

# Patterns in Trees

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## Outline

- Present assumptions and basic techniques.
- Define patterns.
- Find representations as combinatorial objects.
- Convert to corresponding generating functions.
- Compute asymptotics.

## 1. Basics

Labelled universe

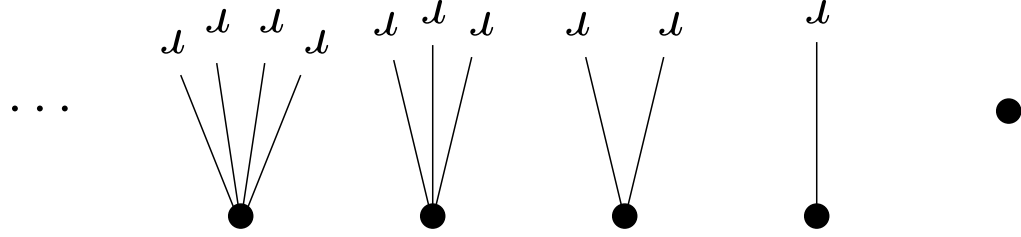
Counting of combinatorial structures by exponential generating functions

$$\sum d \frac{u^d}{d!} = (z)$$

First model: All trees of size  $n$  have equal probability.

# Rooted Trees

Construction (rooted):



Tree function:

$$z^d = z \left( 1 + d(z) + \frac{d^2(z)}{2} + \frac{d^3(z)}{3} + \dots \right)$$

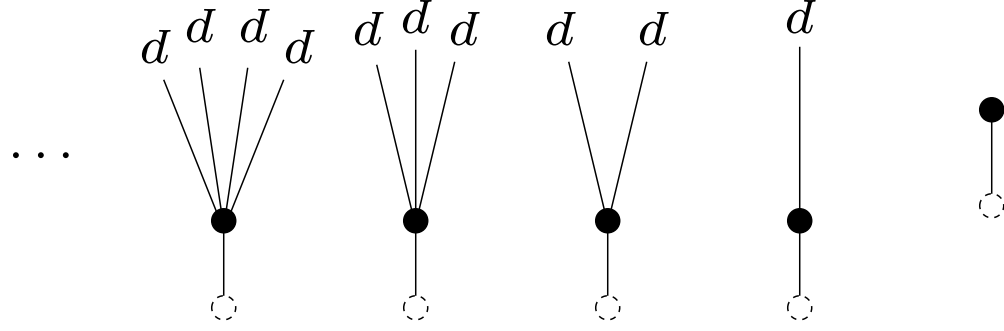
Explicit number of rooted trees of size  $n$  via e.g. Lagrange inversion formula

$$d^n = n^{-1}$$

## Planted Rooted Trees

Sometimes helpful: planted rooted trees.

Node degree does not change during construction.



Results in same function  $p(z)$ .

First order asymptotics

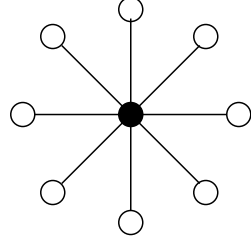
$$d(z) = 1 - \sqrt{2\sqrt{1-az} + O(z)}$$

## Bivariate Generating Functions

Mark special properties with a second variable  $u$

E.g.: Nodes with particular degree  $k + 1$

$$d = d e^{uz} - \frac{uz}{d} + n \frac{uz}{d} + (n-1) \frac{uz}{d}$$



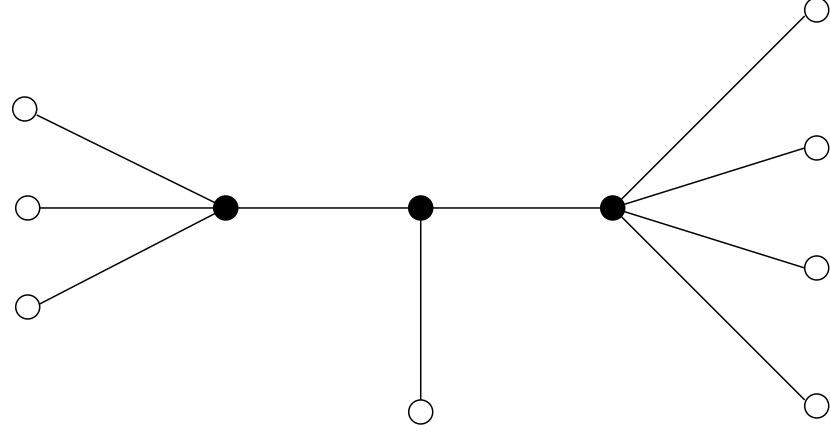
Possible to find limit distribution and compute parameters. In this case [Drmotá and Gittenberger 1997]: asymptotically Gaussian with mean  $uzn$  and variance  $\frac{uz}{d}$ , where  $m_k = \frac{uz}{d}$  und  $\sigma_k^2 = -\frac{uz(1-uz)^2}{1+(k-2)uz} + \frac{uz(1-uz)}{1}$

## 2. Patterns

What is a pattern?

In our case, connected sub-tree  $M$ .

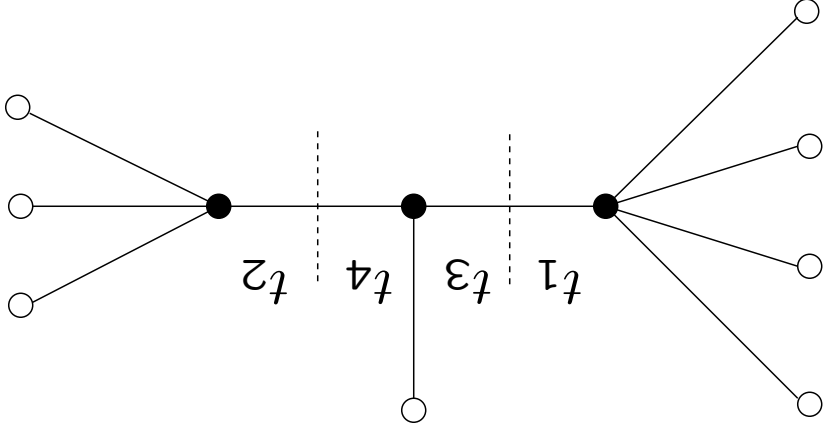
Easy example:



## How to Mark a Pattern?

More difficult than with single nodes.

Split in parts:



Aim: Describe patterns by generating functions for each part  
→ system of functional equations.



$$\begin{aligned}
 & a_1(x, n) = \sum_{j=1}^L a_j(x, n) \\
 & a_2(x, n) = P_1(a_1(x, n), \dots, a_L(x, n)) \cdot x \\
 & \vdots \\
 & a_L(x, n) = P_{L-1}(a_1(x, n), \dots, a_{L-1}(x, n)) \cdot x - \sum_{j=1}^{L-1} a_j(x, n)
 \end{aligned}
 \tag{1}$$

and polynomials  $P_j(y_1, \dots, y_L, n)$  ( $1 \leq j \leq L-1$ ) with non-negative coefficients such that

$$d(x, n) = \sum_{j=1}^L a_j(x, n)$$

Let  $\mathcal{M}$  be a pattern. Then there exist  $L$  auxiliary functions  $a_j(x, n)$  ( $1 \leq j \leq L$ ) with

**Proposition (Planted Rooted Trees)**

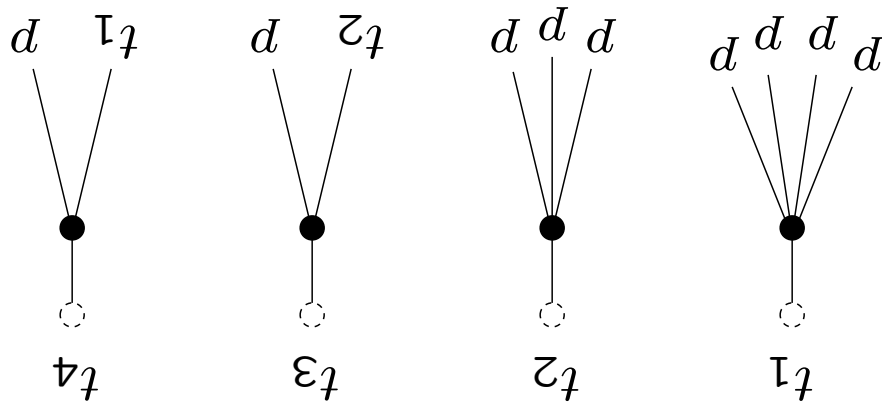
## Notation to Describe Patterns

◦ is a (planted) root node,  $\times$  the Cartesian product,  $\cap$  the intersection,  $\cup$  the union, and  $\dot{\cup}$  the disjoint union.

$\times$  binds stronger than either of  $\cap$ ,  $\cup$ , and  $\dot{\cup}$ .

Note: no immediate one-to-one correspondence to the generating functions (relative probabilities,  $n$ ).

$$\begin{aligned}
 d \times \mathbb{1}t &= t_4 \\
 d \times \mathbb{2}t &= t_3 \\
 d \times d \times d \times \{\circ\} &= t_2 \\
 d \times d \times d \times d \times \{\circ\} &= t_1
 \end{aligned}$$



### 3. Convert to Functions

- Standardise: no duplicate descriptions of the same structure
- Find coefficients
- Sprinkle with  $u$

## Standard Form

Each tree  $a_j$  represented as disjoint union of trees of the kind

$$(2) \quad \{ \circ \} \times a_{l_1} \times \dots \times a_{l_p},$$

( $d =$  degree of the root of  $a_j$ ).

## Standardising the Functions

Build intersections (symbolically):

$$\begin{aligned}
 x_1 &= \{o\} \times x_{1,1} \times \dots \times x_{1,l} \\
 x_2 &= \{o\} \times x_{2,1} \times \dots \times x_{2,l} \\
 &\quad \bigcap_{\substack{\{l', \dots, 1\} = \{l_1, \dots, l_n\} \\ \{l', \dots, 1\} = \{m_1, \dots, m_l\}}} \\
 x_1 \cup x_2 &= (x_{1,1} \cup x_{2,1}) \times \dots \times (x_{1,l} \cup x_{2,l})
 \end{aligned}$$

## Intersection Example

$$\begin{aligned}t_1 &= \{o\} \times d \times d \times d \times d, \\t_2 &= \{o\} \times d \times d \times d, \\t_3 &= \{o\} \times d \times d, \\t_4 &= \{o\} \times d.\end{aligned}$$

Only one non-empty intersection:

$$\begin{aligned}t_3 \cap t_4 &= \{o\} \times (t_1 \cap t_2) \times d \cup \{o\} \times t_1 \times t_2 = \\&= \{o\} \times t_1 \times t_2.\end{aligned}$$





## Coefficients

$A^{j,l_1,\dots,l_T,k}$  := number of possible configurations of type (2)  
( $l_i$  sub-trees of type  $a_i$ ,  $k$  new occurrences of  $\mathcal{M}$ )

Coefficients are computed by simple combinatorics ( $k$  is implicitly given by  $l_i$ ).

## Resulting Functions

$$P_j^j(y_1, \dots, y_L, u) = \sum_{l_1, \dots, l_T \geq 0} \frac{A_{j, l_1, \dots, l_T, k}^{j, l_1, \dots, l_T, i}}{y_1^{l_1} \dots y_T^{l_T} u^k}, \quad 1 \leq j \leq L - 1$$

$$P_L^L(y_1, \dots, y_L) = e^{y_1 + \dots + y_L} - \sum_{j=1}^{L-1} P_j^j(y_1, \dots, y_L, 1).$$

$$a_j^j(x, u) = x \cdot P_j^j(a_1(x, u), \dots, a_L(x, u)).$$

→ proposed structure of the system of functional equations (1).

## How to Find $k = k(l_1, \dots, l_T)$

New patterns occur when all necessary sub-trees are attached to a node of proper degree.

In example, three cases:

1. Node of degree three with a  $t_1$  and  $t_2$ .
2. Node of degree four with a  $t_4$  attached. Each  $t_4$  produces another pattern.
3. Node of degree five with a  $t_3$  attached. Each  $t_3$  produces another pattern.

## Coefficients in Example

$$P_1 = y_5(y_1 + y_2 + y_3 + y_6) + \frac{1}{2}iy_5,$$

$$P_2 = y_4(y_1 + y_2 + y_3 + y_6) + \frac{1}{2}iy_4,$$

$$P_3 = uy_4y_5,$$

$$P_4 = \frac{4i}{(uy_1 + y_2 + uy_3 + y_4 + y_5 + y_6)^4},$$

$$P_5 = \frac{3i}{(y_1 + uy_2 + uy_3 + y_4 + y_5 + y_6)^3}$$

## Coefficients in Example (cont'd)

$$a_1(x, n) = a_1 + a_2 + a_3 + a_6 + x \frac{1}{2} a_5,$$

$$a_2(x, n) = a_2 + a_3 + a_4 + a_6 + x \frac{1}{2} a_4,$$

$$a_3(x, n) = a_3 + x a_4 a_5,$$

$$a_4(x, n) = \frac{4!}{(a_1 + a_2 + a_3 + a_4 + a_5 + a_6)^4} x,$$

$$a_5(x, n) = \frac{3!}{(a_1 + a_2 + a_3 + a_4 + a_5 + a_6)^3} x,$$

$$a_6(x, n) = x + \sum_{i=1}^6 a_i x + \frac{2!}{(a_1 + a_2 + a_3 + a_6)^2} x +$$

$$+ \sum_{i=1}^6 \frac{1}{a_i} \sum_{j=1}^i a_j \cdot$$

## **Strong Connectivity**

$a_6$  depends on itself and all others (last term).

Each  $a_i$  depends on  $a_6$  either directly, or through a chain to a leaf (see pattern).

## Proposition (Rooted Trees)

There exists an analytic function

$$G(x, n, a_1, \dots, a_T)$$

with non-negative Taylor coefficients such that

$$r(x, n) = G(x, n, a_1, \dots, a_T(x, n)).$$

where the  $a_i$  were defined earlier.

## Counting Patterns in Rooted Trees

For the example:

$$r_0 := x e^d - \frac{5i}{x d^5} - \frac{4i}{x d^4} - \frac{3i}{x d^3}.$$

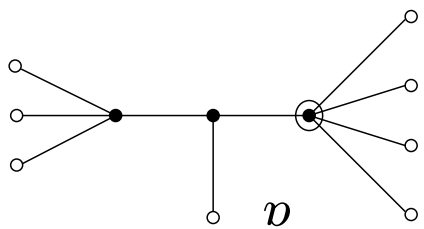
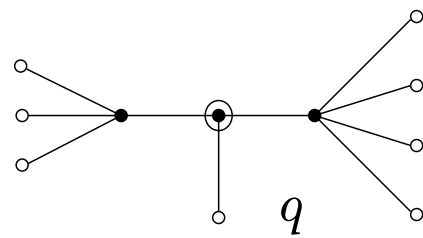
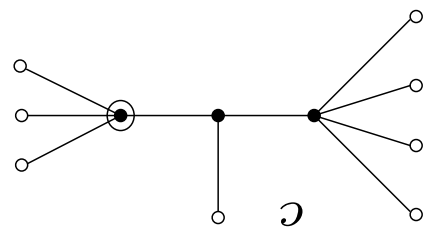
("uninteresting" sub-trees)

Marks in sub-trees already counted correctly, only have to dis-tribute marks for newly appearing patterns.



## Newly Appearing Patterns

$$r = r(x, n) = r_0 + x \sum_{\substack{e_i \geq 0 \\ e_i = 5}} x + n^{e_1 + e_3} \prod_{e_j} a_{j=1}^{e_j} \prod_{e_i} i + x + n^{e_4 e_5} \prod_{e_j} a_{j=1}^{e_j} \prod_{e_i} i + x \sum_{\substack{e_i \geq 0 \\ e_i = 3}} x + n^{e_2 + e_3} \prod_{e_j} a_{j=1}^{e_j} \prod_{e_i} i$$



## Unrooted Trees

$$t_{n,m} = r_{n,m}/n$$

## 4. Asymptotics

Use Drmota's Theorem on systems of functional equations. \*

Paraphrased: Under certain conditions for the system of equations, the coefficients asymptotically follow a Gaussian distribution with mean and variance asymptotically proportional to  $n$ .

Often the difficult part:

Finding the singularity

\*M. Drmota, *Systems of Functional Equations. Random Structures and Algorithms* 10, 103-124, 1997.

## For Our Case

Singularity known:

$$\frac{1}{e}$$

To find: left eigenvector (for the eigenvalue one) of the derivative of the functional matrix with respect to the functions

Application of theorem then gives expectation for examples as

$$\frac{384e - 19}{12e(576e^3 + 24e^2 - 25)} = 0.0026803\dots$$

## Possible Future Extensions

- Different kind of trees or tree distributions
- Logical terms using  $\forall$ ,  $\exists$  and  $\neg$  describing more general patterns.
- Does a 0-1-law hold?