c_n} \#(\text{Fix}_{\sigma}) \end{aligned}$$

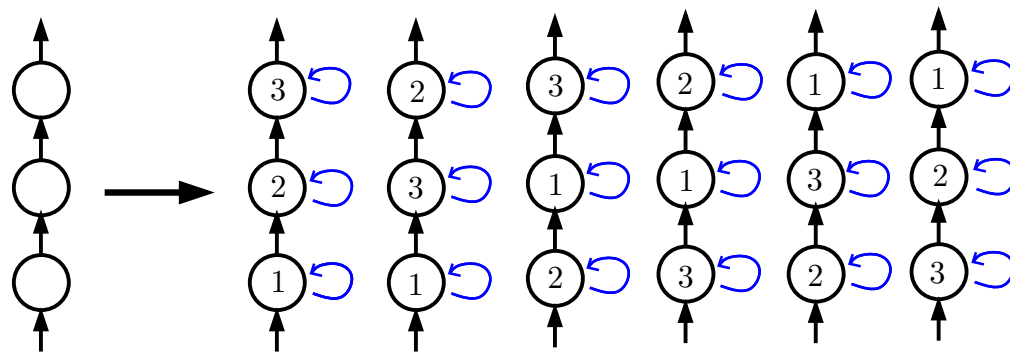
- OGF of $\tilde{\mathcal{A}} = \text{EGF of } \text{Sym}(\mathcal{A}) = Z_{\mathcal{A}}(x, x^2, x^3, \dots)$

Examples of cycle index sums

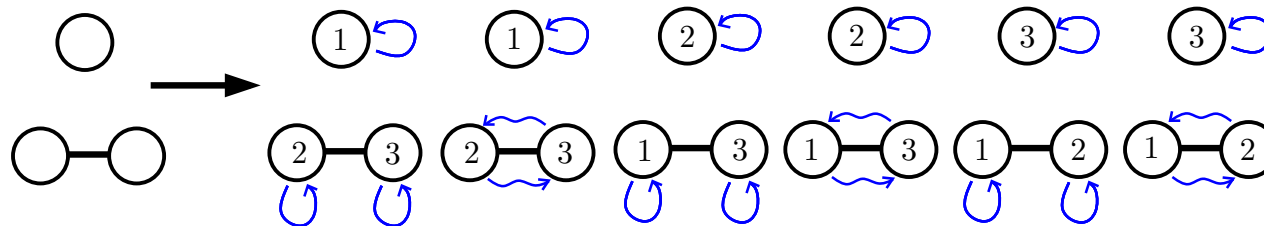


$$\begin{aligned}
 Z &= \frac{1}{6} s_1^3 + \frac{1}{6} s_1^3 + \frac{1}{6} s_1^3 + \frac{1}{6} s_1^3 + \frac{1}{6} s_1^3 + \frac{1}{6} s_1^3 \\
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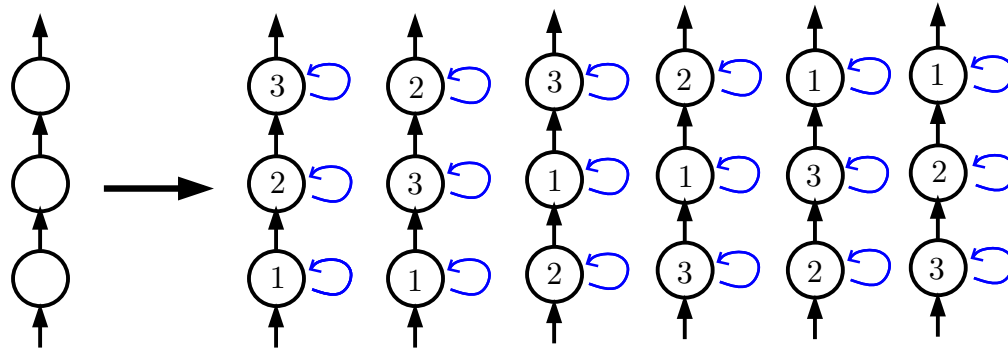


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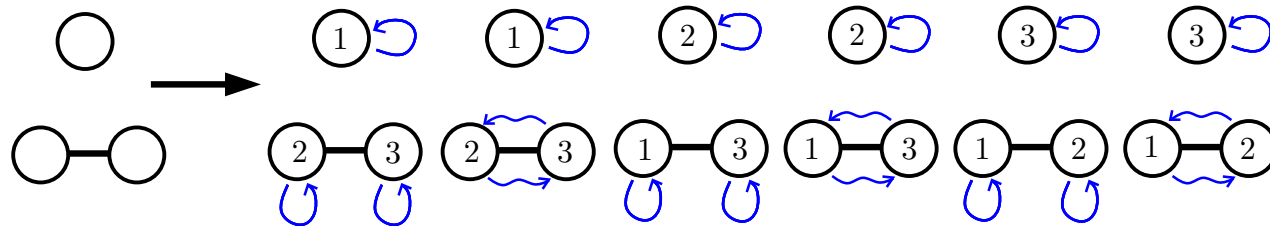
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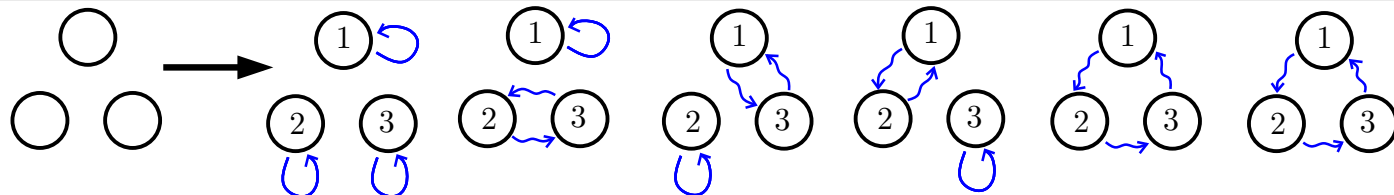
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Dictionary for OGF

- **Unlabelled** class $\tilde{\mathcal{C}} = \cup_n \mathcal{C}_n / \mathfrak{S}_n$ $\tilde{c}_n = \text{Card}(\tilde{\mathcal{C}}_n)$

$$\text{OGF : } \tilde{C}(x) = \sum_{n \geq 0} \tilde{c}_n x^n$$

- **Dictionary** (computation rules):

Disjoint union	$\mathcal{C} = \mathcal{A} + \mathcal{B}$	$\tilde{C}(x) = \tilde{A}(x) + \tilde{B}(x)$
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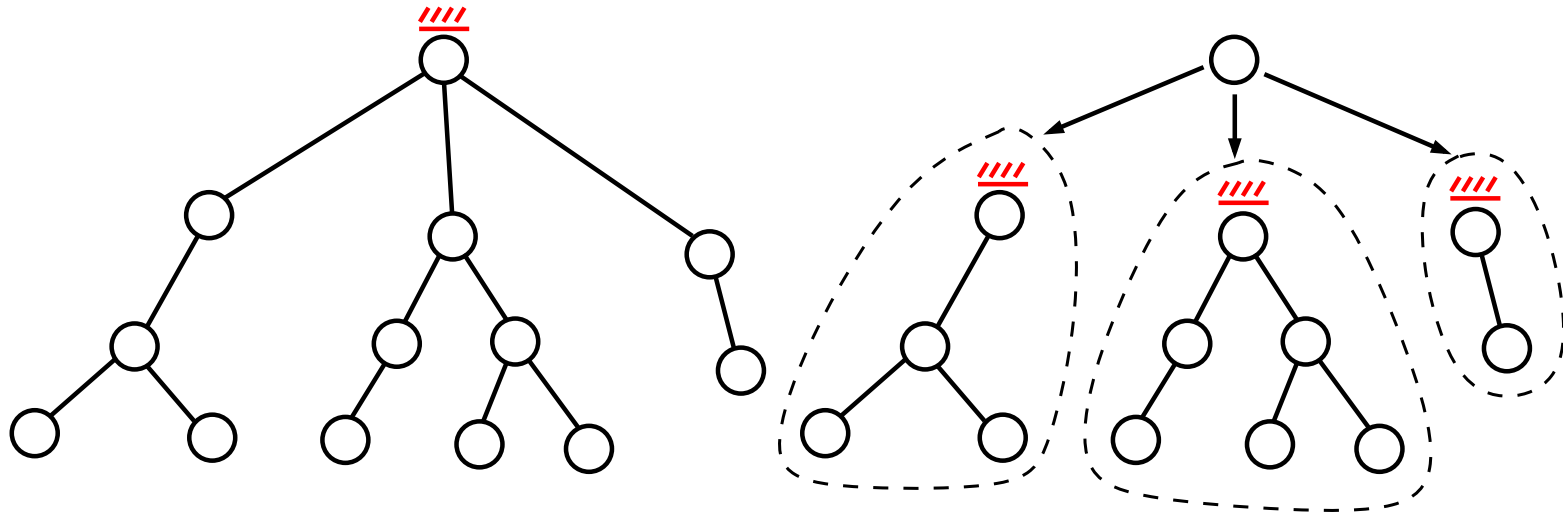
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- **Remark:** $\text{Set}(\mathcal{A}) = \text{Set} \circ \mathcal{A}$, computation rule for \circ implies the one for Set using $Z_{\text{Set}} = \exp(\sum_{i \geq 1} s_i / i)$

Example: rooted trees

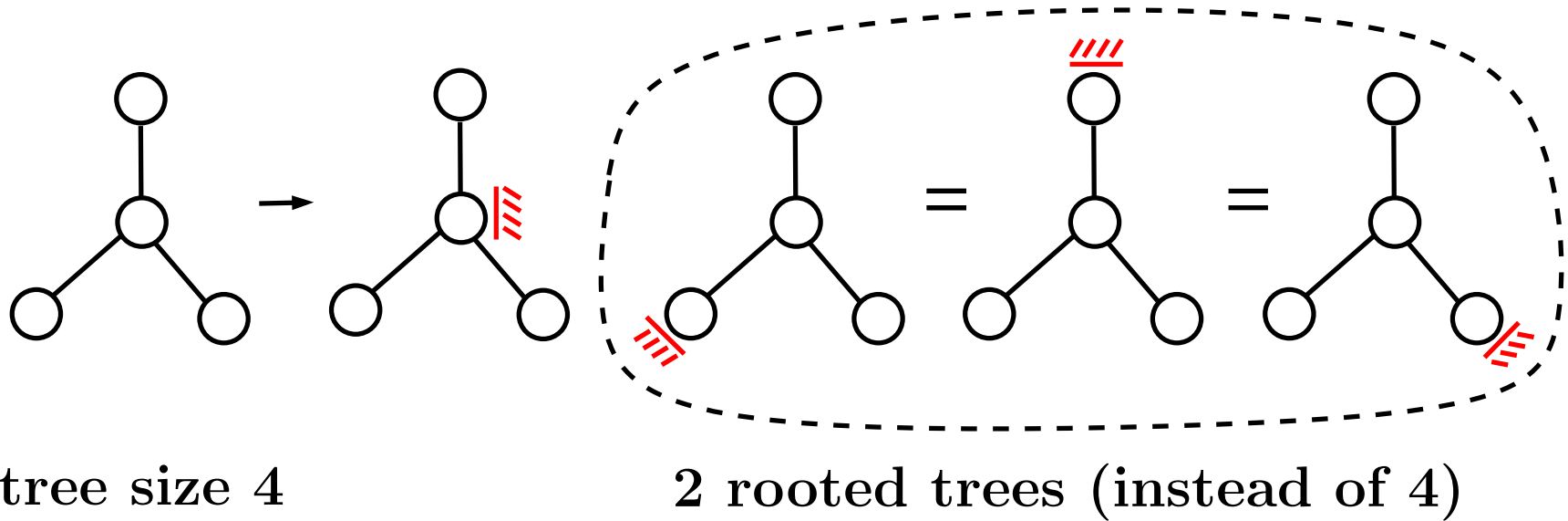
Decomposition at the root:



$$\mathcal{R} = \mathcal{Z} \times \text{Set}(\mathcal{R}) \Rightarrow R(z) = z \exp \left(\sum_{i \geq 1} \frac{1}{i} R(z^i) \right)$$

\Rightarrow **recurrence** formula for $[z^n]R(z)$

Count rooted $\not\Rightarrow$ count unrooted

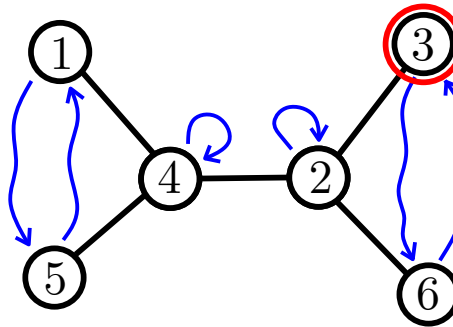


In general $n \cdot a_n^{\text{unrooted}} > a_n^{\text{rooted}}$ (symmetries)

Question: $n \cdot a_n^{\text{unrooted}} = ?$

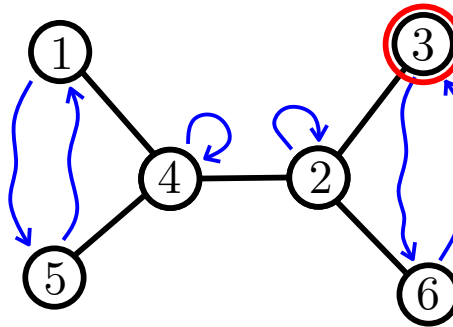
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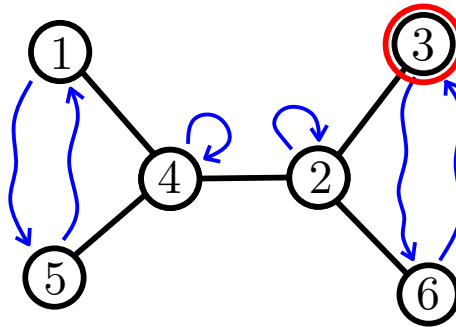


In size n we have

$$\begin{array}{ccc} \tilde{\mathcal{A}} & \xrightarrow{\times n!} & \text{Sym}(\mathcal{A}) \\ & & \downarrow \times n \\ & & (\text{Sym}(\mathcal{A}))^\bullet \end{array}$$

Pointing symmetries

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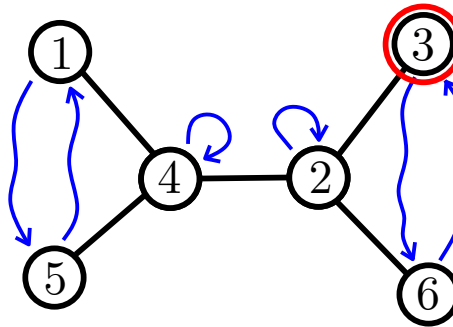
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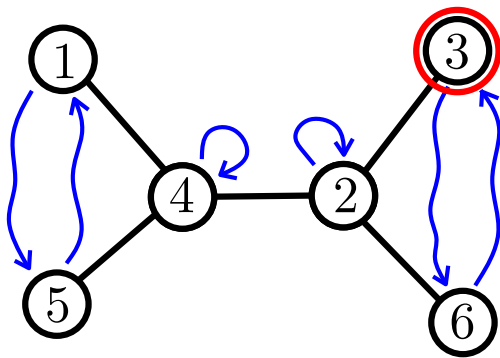
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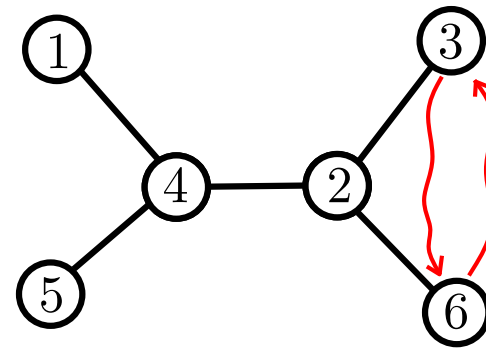
(rk: $\text{Sym}(\mathcal{A}^\bullet) \subseteq (\text{Sym}(\mathcal{A}))^\bullet$)

Cycle-pointed structures

Definition: Cycle-pointed structure = structure A + cycle c such that there exists (at least) one automorphism of A having c as one of its cycles.



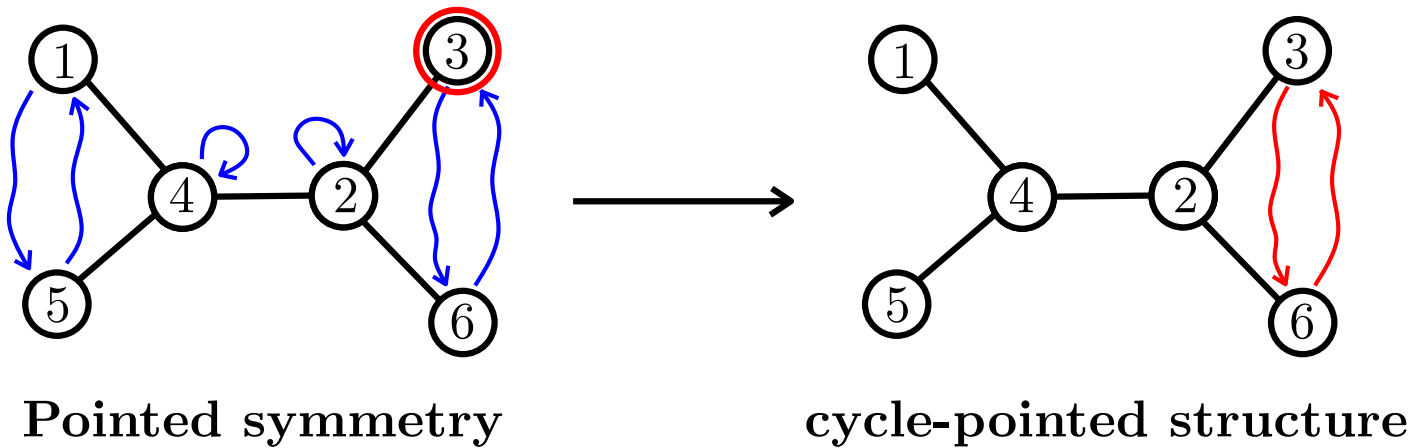
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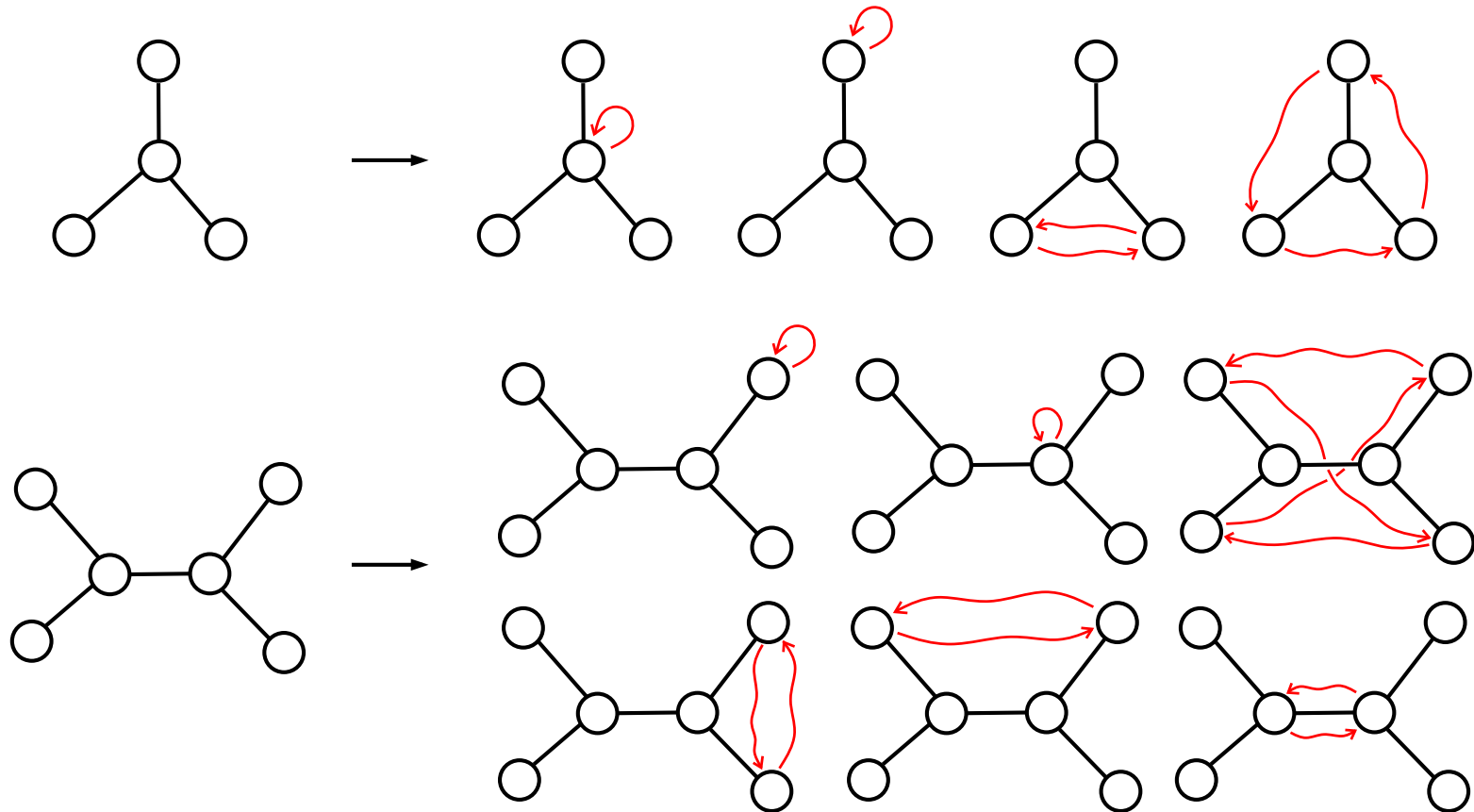
Let $\mathcal{A}^\circ = \{\text{cycle - pointed structures from } \mathcal{A}\}$. Then

$$\text{Sym}(\mathcal{A}^\circ) \simeq (\text{Sym}(\mathcal{A}))^\bullet$$

Cycle-pointing is unbiased

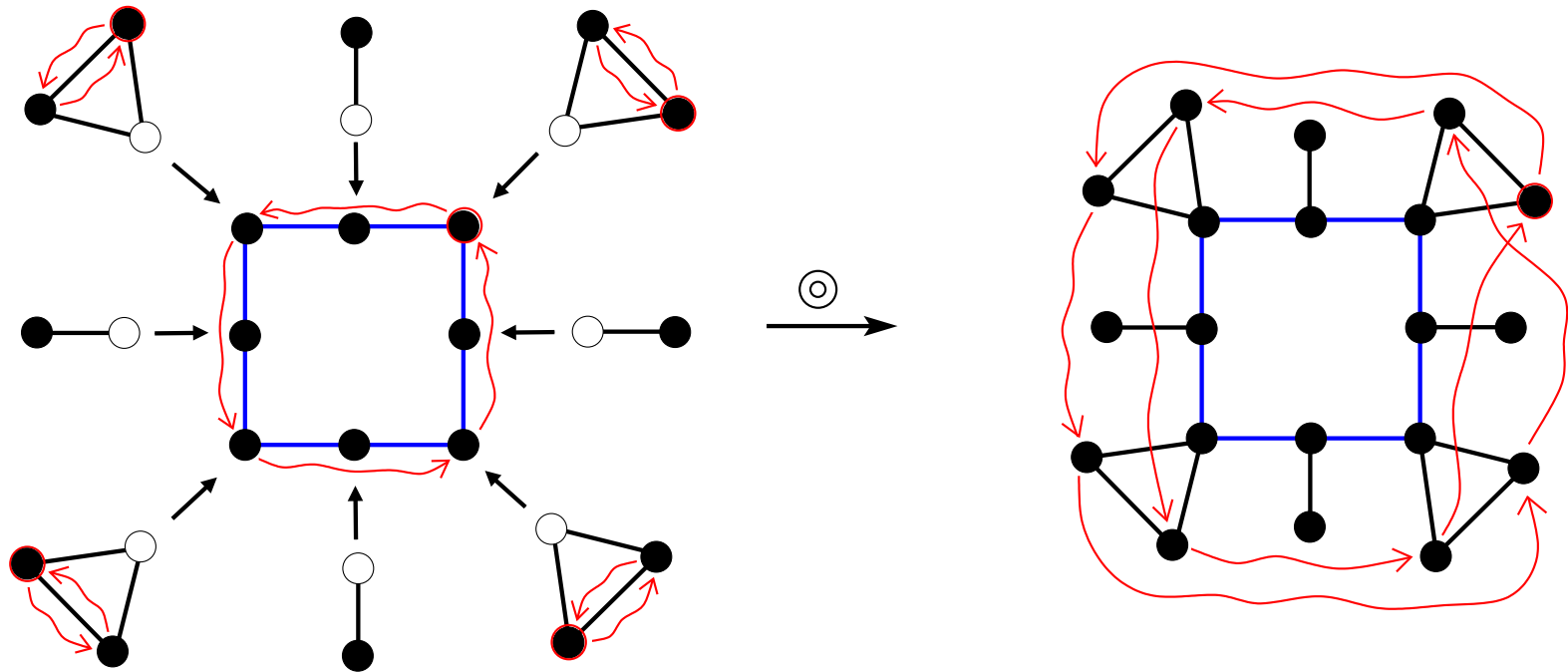
Theorem: An unlabelled structure of size n gives rise to n unlabelled cycle-pointed structures (cf Parker's lemma).

$$\widetilde{\mathcal{A}}_n^\circ \cong n \times \widetilde{\mathcal{A}}_n$$



Pointing the classical constructions

- $(\mathcal{A} + \mathcal{B})^\circ = \mathcal{A}^\circ + \mathcal{B}^\circ$
- $(\mathcal{A} \times \mathcal{B})^\circ = \mathcal{A}^\circ \times \mathcal{B} + \mathcal{A} \times \mathcal{B}^\circ$
- $(\mathcal{A} \circ \mathcal{B})^\circ = \mathcal{A}^\circ \odot \mathcal{B}$



Application: counting trees (1)

Decomposition of cycle-pointed trees (3 lines)

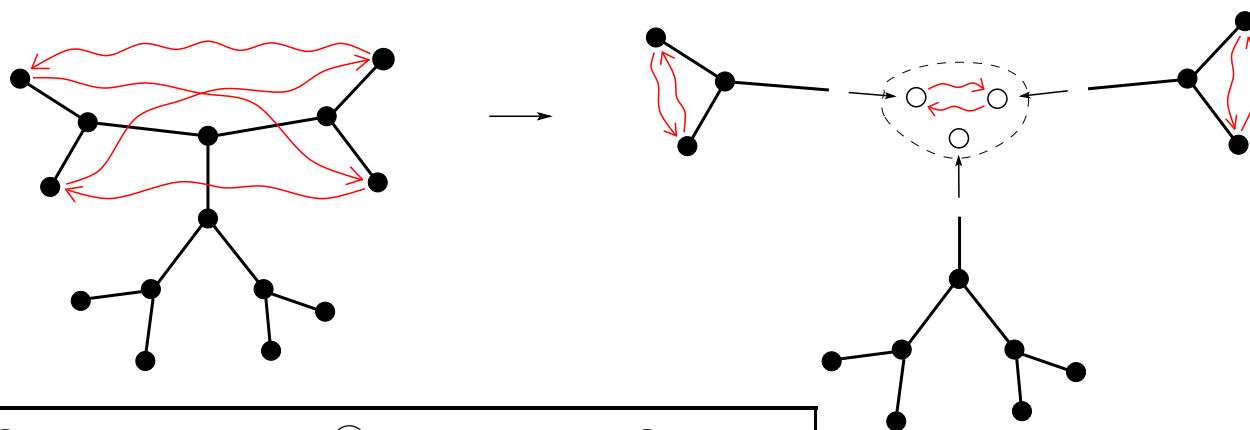
- Dichotomy: pointed cycle has length 1 or ≥ 2 :

$$1) \mathcal{T}^\circ = \mathcal{R} + \mathcal{T}^*$$

- Rooted trees (\mathcal{R}) are decomposed at the root

$$2) \mathcal{R} = \mathcal{Z} \times \text{Set}(\mathcal{R})$$

- Symmetric cycle-pointed trees (\mathcal{T}^*) are decomposed at a centre of symmetry.



$$3) \mathcal{T}^* = \mathcal{Z} \times \text{Set}^* \odot \mathcal{R} + \mathcal{L}^* \odot \mathcal{R}$$

Application: counting trees (2)

$$\begin{cases} T^\circ &= \mathcal{R} + T^* \\ \mathcal{R} &= \mathcal{Z} \times \text{Set}(\mathcal{R}) \\ T^* &= \mathcal{Z} \times \text{Set}^* \odot (\mathcal{R}) + \mathcal{L}^* \odot (\mathcal{R}) \end{cases}$$

↓ translate to equation system
(dictionary rules+Pólya operators)

$$\begin{cases} R(x) &= x \exp \left(\sum_{k \geq 1} \frac{1}{k} R(x^k) \right) \\ xt'(x) &= R(x) + x^2 R'(x^2) + \sum_{i \geq 2} x^i R'(x^i) R(x) \end{cases}$$

↓ extract coefficients

n	1	2	3	4	5	6	7	8	9
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Application: counting trees (2)

$$\begin{cases} \mathcal{T}^\circ & = \mathcal{R} + \mathcal{T}^* \\ \mathcal{R} & = \mathcal{Z} \times \text{Set}(\mathcal{R}) \\ \mathcal{T}^* & = \mathcal{Z} \times \text{Set}^* \odot (\mathcal{R}) + \mathcal{L}^* \odot (\mathcal{R}) \end{cases}$$

↓ translate to equation system
(dictionary rules+Pólya operators)

$$\begin{cases} R(x) & = x \exp \left(\sum_{k \geq 1} \frac{1}{k} R(x^k) \right) \\ xt'(x) & = R(x) + x^2 R'(x^2) + \sum_{i \geq 2} x^i R'(x^i) R(x) \\ & = xR'(x)(1 - R(x)) + x^2 R'(x^2) \end{cases}$$

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Exact counting results

Theorem: [Bergeron, Labelle, Leroux],

For any class \mathcal{A} **decomposable** in terms of

- basic classes $\{\mathbf{1}, \mathcal{Z}, \text{Seq}, \text{Cyc}, \text{Set}\}$,
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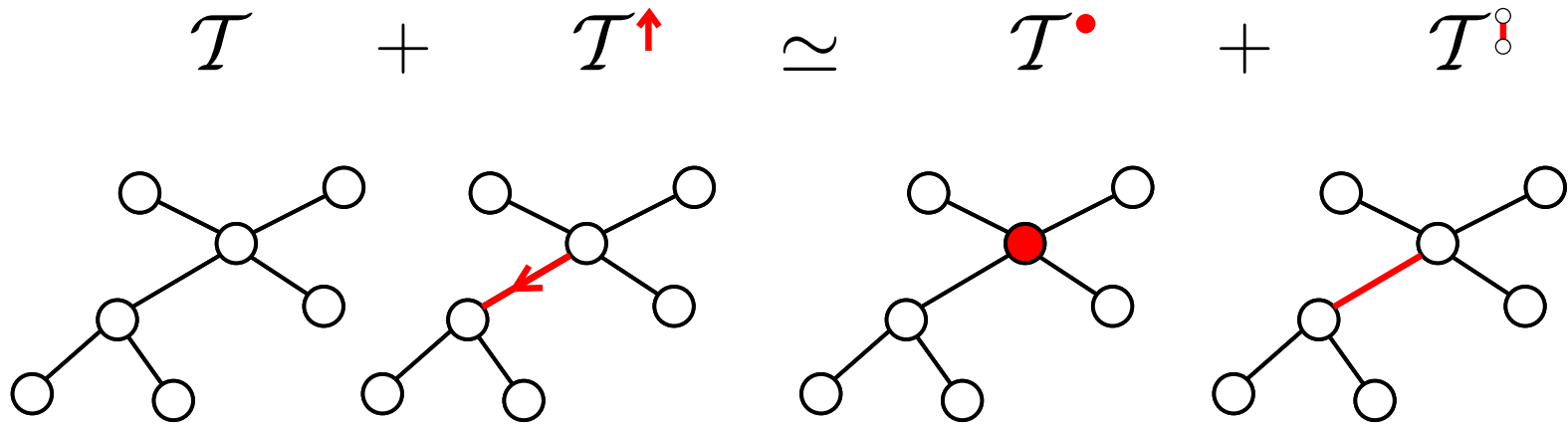
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(includes tree families, outerplanar graphs,...)

Another approach

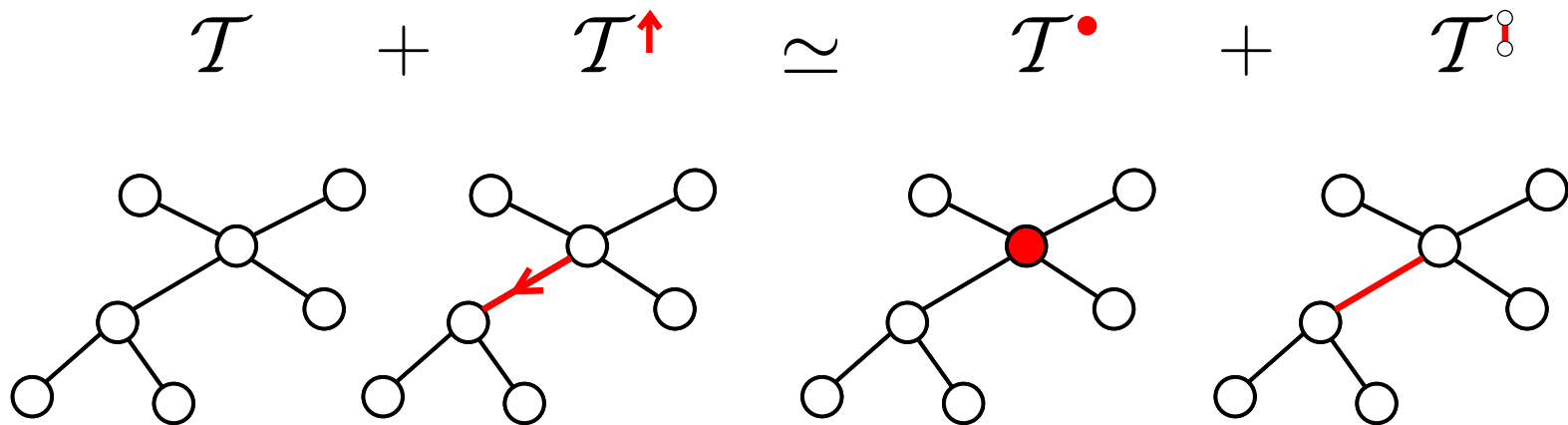
- Dissimilarity characteristic formula (Otter)
- Dissymmetry theorem (Robinson, Leroux):



$$\mathbf{t}(\mathbf{x}) = \mathbf{t}^{\bullet}(\mathbf{x}) + \mathbf{t}^{\circ}(\mathbf{x}) - \mathbf{t}^{\uparrow}(\mathbf{x})$$

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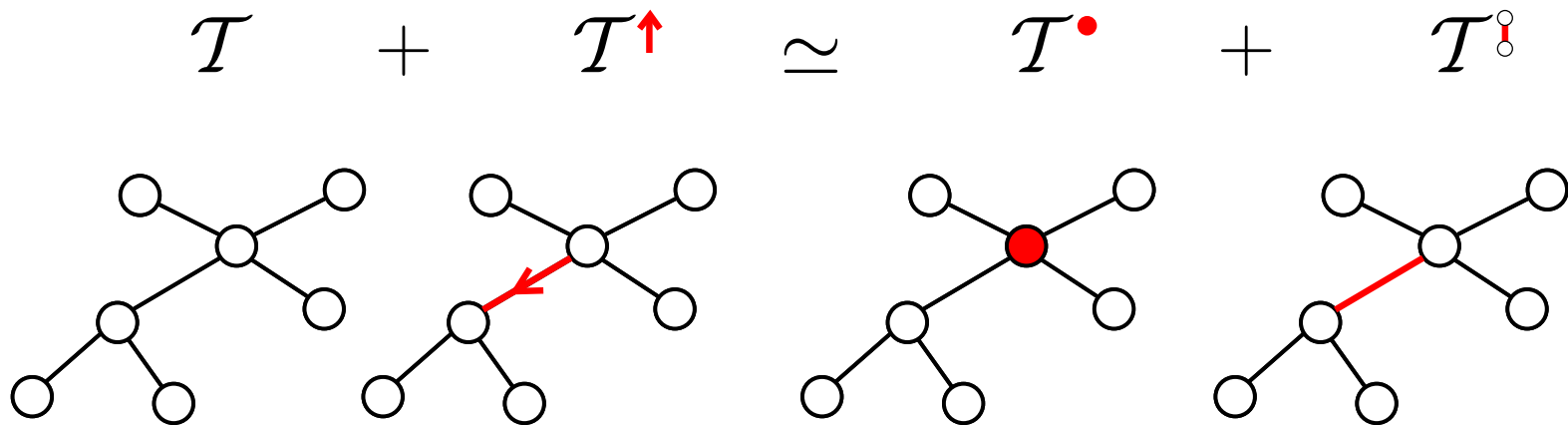


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Application to asymptotic enumeration

Asymptotic scheme

Main result: “universality” of asymptotic behaviour

$$|\widetilde{\mathcal{A}}_n| \sim c \gamma^n n^{-5/2}$$

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Scheme:

- Decompose cycle-pointed class \mathcal{A}°
 \Rightarrow Equation for $\widetilde{A}^\circ(x)$
- Drmota-Lalley-Woods $\Rightarrow \widetilde{A}^\circ(x)$ has square-root sing.
- Transfer theorem [Flajolet-Odlyzko] $\Rightarrow |\widetilde{\mathcal{A}}_n^\circ| \sim c \gamma^n n^{-3/2}$
- Pointing relation: $|\widetilde{\mathcal{A}}_n| = \frac{1}{n} |\widetilde{\mathcal{A}}_n^\circ| \sim c \gamma^n n^{-5/2}$

Illustration on trees

- Rooted labelled trees: $y = L(z)$ satisfies

$$y = z \exp(y)$$

Inverse is $g(y) = y \exp(-y)$,

$$g'(y) = 0 \Rightarrow y = 1 \Rightarrow z = 1/e \Rightarrow L(z) = 1 - c\sqrt{1 - ze} + \dots$$

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$$zt'(z) = \underbrace{z^2 R'(z^2)}_{\text{analytic at } \rho} + \underbrace{\left(1 + \sum_{i \geq 2} z^i R'(z^i)\right)}_{G(z) \text{ analytic at } \rho} \cdot R(z)$$

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(Rk: $\mathbb{E}_n(\#\text{dissimilar vertices in tree}) \sim n/G(\rho)$)

Asymptotic using dissym. theorem

$$t(z) = R(z) - \frac{1}{2}(R(z)^2 - R(z^2))$$

Square-root expansion: $R(z) = 1 - * \sqrt{1 - z/\rho} + \dots$

\Rightarrow square-root terms cancel out, $t(z)$ “ \leq ” $(1 - z/\rho)^{3/2}$

But $zt'(z) \geq R(z)$, so $t(z)$ “ \geq ” $(1 - z/\rho)^{3/2}$

Hence (transfer theorem):

$$[z^n]t(z) \sim c \rho^{-n} n^{-5/2}$$

(Rk: Cancellation proof uneasy if “big” functional equation)

Application to random generation

General methods

Two sampling methods giving uniform distribution

- **Recursive method** (Nijenhuis-Wilf'78):

$$\mathbb{P}(\gamma \in \mathcal{C}_n) = \frac{1}{c_n} \quad [\text{Fixed size}]$$

- Based on **coefficients**:

$$\mathcal{C} = \mathcal{A} + \mathcal{B} \Rightarrow \mathbb{P}(\gamma \in \mathcal{A}_n) = \frac{a_n}{c_n}$$

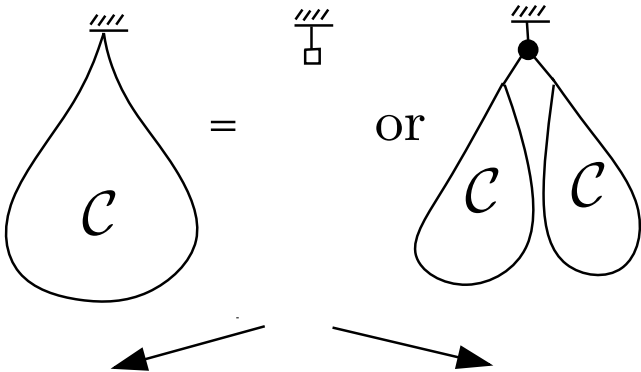
- **Boltzmann samplers** (Duchon, Flajolet, Louchard, Schaeffer'02)

$$\mathbb{P}(\gamma \in \mathcal{C}) = \frac{x^{|\gamma|}}{C(x)} \quad [\text{Whole class}]$$

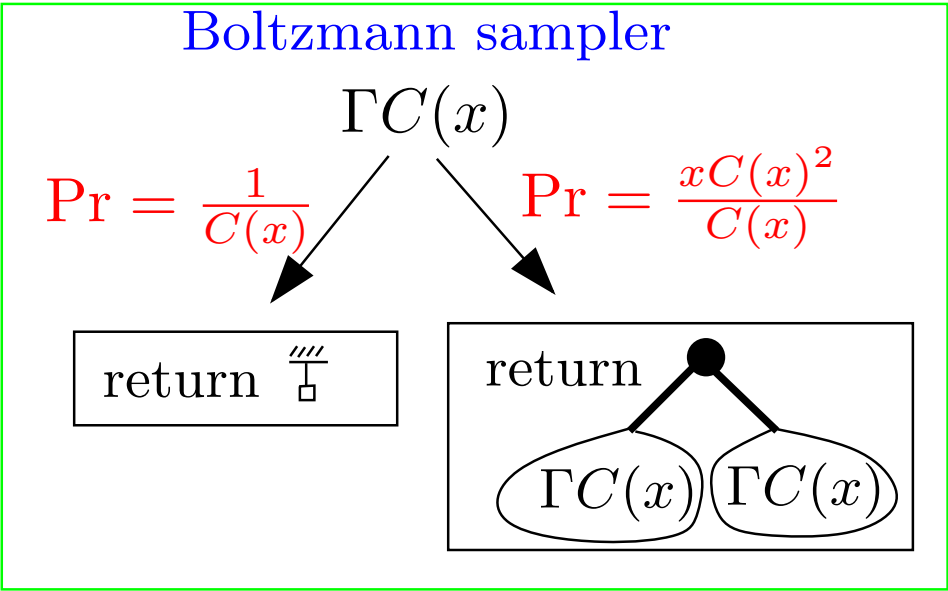
- Based on **gen. funct.**

$$\mathcal{C} = \mathcal{A} + \mathcal{B} \Rightarrow \mathbb{P}(\gamma \in \mathcal{A}) = \frac{A(x)}{C(x)}$$

Boltzmann samplers: example



Generating function
 $C(x) = 1 + xC(x)^2$



Results

Theorem: [Duchon et al'02], [Flajolet et al'07]

For any class \mathcal{A} decomposable in terms of

- basic classes $\{\mathbf{1}, \mathcal{Z}\}$
- constructions $\{+, \times, \text{Seq}, \text{Cyc}, \text{Set}\}$,

there is a linear-time Boltzmann sampler $\Gamma_{\tilde{\mathcal{A}}}(x)$.

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Pólya-Boltzmann samplers

- Ordinary Boltzmann samplers:

$$\tilde{A}(x) = \sum_{\gamma \in \tilde{\mathcal{A}}} x^{|\gamma|} \Rightarrow \mathbb{P}(\gamma) = \frac{x^{|\gamma|}}{\tilde{A}(x)}$$

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- Recover ordinary Boltzmann sampler using specialization

$$Z_{\mathcal{C}}(x, x^2, \dots) = C(x) \Rightarrow \Gamma Z_{\mathcal{C}}(x, x^2, \dots) = \Gamma \tilde{C}(x)$$

Sampler for trees

Let $t(x)$ be the OGF of (unrooted) trees.

1) Translate the equation

$$R(x) = \exp\left(\sum_{i \geq 1} \frac{1}{i} R(x^i)\right)$$

into a **Boltzmann sampler for \mathcal{R}** cf [Flajolet et al'07]
(superposition of Poisson laws)

2) Translate the equation

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(Also recursive sampler by Wilf using “centre of gravity”)

Open problems

- Refined **complexity analysis** of Pólya-Boltzmann samplers, cf [Pivoteau, Salvy, Soria'08]
- Boltzmann sampling with **catalytic variables**
- For which recursive specification can we **determine the growth rate** automatically ?
- Use Boltzmann samplers to **study random unlabelled structures**, cf [Panagiotou, Steger].