

# On the Structure of Compatible Rational Functions

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# Outline

1. Bivariate compatible rational functions
2. General case
3. A structure theorem
4. Multiplicative decomposition of hyperexp.-hypergeom. elements
5. Applications
6. Algebraic dependence of hyperexp.-hypergeom. elements

**Note.**  $\mathbb{F}$  is an algebraically closed field of characteristic zero.

# Bivariate case 1: differential and difference

Let

$$\delta : \mathbb{F}(t, x) \rightarrow \mathbb{F}(t, x)$$

$$\sigma : \mathbb{F}(t, x) \rightarrow \mathbb{F}(t, x)$$

and

$$f(t, x) \mapsto \frac{\partial f(t, x)}{\partial t}$$

$$f(t, x) \mapsto f(t, x + 1)$$

A first-order system

$$\begin{cases} \delta(z) = u z, \\ \sigma(z) = v z, \end{cases} \quad \text{where } u, v \in \mathbb{F}(t, x) \text{ with } v \neq 0,$$

has a nonzero solution iff the following **compatibility condition (CC)** holds:

$$\frac{\delta(v)}{v} = \sigma(u) - u.$$

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$$\begin{cases} \sigma \circ \delta(z) = \sigma(u) v z \\ \delta \circ \sigma(z) = (\delta(v) + uv) z \end{cases} \implies \text{the CC.}$$

**Definition.**  $u, v \in \mathbb{F}(t, x)$  are **compatible** w.r.t.  $\{\delta, \sigma\}$  if  $v \neq 0$  and

$$\frac{\delta(v)}{v} = \sigma(u) - u.$$

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**Lemma.** (R. Feng, M. Singer, and M. Wu, 2010)

$u, v \in \mathbb{F}(t, x)$  are compatible w.r.t.  $\{\delta, \sigma\}$  iff  $\exists f \in \mathbb{F}(t, x), \alpha, \beta \in \mathbb{F}(t), \lambda \in \mathbb{F}(x)$  s.t.

$$u = \frac{\delta(f(t, x))}{f(t, x)} + x \frac{\delta(\alpha(t))}{\alpha(t)} + \beta(t)$$

and

$$v = \frac{\sigma(f(t, x))}{f(t, x)} \alpha(t) \lambda(x).$$

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**Corollary.** A solution of  $\{\delta(z) = u z, \sigma(z) = v z\}$  can be written as

$$f(t, x) \alpha(t)^x \exp\left(\int \beta(t) dt\right) T(x),$$

where  $\delta(T) = 0$  and  $\sigma(T) = \lambda T$ .

## Bivariate case 2: differential and q-difference

Let

$$\delta : \mathbb{F}(t, y) \rightarrow \mathbb{F}(t, y)$$

$$\tau : \mathbb{F}(t, y) \rightarrow \mathbb{F}(t, y)$$

and

$$f(t, y) \mapsto \frac{\partial f(t, y)}{\partial t}$$

$$f(t, y) \mapsto f(t, qy),$$

where  $q \in \mathbb{F}$  is not a root of unity.

A first-order system:

$$\begin{cases} \delta(z) = uz, \\ \tau(z) = wz, \end{cases} \quad u, w \in \mathbb{F}(t, y) \text{ with } w \neq 0,$$

has a nonzero solution iff the following **CC** holds:

$$\frac{\delta(w)}{w} = \tau(u) - u.$$

**Definition.**  $u, w \in \mathbb{F}(t, y)$  are **compatible** w.r.t.  $\{\delta, \tau\}$  if  $w \neq 0$  and

$$\frac{\delta(w)}{w} = \tau(u) - u.$$

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**Lemma.** (a  $q$ -analogue of FSW Lemma)

$u, w \in \mathbb{F}(t, y)$  are compatible w.r.t.  $\{\delta, \tau\}$  iff  $\exists f \in \mathbb{F}(t, y), \beta \in \mathbb{F}(t), \mu \in \mathbb{F}(y)$  s.t.

$$u = \frac{\delta(f(t, y))}{f(t, y)} + \beta(t) \quad \text{and} \quad v = \frac{\tau(f(t, y))}{f(t, y)} \mu(y).$$

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**Corollary.** A solution of  $\{\delta(z) = u z, \tau(z) = w z\}$  can be written as

$$f(t, y) \exp\left(\int \beta(t) dt\right) Q(y),$$

where  $\delta(Q) = 0$  and  $\tau(Q) = \mu Q$ .

# Comparison

## Differential and difference.

**Fact.** For  $g \in \mathbb{F}(t, x)$ ,

$$\underline{\sigma(g) - g \in \mathbb{F}(t) \Leftrightarrow g \in \mathbb{F}(t)[x] \text{ and } \deg_x g \leq 1.}$$

⇓

$$u = \frac{\delta(f(t, x))}{f(t, x)} + x \frac{\delta(\alpha(t))}{\alpha(t)} + \beta(t) \quad \text{and} \quad v = \frac{\sigma(f(t, x))}{f(t, x)} \alpha(t) \lambda(x).$$

## Differential and $q$ -difference.

**Fact.** For  $g \in \mathbb{F}(t, y)$ ,

$$\underline{\tau(g) - g \in \mathbb{F}(t) \Leftrightarrow g \in \mathbb{F}(t).}$$

⇓

$$u = \frac{\delta(f(t, y))}{f(t, y)} + \beta(t) \quad \text{and} \quad w = \frac{\tau(f(t, y))}{f(t, y)} \mu(y).$$

## Bivariate Case III: difference and q-difference

Let

$$\begin{aligned} \sigma : \mathbb{F}(x, y) &\rightarrow \mathbb{F}(x, y) & \tau : \mathbb{F}(x, y) &\rightarrow \mathbb{F}(x, y) \\ f(x, y) &\mapsto f(x + 1, y) & \text{and} & \\ f(x, y) &\mapsto f(x, qy), & & \end{aligned}$$

where  $q \in \mathbb{F}$  is not a root of unity.

For  $v, w \in \mathbb{F}(x, y)$  with  $vw \neq 0$ , the system:

$$\begin{cases} \sigma(z) = v z \\ \tau(z) = w z \end{cases}$$

has a nonzero solution iff the following **CC** holds:

$$\frac{\sigma(w)}{w} = \frac{\tau(v)}{v}.$$

**Definition.**  $v, w \in \mathbb{F}(t, x)$  are **compatible** w.r.t.  $\{\sigma, \tau\}$  if  $vw \neq 0$  and

$$\frac{\sigma(w)}{w} = \frac{\tau(v)}{v}.$$

**Lemma.**  $v$  and  $w$  are compatible w.r.t.  $\{\sigma, \tau\}$  iff there exist  $f \in \mathbb{F}(x, y)$ ,  $\lambda \in \mathbb{F}(x)$ ,  $\mu \in \mathbb{F}(y)$  s.t.

$$v = \frac{\sigma(f(x, y))}{f(x, y)} \lambda(x) \quad \text{and} \quad w = \frac{\tau(f(x, y))}{f(x, y)} \mu(y).$$

**Corollary.** A solution of

$$\begin{cases} \sigma(z) = v z \\ \tau(z) = w z \end{cases}$$

can be written as

$$c f(x, y) T(x) Q(y),$$

where  $c \in \mathbb{F}$ ,  $\sigma(T) = \lambda T$  and  $\tau(Q) = \mu Q$ .

# Survey on mixed cases

| Operators  | CC  | Rational solutions  |
|--|---|---|
| $\delta = \frac{\partial}{\partial t}$<br>$\sigma : x \mapsto x + 1$ | $\frac{\delta(v)}{v} = \sigma(u) - u$     | $u = \frac{\delta(f)}{f} + x \frac{\delta(\alpha)}{\alpha} + \beta, \quad v = \frac{\sigma(f)}{f} \alpha \lambda$<br>where $f \in \mathbb{F}(t, x), \alpha, \beta \in \mathbb{F}(t), \lambda \in \mathbb{F}(x)$ |
| $\delta = \frac{\partial}{\partial t},$<br>$\tau : y \mapsto q y$    | $\frac{\delta(w)}{w} = \tau(u) - u$       | $u = \frac{\delta(f)}{f} + \beta, \quad w = \frac{\tau(f)}{f} \mu$<br>where $f \in \mathbb{F}(t, y), \beta \in \mathbb{F}(t), \mu \in \mathbb{F}(y)$  |
| $\sigma : x \mapsto x + 1,$<br>$\tau : y \mapsto q y$                | $\frac{\sigma(w)}{w} = \frac{\tau(v)}{v}$ | $v = \frac{\sigma(f)}{f} \lambda, \quad w = \frac{\tau(f)}{f} \mu,$<br>where $f \in \mathbb{F}(x, y), \lambda \in \mathbb{F}(x), \mu \in \mathbb{F}(y).$  |

# General case

Let  $\mathbf{t} = (t_1, \dots, t_\ell)$ ,  $\mathbf{x} = (x_1, \dots, x_m)$  and  $\mathbf{y} = (y_1, \dots, y_n)$ .

For  $i = 1, \dots, \ell$ ,

$$\delta_i : \mathbb{F}(\mathbf{t}, \mathbf{x}, \mathbf{y}) \longrightarrow \mathbb{F}(\mathbf{t}, \mathbf{x}, \mathbf{y})$$

$$f \mapsto \frac{\partial f}{\partial t_i}.$$

---

For  $j = 1, \dots, m$ ,

$$\sigma_j : \mathbb{F}(\mathbf{t}, \mathbf{x}, \mathbf{y}) \longrightarrow \mathbb{F}(\mathbf{t}, \mathbf{x}, \mathbf{y})$$

$$f(\mathbf{t}, x_1, \dots, x_m, \mathbf{y}) \mapsto f(\mathbf{t}, x_1, \dots, x_{j-1}, \mathbf{x}_j + \mathbf{1}, x_{j+1}, \dots, x_m, \mathbf{y}).$$

---

For  $k = 1, \dots, n$ ,

$$\tau_k : \mathbb{F}(\mathbf{t}, \mathbf{x}, \mathbf{y}) \longrightarrow \mathbb{F}(\mathbf{t}, \mathbf{x}, \mathbf{y})$$

$$f(\mathbf{t}, \mathbf{x}, y_1, \dots, y_n) \mapsto f(\mathbf{t}, \mathbf{x}, y_1, \dots, y_{k-1}, q_k y_k, y_{k+1}, \dots, y_n),$$

where  $q_1, \dots, q_n \in \mathbb{F}^\times$  s.t.  $q_1^{e_1} \cdots q_n^{e_n} \neq 1$  unless  $e_1 = \dots = e_n = 0$ .

# Compatible rational functions

$$\Delta = \{\delta_1, \dots, \delta_\ell, \sigma_1, \dots, \sigma_m, \tau_1, \dots, \tau_n\}.$$

**Definition.** A sequence of rational functions

$$u_1, \dots, u_\ell, v_1, \dots, v_m, w_1, \dots, w_n$$

is **compatible w.r.t.  $\Delta$**  if

$$v_1 \cdots v_m w_1 \cdots w_n \neq 0,$$

and the following six sets of CC's hold:

| Unmixed  | Mixed   |
|--|---|
| $\delta_i(u_j) = \delta_j(u_i), \quad 1 \leq i < j \leq \ell$                      | $\frac{\delta_i(v_j)}{v_j} = \sigma_j(u_i) - u_i, \quad i = 1, \dots, \ell, \quad j = 1, \dots, m$  |
| $\frac{\sigma_i(v_j)}{v_j} = \frac{\sigma_j(v_i)}{v_i}, \quad 1 \leq i < j \leq m$ | $\frac{\delta_i(w_k)}{w_k} = \tau_k(u_i) - u_i, \quad i = 1, \dots, \ell, \quad k = 1, \dots, n$    |
| $\frac{\tau_i(w_j)}{w_j} = \frac{\tau_j(w_i)}{w_i}, \quad 1 \leq i < j \leq n$     | $\frac{\sigma_j(w_k)}{w_k} = \frac{\tau_k(v_j)}{v_j}, \quad j = 1, \dots, m, \quad k = 1, \dots, n$ |

## A structure theorem

**Theorem.** A sequence of rational functions  $u_1, \dots, u_\ell, v_1, \dots, v_m, w_1, \dots, w_n$  is compatible iff, for all  $i, j, k$  with  $1 \leq i \leq \ell$ ,  $1 \leq j \leq m$  and  $1 \leq k \leq n$ ,

$$u_i = \frac{\delta_i(f(\mathbf{t}, \mathbf{x}, \mathbf{y}))}{f(\mathbf{t}, \mathbf{x}, \mathbf{y})} + \sum_{j=1}^m x_j \frac{\delta_i(\alpha_j(\mathbf{t}))}{\alpha_j(\mathbf{t})} + \beta_i(\mathbf{t}),$$

$$v_j = \frac{\sigma_j(f(\mathbf{t}, \mathbf{x}, \mathbf{y}))}{f(\mathbf{t}, \mathbf{x}, \mathbf{y})} \alpha_j(\mathbf{t}) \lambda_j(\mathbf{x}),$$

$$w_k = \frac{\tau_k(f(\mathbf{t}, \mathbf{x}, \mathbf{y}))}{f(\mathbf{t}, \mathbf{x}, \mathbf{y})} \mu_k(\mathbf{y}),$$

where  $f \in \mathbb{F}(\mathbf{t}, \mathbf{x}, \mathbf{y})$ ,  $\alpha_j, \beta_i \in \mathbb{F}(\mathbf{t})$ ,  $\lambda_j \in \mathbb{F}(\mathbf{x})$ ,  $\mu_k \in \mathbb{F}(\mathbf{y})$ ,

$\beta_1(\mathbf{t}), \dots, \beta_\ell(\mathbf{t})$  are compatible w.r.t.  $\{\delta_1, \dots, \delta_\ell\}$ ,

$\lambda_1(\mathbf{x}), \dots, \lambda_m(\mathbf{x})$  are compatible w.r.t.  $\{\sigma_1, \dots, \sigma_m\}$ , and

$\mu_1(\mathbf{y}), \dots, \mu_n(\mathbf{y})$  are compatible w.r.t.  $\{\tau_1, \dots, \tau_n\}$ .

# Survey on unmixed cases

| Systems   | CC's  | Structure   |
|---|---|---|
| $\begin{cases} \delta_1(z) = u_1 z \\ \vdots \\ \delta_\ell(z) = u_\ell z, \end{cases}$ $u_1, \dots, u_\ell \in \mathbb{F}(\mathbf{t})$ | $\delta_i(u_j) = \delta_j(u_i),$ $1 \leq i < j \leq \ell$                     | <p style="color: purple;">Christoper's theorem:</p> $u_i = \delta_i(g) + \sum_k c_k \frac{\delta_i(g_k)}{g_k},$ <p>where <math>c_k \in \mathbb{F}</math> and <math>f, g_k \in \mathbb{F}(\mathbf{t})</math></p>   |
| $\begin{cases} \sigma_1(z) = v_1 z \\ \vdots \\ \sigma_m(z) = v_m z, \end{cases}$ $v_1, \dots, v_m \in \mathbb{F}(\mathbf{x})$          | $\frac{\sigma_i(v_j)}{v_j} = \frac{\sigma_j(v_i)}{v_i}$ $1 \leq i < j \leq m$ | <p style="color: purple;">Ore-Sato's theorem:</p> $v_j = \frac{\sigma_j(g)}{g} \prod_{\mathbf{s} \in S} \prod_0^{s_j} r_{\mathbf{s}}(\mathbf{s} \cdot \mathbf{x} + k)$ <p>where <math>f \in \mathbb{F}(\mathbf{x})</math>, <math>S \subset \mathbb{Z}^m</math> and <math>r_{\mathbf{s}} \in \mathbb{F}(z)</math>.</p> |
| $\begin{cases} \tau_1(z) = w_1 z \\ \vdots \\ \tau_n(z) = w_n z, \end{cases}$ $w_1, \dots, w_n \in \mathbb{F}(\mathbf{y})$              | $\frac{\tau_i(w_j)}{w_j} = \frac{\tau_j(w_i)}{w_i}$ $1 \leq i < j \leq n$     | <p>A <math>q</math>-analogue<br/>by I. Gel'fand, M. Graev, and V. Retakh.<br/>when <math>q_1 = \dots = q_n</math></p>   |

# Hyperexponential-Hypergeometric elements

**Definition.** A first-order system

$$\{\delta_i(z) = u_i z, \sigma_j(z) = v_j z, \tau_k(z) = w_k z \mid i = 1, \dots, \ell, j = 1, \dots, m, k = 1, \dots, n\} \quad (*)$$

is **compatible** if  $u_1, \dots, u_\ell, v_1, \dots, v_m, w_1, \dots, w_n$  are compatible w.r.t.  $\Delta$ .

**Definition.** A nonzero solution of (\*) is called a **hyperexponential-hypergeometric element**, abbr. as *H*-element.

# Multiplicative decompositions

**Corollary.** An  $H$ -element can be written as

$$\text{a rational function} \times (\alpha_1(\mathbf{t})^{x_1} \cdots \alpha_m(\mathbf{t})^{x_m}) \times \mathcal{E}(\mathbf{t}) \mathcal{G}(\mathbf{x}) \mathcal{Q}(\mathbf{y}),$$

where

$\mathcal{E}$  is hyperexp. w.r.t.  $\{\delta_1, \dots, \delta_\ell\}$ ,

$\mathcal{G}$  hypergeom. w.r.t.  $\{\sigma_1, \dots, \sigma_m\}$ ,

$\mathcal{Q}$   $q$ -hypergeom. w.r.t.  $\{\tau_1, \dots, \tau_n\}$ .

# Standard representations

Let  $\prec$  be a monomial ordering in  $\mathbb{F}[\mathbf{t}, \mathbf{x}, \mathbf{y}]$ .

**Definition.** Given a sequence  $R$  of compatible rational functions

$$u_1, \dots, u_\ell, v_1, \dots, v_m, w_1, \dots, w_n$$

in  $\mathbb{F}(\mathbf{t}, \mathbf{x}, \mathbf{y})$ , a sequence  $S$  of rational functions:

$$f, \alpha_1, \dots, \alpha_m, \beta_1, \dots, \beta_\ell, \lambda_1, \dots, \lambda_m, \mu_1, \dots, \mu_n$$

is the **standard representation** of  $R$  if

- (i) the members of  $R$  and  $S$  satisfy the relations in the structure theorem;
- (ii)  $f$  has no nontrivial factors in  $\mathbb{F}[\mathbf{t}] \cup \mathbb{F}[\mathbf{x}] \cup \mathbb{F}[\mathbf{y}]$ ;
- (iii)  $f, \alpha_1, \dots, \alpha_m$  are monic w.r.t.  $\prec$ .

## Computing a representation

For brevity, let  $n = 0$ .

**Algorithm.** Given compatible rational functions  $u_1, \dots, u_\ell, v_1, \dots, v_m$ , compute  $f \in \mathbb{F}(\mathbf{t}, \mathbf{x})$ ,  $\alpha_1, \dots, \alpha_m, \beta_1, \dots, \beta_\ell \in \mathbb{F}(\mathbf{t})$ , and  $\lambda_1, \dots, \lambda_m \in \mathbb{F}(\mathbf{x}) \in \mathbb{F}(\mathbf{x})$  s.t.

$$u_i = \frac{\delta_i(f)}{f} + \sum_{j=1}^m x_j \frac{\delta_i(\alpha_j)}{\alpha_j} + \beta_i \quad \text{for } 1 \leq i \leq \ell,$$

and

$$v_j = \frac{\sigma_j(f)}{f} \alpha_j \lambda_j \quad \text{for } 1 \leq j \leq m.$$

(1) [**Determine  $\alpha_j$  and  $\lambda_j$ .**]

(1.1) For all  $j$ , compute  $f_j, a_j, b_j$  s.t.

$$v_j = f_j a_j b_j,$$

where  $f_j \in \mathbb{F}(\mathbf{t}, \mathbf{x})$  has no factors in  $\mathbb{F}(\mathbf{t}) \cup \mathbb{F}(\mathbf{x})$ ,  $a_j \in \mathbb{F}(\mathbf{t})$  is monic, and  $b_j \in \mathbb{F}(\mathbf{x})$ ;

(1.2) Compute  $g_j, r_j \in \mathbb{F}(\mathbf{t}, \mathbf{x}, \mathbf{y})$  s.t.

$$f_j = \frac{\sigma_j(g_j)}{g_j} r_j \quad \text{with } r_j \text{ being } \sigma_j\text{-reduced;}$$

(1.3) Set  $\alpha_j = a_j$  and  $\lambda_j = r_j b_j$ ;

(2) [**Determine  $f$ .**] Set  $f$  to be a nonzero solution of the system:

$$\left\{ \sigma_j(z) = \frac{\sigma_j(g_j)}{g_j} z \mid j = 1, \dots, m \right\}.$$

(3) [**Determine  $\beta_i$** ] For all  $i$ , set

$$\beta_i = u_i - \frac{\delta_i(f)}{f} - \sum_{j=1}^m x_j \frac{\delta_i(\alpha_j)}{\alpha_j};$$

(4) Return:  $f, \alpha_1, \dots, \alpha_m, \beta_1, \dots, \beta_\ell, \lambda_1, \dots, \lambda_m$ .

# Example

Consider an  $H$ -element  $h(t, x_1, x_2, x_3, x_4)$  with certificates

$$u = \frac{2x_1 + x_2 t - x_2}{t^2 - 1}$$

$$u = x_1 \frac{\delta(\alpha_1)}{\alpha_1} + x_2 \frac{\delta(\alpha_2)}{\alpha_2}$$

$$v_1 = \frac{(t-1)(x_2-x_1)(x_2+x_4-x_1)}{(x_3+x_1+1)(x_1+1)(t+1)}$$

$$v_1 = \alpha_1 \lambda_1$$

$$v_2 = \frac{(t+1)(x_2+x_4+1)(x_2+x_3+1)}{2(x_2-x_1+1)(x_2+x_4-x_1+1)}$$

$$\rightsquigarrow v_2 = \alpha_2 \cdot \lambda_2$$

$$v_3 = \frac{x_2+x_3+1}{x_3+x_1+1}$$

$$v_3 = \lambda_3$$

$$v_4 = \frac{x_2+x_4+1}{x_2+x_4-x_1+1}$$

$$v_4 = \lambda_4$$

where

$$\alpha_1 = \frac{t-1}{t+1}, \quad \alpha_2 = \frac{t+1}{2},$$

$$\lambda_1 = \frac{(x_2-x_1)(x_2+x_4-x_1)}{(x_3+x_1+1)(x_1+1)}, \quad \lambda_2 = \frac{(x_2+x_4+1)(x_2+x_3+1)}{(x_2-x_1+1)(x_2+x_4-x_1+1)}, \quad \lambda_3 = \frac{x_2+x_3+1}{x_3+x_1+1}, \quad \lambda_4 = \frac{x_2+x_4+1}{x_2+x_4-x_1+1}.$$

A multiplicative form of  $h$  is

$$h = \left( \frac{t-1}{t+1} \right)^{x_1} \left( \frac{t+1}{2} \right)^{x_2} \binom{x_2+x_3}{x_2-x_1} \binom{x_2+x_4}{x_1}.$$

## Is $h$ in $\mathbb{F}(\mathbf{t}, \mathbf{x}, \mathbf{y})$ ?

Let  $h$  be an  $H$ -element of the form:

$$h = f(\mathbf{t}, \mathbf{x}, \mathbf{y}) \alpha_1(\mathbf{t})^{x_1} \cdots \alpha_m(\mathbf{t})^{x_m} \mathcal{E}(\mathbf{t}) \mathcal{G}(\mathbf{x}) \mathcal{Q}(\mathbf{y}),$$

where the  $\alpha_j$ 's are monic

**Prop.**  $h \in \mathbb{F}(\mathbf{t}, \mathbf{x}, \mathbf{y}) \Leftrightarrow \alpha_1 = \cdots = \alpha_m = 1, \mathcal{E}(\mathbf{t}) \in \mathbb{F}(\mathbf{t}), \mathcal{G}(\mathbf{x}) \in \mathbb{F}(\mathbf{x})$  and  $\mathcal{Q}(\mathbf{y}) \in \mathbb{F}(\mathbf{y})$ .

**Facts:**

1.  $\mathcal{E}(\mathbf{t}) \in \mathbb{F}(\mathbf{t}) \Leftrightarrow \frac{\delta_i(\mathcal{E})}{\mathcal{E}} = \frac{\delta_i(g_i)}{g_i}$  for some  $g_i \in \mathbb{F}(\mathbf{t}), i = 1, \dots, \ell$ .
2.  $\mathcal{G}(\mathbf{x}) \in \mathbb{F}(\mathbf{x}) \Leftrightarrow \frac{\sigma_j(\mathcal{G})}{\mathcal{G}} = \frac{\sigma_j(g_j)}{g_j}$  for some  $g_j \in \mathbb{F}(\mathbf{x}), j = 1, \dots, m$ .
3.  $\mathcal{Q}(\mathbf{y}) \in \mathbb{F}(\mathbf{y}) \Leftrightarrow \frac{\tau_k(\mathcal{Q})}{\mathcal{Q}} = \frac{\tau_k(g_k)}{g_k}$  for some  $g_k \in \mathbb{F}(\mathbf{y}), k = 1, \dots, n$ .

## Is $h$ algebraic over $\mathbb{F}(\mathbf{t}, \mathbf{x}, \mathbf{y})$ ?

**Prop.**  $h$  is algebraic over  $\mathbb{F}(\mathbf{t}, \mathbf{x}, \mathbf{y})$  iff  $h^d \in \mathbb{F}(\mathbf{t}, \mathbf{x}, \mathbf{y})$  for some  $d \in \mathbb{Z}$  with  $d \neq 0$ .

**Corollary.** Let  $h$  be an  $H$ -element of the form:

$$h = f(\mathbf{t}, \mathbf{x}) \alpha_1(\mathbf{t})^{x_1} \cdots \alpha_m(\mathbf{t})^{x_m} \mathcal{E}(\mathbf{t}) \mathcal{G}(\mathbf{x}),$$

where the  $\alpha_j$ 's are monic. Let  $\mathbb{U}$  denote the set of roots of unity in  $\mathbb{F}$ . Then  $h$  is algebraic over  $\mathbb{F}(\mathbf{t}, \mathbf{x})$  iff

- (i)  $\alpha_1 = \cdots = \alpha_m = 1$ ,
- (ii)  $\frac{\delta_i(\mathcal{E})}{\mathcal{E}} = r_i \frac{\delta_i(g_i(\mathbf{t}))}{g_i(\mathbf{t})}$ , where  $r_i \in \mathbb{Q}$ ,  $g_i \in \mathbb{F}(\mathbf{t})$ ,  $i = 1, \dots, \ell$ ;
- (iii)  $\frac{\sigma_j(\mathcal{G})}{\mathcal{G}} = c_j \frac{\sigma_j(g_j(\mathbf{x}))}{g_j(\mathbf{x})}$ , where  $c_j \in \mathbb{U}$ ,  $g_j \in \mathbb{F}(\mathbf{x})$ ,  $j = 1, \dots, m$ ;

# Example

Let  $h$  satisfy  $\delta_1(h) = u_1h$ ,  $\delta_2(h) = u_2h$  and  $\sigma(h) = vh$ , where

$$u_1 = \frac{3t_1 + 1}{2t_1(t_1 + 1)}, \quad u_2 = \frac{x - 2t_2}{3t_2(x + t_2)}, \quad v = -\frac{x + t_2}{x + t_2 + 1}$$

- By the structure theorem,

$$f = \frac{1}{x + t_2}, \quad \alpha = 1, \quad \beta_1 = \frac{3t_1 + 1}{2t_1(t_1 + 1)}, \quad \beta_2 = \frac{1}{3t_2}, \quad \lambda = -1.$$

- Christopher's theorem

$$\beta_1 = \frac{\delta_1(t_1 + 1)}{t_1 + 1} + \frac{1}{2} \frac{\delta_1(t_1)}{t_1}, \quad \beta_2 = \frac{\delta_2(t_1 + 1)}{t_1 + 1} + \frac{1}{3} \frac{\delta_2(t_2)}{t_2}.$$

- Observe that  $\lambda = -1$ .

So  $h$  is algebraic over  $\mathbb{F}(t_1, t_2, x)$ .

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$h$  is of the form

$$h = c \frac{1}{x + t_2} \left( (t_1 + 1) t_1^{\frac{1}{2}} t_2^{\frac{1}{3}} \right) (-1)^x.$$

## Algebraic dependence (ongoing project)

**Problem.** Given  $H$ -elements  $h_1, \dots, h_s$  over  $\mathbb{F}(\mathbf{t}, \mathbf{x}, \mathbf{y})$ , determine whether they are algebraically dependent over  $\mathbb{F}(\mathbf{t}, \mathbf{x}, \mathbf{y})$ .

For  $i = 1, \dots, s$ , write

$$h_i = f_i \alpha_{i,1}^{x_1} \cdots \alpha_{i,m}^{x_m} \mathcal{E}_i(\mathbf{t}) \mathcal{G}_i(\mathbf{x}) \mathcal{Q}_i(\mathbf{y}),$$

where  $\alpha_{i,j} \in \mathbb{F}(\mathbf{t})$  is monic,  $\mathcal{E}_i(\mathbf{t})$  hyperexp.,  $\mathcal{G}_i(\mathbf{x})$  hypergeom., and  $\mathcal{Q}_i$   $q$ -hypergeom..

**Prop.**  $h_1, \dots, h_s$  are a.d. over  $\mathbb{F}(\mathbf{t}, \mathbf{x}, \mathbf{y})$  iff there exist  $e_1, \dots, e_s \in \mathbb{Z}$ , not all zero, s.t.

(i)

$$\alpha_{1,j}^{e_1} \cdots \alpha_{s,j}^{e_s} = 1 \quad \text{for } 1 \leq j \leq m;$$

(ii)

$$\mathcal{E}_1^{e_1} \cdots \mathcal{E}_s^{e_s} \in \mathbb{F}(\mathbf{t}), \quad \mathcal{G}_1^{e_1} \cdots \mathcal{G}_s^{e_s} \in \mathbb{F}(\mathbf{x}), \quad \mathcal{Q}_1^{e_1} \cdots \mathcal{Q}_s^{e_s} \in \mathbb{F}(\mathbf{y}).$$

## Symbolic power functions

**Problem.** Given  $\alpha_{i,j} \in \mathbb{F}(t)$  is monic, where  $i = 1, \dots, s$  and  $j = 1, \dots, m$ , compute the  $\mathbb{Z}$ -submodule  $M_{\alpha,j}$  of  $\mathbb{Z}^s$  s.t.

$$\alpha_{1,j}^{e_{1,j}} \cdots \alpha_{s,j}^{e_{s,j}} = 1$$

for all  $(e_{1,j}, \dots, e_{s,j}) \in \mathbb{Z}^s$ .

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### A method

1. Determine  $M_{\alpha,j}$  — computing polynomial gcds and solving linear systems over  $\mathbb{Z}$ .
2.  $M_{\alpha} = M_{\alpha,1} \cap \cdots \cap M_{\alpha,m}$ .

## Example

Let

$$f_1 = \frac{1}{(t^2 - 1)(t^3 - 1)} \quad \text{and} \quad f_2 = \frac{(t - 1)^2}{(t + 1)^3(t^2 + t + 1)}.$$

Decide if  $f_1^x$  and  $f_2^x$  are a.d. over  $\mathbb{F}(t, x)$ .

By gcd-computation

$$f_1 = \frac{1}{(t - 1)^2(t + 1)(t^2 + t + 1)} \quad \text{and} \quad f_2 = \frac{(t - 1)^2}{(t + 1)^3(t^2 + t + 1)}.$$

$$f_1^{e_1} f_2^{e_2} = 1 \text{ iff}$$

$$\begin{cases} 2e_1 - 2e_2 = 0 \\ e_1 + 3e_2 = 0 \\ e_1 + e_2 = 0 \end{cases}$$

$$\text{iff } e_1 = e_2 = 0.$$

**Answer.**  $f_1^x$  and  $f_2^x$  are a.i. over  $\mathbb{F}(t, x)$ .

## Hyperexponential case

Let  $\delta = d/dz$

**Lemma.** Let  $\mathcal{E}_1, \dots, \mathcal{E}_s$  be hyperexponential over  $\mathbb{F}(z)$ , and

$$\frac{\delta(\mathcal{E}_i)}{\mathcal{E}_i} = \beta_i, \quad i = 1, \dots, s.$$

Then

$$\exists e_1, \dots, e_s \in \mathbb{Z}, \quad \prod_{i=1}^s \mathcal{E}_i^{e_i} \in \mathbb{F}(z)$$

$\Updownarrow$

$\sum_{i=1}^s e_i \beta_i$  is a logarithmic derivative of some element in  $\mathbb{F}(z)$ .

## Example

Let

$$\mathcal{E}_1 = \exp \left( \int \left( \frac{1}{z} - \frac{1}{z-1} + 4z \right) dz \right) \quad \text{and} \quad \mathcal{E}_2 = \exp \left( \int \left( \frac{1}{z} - \frac{2z}{z^2+1} - 2z \right) dt \right).$$

Decide if  $\mathcal{E}_1$  and  $\mathcal{E}_2$  are a.d. over  $\mathbb{F}(z)$ .

By partial fraction decomposition,

$$\frac{\delta(\mathcal{E}_1)}{\mathcal{E}_1} = \left( \frac{1}{z} - \frac{1}{z-1} \right) + 4z \quad \text{and} \quad \frac{\delta(\mathcal{E}_2)}{\mathcal{E}_2} = \left( \frac{1}{z} - \frac{2z}{z^2+1} \right) - 2z.$$

$\mathcal{E}_1^{e_1} \mathcal{E}_2^{e_2} \in \mathbb{F}(z)$  iff

$$e_1 \frac{\delta(\mathcal{E}_1)}{\mathcal{E}_1} + e_2 \frac{\delta(\mathcal{E}_2)}{\mathcal{E}_2} \text{ is a logarithmic derivative}$$

iff  $4e_1 - 2e_2 = 0$ .

**Answer.**  $\mathcal{E}_1$  and  $\mathcal{E}_2$  are a.d. over  $\mathbb{F}(z)$ .

# Hypergeometric case

Let  $\sigma(f(z)) = f(z + 1)$  for all  $f \in \mathbb{F}(z)$ .

**Lemma.** Let  $\mathcal{G}_1, \dots, \mathcal{G}_s$  be hypergeometric over  $\mathbb{F}(z)$ , and

$$\frac{\sigma(\mathcal{G}_i)}{\mathcal{G}_i} = \lambda_i, \quad i = 1, \dots, s.$$

Write  $\lambda_i = c_i \tilde{\lambda}_i$ , where  $c_i$  is in  $\mathbb{F}$  and  $\tilde{\lambda}_i$  is monic. Then

$$\exists e_1, \dots, e_s \in \mathbb{Z}, \quad \prod_{i=1}^s \mathcal{G}_i^{e_i} \in \mathbb{F}(z)$$

$\Leftrightarrow$

$$\prod_{i=1}^s c_i^{e_i} = 1 \quad \text{and} \quad \sum_{i=1}^s e_i \frac{\delta(\lambda_i)}{\lambda_i} = \sigma \left( \frac{\delta(g)}{g} \right) - \frac{\delta(g)}{g}$$

for some  $g \in \mathbb{F}(z)$ .

**Reference.** M. Singer. *Deciding if solutions of  $\sigma(y_1) = f_1 y_1, \dots, \sigma y_n = f_n y_n$  are algebraically dependent over  $\mathbb{C}(x)$* . A note on discussions at the 2nd NCSU-China Symb. Comput. Collaboration Workshop, Hangzhou, 2007.

# Subproblems

1. Given  $c_1, \dots, c_s$  in a finite algebraic extension over  $\mathbb{Q}$ , find all vectors  $(e_1, \dots, e_s) \in \mathbb{Z}^s$  s.t.

$$\prod_{i=1}^s c_i^{e_i} = 1.$$

**Reference.** M. Kauers. *Algorithms for Nonlinear Higher Order Difference Equations*. PhD thesis, RISC-Linz, Linz, Austria, 2005.

2. Given  $\lambda_1, \dots, \lambda_s \in \mathbb{F}(z)$ , find all vectors  $(e_1, \dots, e_s) \in \mathbb{Z}^s$  s.t.

$$\sum_{i=1}^s e_i \frac{\delta(\lambda_i)}{\lambda_i} = \sigma \left( \frac{\delta(g)}{g} \right) - \frac{\delta(g)}{g}$$

for some  $g \in \mathbb{F}(z)$ .

**Reference.** Singer's note

## $q$ -Hypergeometric case

Let  $\tau(f(z)) = f(qz)$  for all  $f \in \mathbb{F}(z)$ , where  $q$  is not a root of unity.

Let  $\mathcal{Q}_1, \dots, \mathcal{Q}_s$  be  $q$ -hypergeometric over  $\mathbb{F}(z)$ , and

$$\frac{\tau(\mathcal{Q}_i)}{\mathcal{Q}_i} = \mu_i, \quad i = 1, \dots, s.$$

Write  $\mu_i = c_i \tilde{\mu}_i$ , where  $c_i$  is in  $\mathbb{F}$  and  $\tilde{\mu}_i$  is monic.

**Lemma.** Assume that there exist  $e_1, \dots, e_s \in \mathbb{Z}$  s.t.  $\prod_{i=1}^s \mathcal{Q}_i^{e_i} \in \mathbb{F}(z)$ . Then

(i) there exists  $g \in \mathbb{F}(z)$  s.t.

$$\sum_{i=1}^s e_i \frac{\delta(\mu_i)}{\mu_i} = \tau\left(\frac{\delta(g)}{g}\right) - \frac{\delta(g)}{g};$$

(ii)

$$\prod_{i=1}^s c_i^{e_i} = q^d,$$

where  $d = \deg \text{num}(g) - \deg \text{den}(g)$ .

## Summary

$$h = \text{a rational function} \times \alpha_1(\mathbf{t})^{x_1} \cdots \alpha_m(\mathbf{t})^{x_m} \mathcal{E}(\mathbf{t}) \mathcal{G}(\mathbf{x}) \mathcal{Q}(\mathbf{y}).$$

## Future work

1. Implementation of algorithms for decomposing  $H$ -elements;
2. An algorithm for determining algebraic dependence of  $H$ -elements;
3. Criteria for the existence of telescopers.
4. ...