

Singular Combinatorics

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Combinatorics: discrete structures by finitary rules
 \leadsto Enumerative & Quantitative aspects

- Counting and asymptotics

$$n! \sim n^n e^{-n} \sqrt{2\pi n}$$

- Asymptotic laws (e.g., Monkey and typewriter!)

$$\Omega_n \xrightarrow{\mathcal{D}} \frac{1}{\sqrt{2\pi}} \int_{-\infty}^x e^{-t^2/2} dt.$$

Approaches

- Probabilistic, stochastic
- Analytic: Generating Functions

1. Introduction

“Symbolic” Methods

Rota-Stanley; Foata-Schutzenberger; Joyal and UQAM group;
Jackson-Goulden, &c; ca 1980[±].

Basic combinatorial constructions admit of direct
translations as operators over generating functions (GF's).

\mathcal{C} : class of comb. structures;

C_n : # objects of size n

$$\begin{array}{ll} \text{(counting)} & \left\{ \begin{array}{l} C(z) := \sum C_n z^n \\ \hat{C}(z) := \sum C_n \frac{z^n}{n!} \end{array} \right. \\ \text{(params)} & \left\{ \begin{array}{l} C(z, u) := \sum C_{n,k} z^n u^k \\ \hat{C}(z, u) := \sum C_{n,k} u^k \frac{z^n}{n!} \end{array} \right. \end{array}$$

↓↓↓

Ordinary GF's for unlabelled structures

Exponential GF's for labelled structures.

“Dictionaries”

= **Constructions** viewed as **Operators** over GF's.

Constr.	Operations	
Union	$+$	$+$
Product	\times	\times
Sequence	$(1 - f)^{-1}$	$(1 - f)^{-1}$
MultiSet	Pólya Exp.	e^f
Cycle	Pólya Log.	$\log(1 - f)^{-1}$
	(unlab.)	(lab.)

Books: Goulden-Jackson, Bergeron-LL, Stanley, F-Sedgewick

\Rightarrow Formal power series w/ rich set of functional eqns!!

\Rightarrow How to extract coeff., especially, asymptotically??

$$\text{Exp}(f) := \exp \left(f(z) + \frac{1}{2}f(z^2) + \cdots \right); \quad \text{Log}(f) := \log \frac{1}{1 - f(z)} + \cdots$$

“Complex–analytic Structures”

Interpret:

♡ Counting GF as analytic transformation of \mathbb{C} ;

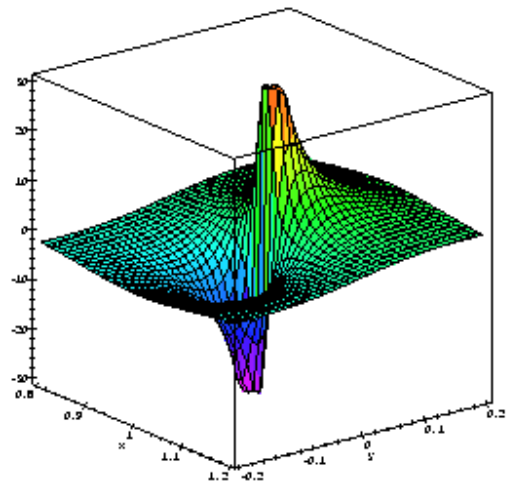
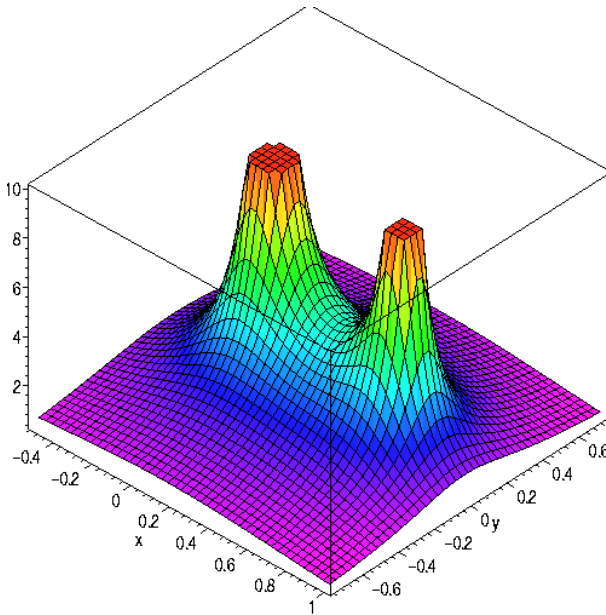
♡ Comb. Construction as analytic functional.

Singularities are crucial to asymptotic prop's!

(cf. analytic number theory, complex analysis, etc)

Asymptotic counting via Singularity Analysis (S.A.)

Asymptotic laws via Perturbation + S.A.



Refs: F–Odlyzko, SIAM A&DM, 1990; Odlyzko’s 1995 survey in *Handbook of Combinatorics*

+ Banderier, Fill, J. Gao, Gonnet, Gourdon, Kapur, G. Labelle, Laforest, T. Lafforgue, Noy, Odlyzko, Panario, Pouyanne, Prodinger, Puech, Richmond, Robson, Salvy, Schaeffer, Sipala, Soria, Steyaert, Szpankowski, B. Vallée.

♠ Location of singularity at $z = \rho$:

$$\text{coeff. } [z^n]f(z) = \rho^{-n} \cdot \text{coeff. } [z^n]f(\rho z),$$

where latter is singular at $z = 1$.

♠ Nature of singularity at $z = 1$:

$$\frac{1}{(1-z)^2} \longrightarrow n+1 \sim n$$

$$\frac{1}{1-z} \log \frac{1}{1-z} \longrightarrow H_n \equiv \frac{1}{1} + \dots + \frac{1}{n} \sim \log n$$

$$\frac{1}{1-z} \longrightarrow 1 \sim 1$$

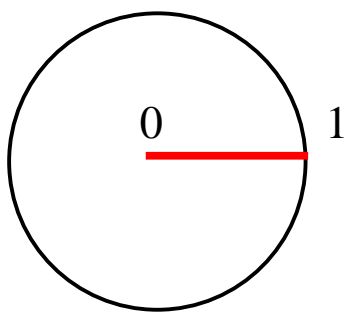
$$\frac{1}{\sqrt{1-z}} \longrightarrow \frac{1}{2^{2n}} \binom{2n}{n} \sim \frac{1}{\sqrt{\pi n}}$$

{	Location	of sing's :	Exponential factor	ρ^{-n}
	Nature	of sing's :	"Polynomial" factor	$\vartheta(n)$

Generating Function \leadsto Coefficients

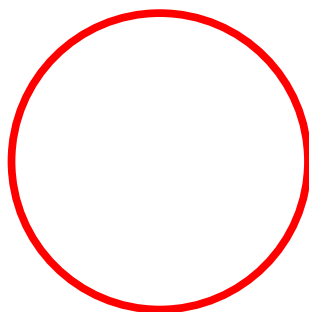
Solving a “Tauberian” problem

Real-Tauberian



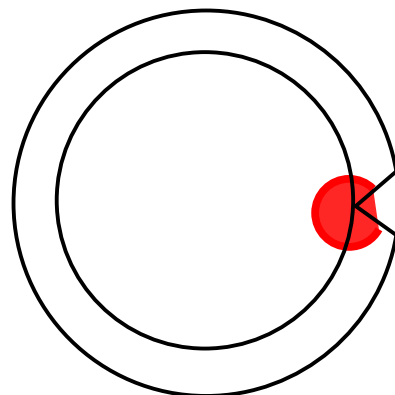
(large \implies large)

Darboux-Pólya



(smooth \implies small)

Singularity An.



(Full mappings)

Combinatorial constructions \leadsto Analytic Functionals

\implies Analytic continuation prevails for comb. GF's

2. Basic Singularity Analysis

Theorem 1. *Basic scale translates:*

$$\sigma_{\alpha,\beta}(z) := (1-z)^{-\alpha} \left(\frac{1}{z} \log \frac{1}{1-z} \right)^\beta$$

$$\implies [z^n] \sigma_{\alpha,\beta} \underset{n \rightarrow \infty}{\sim} \frac{n^{\alpha-1}}{\Gamma(\alpha)} (\log n)^\beta.$$

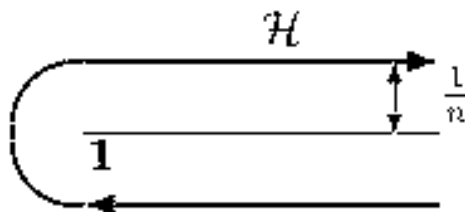
PROOF. Cauchy's coefficient integral, $f(z) = (1-z)^{-\alpha}$

$$[z^n] f(z) = \frac{1}{2i\pi} \int_{\gamma} f(z) \frac{dz}{z^{n+1}}$$

$$\Downarrow \quad \left(z = 1 + \frac{t}{n} \right) \quad \Downarrow$$

$$\frac{1}{2i\pi} \int_{\mathcal{H}} \left(-\frac{t}{n} \right)^{-\alpha} e^{-t} \frac{dt}{n}$$

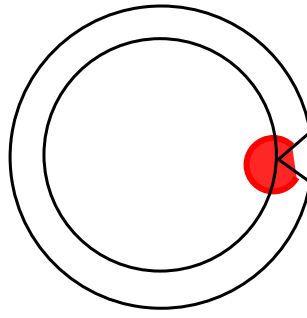
$$n^{\alpha-1} \times \frac{1}{\Gamma(\alpha)}.$$



♡ Slowly varying \implies slowly varying

♡ Log-log \implies Log-Log

♡ Full asymptotic expansions

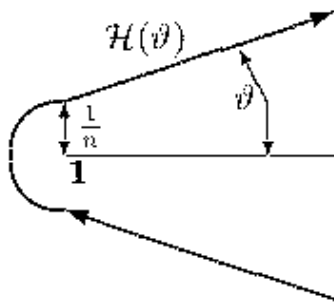


Theorem 2. *\mathcal{O} -transfers:*

Under continuation in a Δ -domain,

$$f(z) = O(\sigma_{\alpha,\beta}(z)) \implies [z^n]f(z) = O([z^n]\sigma_{\alpha,\beta}(z)).$$

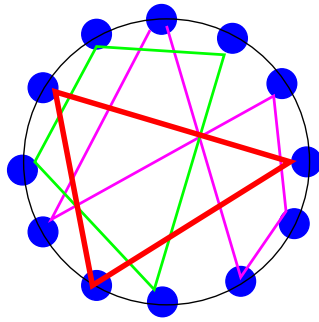
PROOF:



$$\text{Usage: } \left\{ \begin{array}{l} f(z) = \lambda\sigma(z) + \mu\tau(z) + \dots + O(\omega(z)) \\ \implies \\ f_n = \lambda\sigma_n + \mu\tau_n + \dots + O(\omega_n). \end{array} \right.$$

Similarly: *\mathcal{o} -transfer.*

- Dominant singularity at ρ gives factor ρ^{-n} .
- Finitely many singularities work fine
- Some cases with ∞ of sing's.



EXAMPLE 1. *2-regular graphs* [Comtet]

$$\mathcal{G} = \mathfrak{M} \left(\frac{1}{2} \mathfrak{C}_{\geq 3}(\mathcal{Z}) \right)$$

$$\hat{G}(z) = \exp \left(\frac{1}{2} \log \frac{1}{1-z} - \frac{z}{2} - \frac{z^2}{4} \right)$$

$$\hat{G}(z) \underset{z \rightarrow 1}{\sim} \frac{e^{-3/4}}{\sqrt{1-z}}$$

$$\frac{G_n}{n!} \underset{n \rightarrow \infty}{\sim} \frac{e^{-3/4}}{\sqrt{\pi n}}.$$

(Originally by Darboux-Pólya.)

□

EXAMPLE 2. *Richness index of trees* [FSS,90] = Number of different terminal subtrees. Catalan case:

$$K(z) = \frac{1}{2z} \sum_{k \geq 0} \frac{1}{k+1} \binom{2k}{k} \left(\sqrt{1-4z-4z^{k+1}} - \sqrt{1-4z} \right)$$

$$K(z) \underset{z \rightarrow 1/4}{\approx} \frac{1}{\sqrt{X \log X}}, \quad X := 1-4z$$

$$\text{Mean index} \underset{n \rightarrow \infty}{\sim} C \frac{n}{\sqrt{\log n}}, \quad C \equiv \sqrt{\frac{8 \log 2}{\pi}}.$$

Related to compact tree representations as DAGs.

□

3. Closure Properties

Function of S.A.-type = amenable to singularity analysis

- is continuable in a Δ -domain,
- admits singular expansion in scale $\{\sigma_{\alpha,\beta}\}$.

Theorem 3. *Generalized polylogarithms*

$$\text{Li}_{\alpha,k} := \sum (\log n)^k n^{-\alpha} z^n$$

are of S.A.-type.

PROOF. Cauchy-Lindelöf representations

$$\sum \varphi(n)(-z)^n = -\frac{1}{2i\pi} \int_{1/2-i\infty}^{1/2+i\infty} \varphi(s) z^s \frac{\pi}{\sin \pi s} ds.$$

+ Mellin transform techniques (Ford, Wong, F.).

EXAMPLE 3. *Entropy of Bernoulli distribution*

$$H_n := - \sum_k \pi_{n,k} \log \pi_{n,k}, \quad \pi_{n,k} \equiv \binom{n}{k} p^k (1-p)^{n-k}$$

$$\text{involves } \sum \log(k!) z^k = (1-z)^{-1} \text{Li}_{0,1}(z)$$

$$\frac{1}{2} \log n + \frac{1}{2} + \log \sqrt{2\pi p(1-p)} + \dots$$

(Redundancy, coding, information th.; Jacquet-Szpankowski.) □

- Elements like $\log n, \sqrt{n}$ in combinatorial sums

Theorem 4. *Functions of S.A.-type are closed under integration and differentiation.*

PROOF. Adapt from Olver, Henrici, etc.

Theorem 5. *Functions of S.A.-type are closed under Hadamard product*

$$f(z) \odot g(z) := \sum_n (f_n g_n) z^n.$$

PROOF. Start from Hadamard's formula

$$f(z) \odot g(z) = \frac{1}{2i\pi} \int_{\gamma} f(t) g\left(\frac{w}{t}\right) \frac{dt}{t}.$$

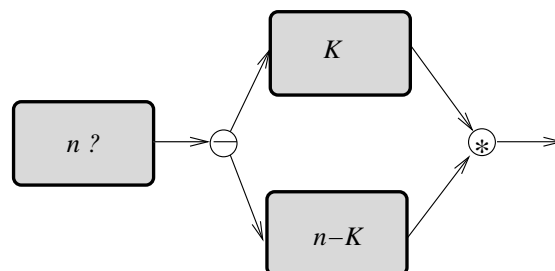
+ adapt Hankel contours [H., Jungen, R. Wilson; Fill-F-Kapur]

EXAMPLE 4. *Divide-and-conquer recurrences*

$$\begin{aligned} f_n &= t_n + \sum \pi_{n,k} (f_k + f_{n-k}) \\ \text{Sing}(f(z)) &= \Phi(\text{Sing}(t(z))) \\ \text{Asympt}[f_n] &= \Psi(\text{Sing}(t)). \end{aligned}$$

E.g., Catalan statistics: need $\sum \binom{2n}{n} \log n \cdot z^n$.

Useful in computer science applications [FFK, 2002⁺]. □



4. Functional Equations

- Rational functions. Combinatorics: linear system $\mathbb{Q}_{\geq 0}[z]$ implies **polar singularities** (X^{-k}):

$$[z^n]f(z) \approx \sum \omega^n n^k, \quad \omega \in \overline{\mathbb{Q}}, \quad k \in \mathbb{Z}_{\geq 0}.$$

+ irreducibility: Perron-Frobenius \implies **simple dom. pole**.

- Word problems from regular language models;
- Transfer matrices [Bender-Richmond]: easy dimer coverings in $k \times n$ strip, knight tours, etc.

- Algebraic functions, by Puiseux expansions ($X^{p/q}$):

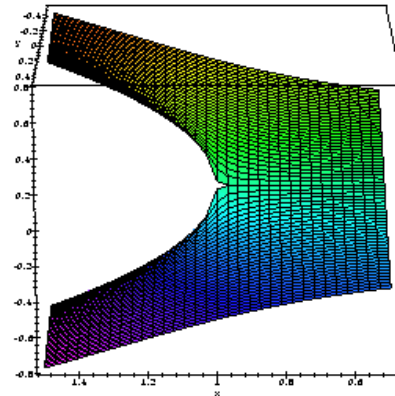
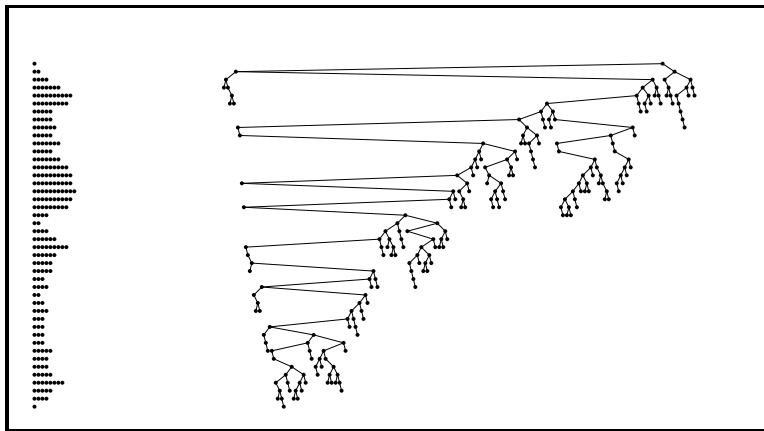
$$[z^n]f(z) \approx \sum \sum \omega^n n^{p/q}, \quad \omega \in \overline{\mathbb{Q}}, \quad p/q \in \mathbb{Q},$$

results from S.A. (or Darboux!)

Asymptotics of coeff. is decidable [Chabaud-F-Salvy].

- Word problems and context-free models;
- Trees; geom. configurations (non-crossing graphs, polygonal triang.); planar maps [Tutte].

$$(1 - \sqrt{1 - 4z}) / (2z)$$



Square-root singularity is “universal” for many recursive classes (irreducibility). Owing to controlled “failure” of Implicit Function Theorem where **quadratic** dependency replaces linear (=analytic) dependency. Entails coeff. asymptotic $\approx \omega^n n^{-3/2}$ with critical exponent $-3/2$ that **is universal**.

E.g., 2–3 trees (Meir-Moon):

$$f = z\phi(f), \quad \phi(u) = 1 + u^2 + u^3.$$

Pólya’s combinatorial chemistry programme:

$$f(z) = z \operatorname{Exp}(f(z)) \equiv ze^{f(z) + \frac{1}{2}f(z^2) + \frac{1}{3}f(z^3) + \dots}$$

Starting with Pólya 1937; Otter 1949; Harary-Robinson et al. 1970’s; Meir-Moon 1978; Drmota-Lalley-Woods thm.1990⁺

- “Holonomic” functions. Defined as solutions of linear ODE’s with coeffs in $\mathbb{C}(z)$ [Zeilberger] $\equiv \mathcal{D}$ -finite.

$$\mathcal{L}[f(z)] = 0, \quad \mathcal{L} \in \mathbb{C}(z)[\partial_z].$$

- Stanley, Zeilberger, Gessel: great importance for combinatorial enumeration: Young tableaux and permutation statistics; regular graphs, constrained matrices, etc.

Fuchsian case (or “regular” singularity) $(X^\beta \log^k X)$:

$$[z^n]f(z) \approx \sum \omega^n n^\beta (\log n)^k, \quad \omega, \beta \in \overline{\mathbb{Q}}, \quad k \in \mathbb{Z}_{\geq 0}.$$

S.A. applies automatically to classical classification.

Asymptotics of coeff is decidable

- general case: modulo oracle for connection problem;
- strictly positive case: unconditionally.

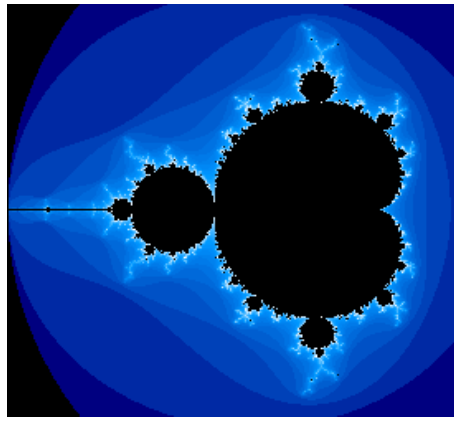
EXAMPLE 5. *Quadrees—Partial Match* [FGPR’92]

Divide-and-conquer recurrence with coeff. in $\mathbb{Q}(n)$

Fuchsian equation of order d (dimension) for GF

$$Q_n^{(d=2)} \approx n^{(\sqrt{17}-3)/2}.$$

E.g., $d = 2$: Hypergeom ${}_2F_1$ with algebraic arguments. □



- Functional Equations and Substitution.

- Early example of *balanced 2–3 trees* by Odlyzko, 1979.

$$T(z) = z + T(\tau(z)), \quad \tau(z) := z^2 + z^3.$$

Infinitely many exponents with common real part implies periodicities: $T_n \sim \frac{\phi^n}{n} \Omega(\log n)$.

- Singular iteration for *height of trees* (binary and other simple varieties; F-Gao-Odlyzko-Richmond; cf Rényi-Szekeres):

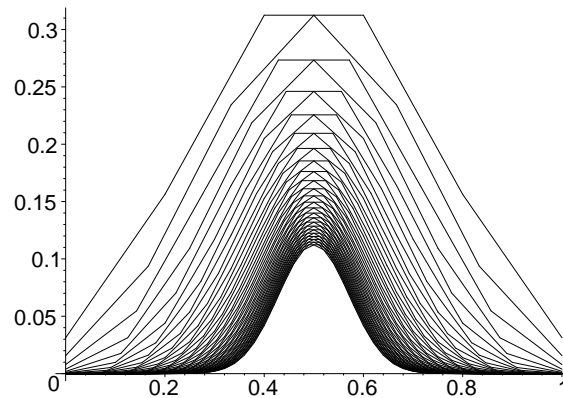
$$y_h = z + y_{h-1}^2, \quad y_0 = z.$$

— Moments and convergence in law; Local limit law of *ϑ-type*.
(Applies to branching processes conditioned on total progeny.)

- *Digital search trees* via *q-hypergeometrics*: singularities accumulate geometrically \leadsto *periodicities* [F-Richmond]:

$$\partial_z^k f(z) = t(z) + 2e^{z/2} f\left(\frac{z}{2}\right).$$

- *Order of binary trees* (Horton-Strahler, Register function; F-Prodinger) via *Mellin* tr. of GF and & singularities.



5. Limit Laws

Classical probability theory: sums of independent Random Variables \leadsto powers of fixed function (probability GF, Fourier tr.) \leadsto **Normal Law**.

For problems expressed by **Bivariate GF** (BGF): field founded by E. Bender et al. + developments by F, Soria, Hwang, . . .

Idea: BGF $F(z, u) = \sum f_n(u)z^n$, where $f_n(u)$ describes parameter on objects of size n . If (for u near 1)

$$f_n(u) \approx \omega(u)^{\kappa_n}, \quad \kappa_n \rightarrow \infty,$$

then speak of **Quasi-Powers** approximation. Recycle continuity theorem, Berry-Esseen, Chernov, etc. \implies **Normal law** and many goodies. . .

(speed of convergence, large deviation fn, local limits)

Two important cases:

- Movable singularity:

$$F(z, u) \approx \left(1 - \frac{z}{\rho(u)}\right)^{-\alpha} \implies \frac{f_n(u)}{f_n(1)} \approx \left(\frac{\rho(1)}{\rho(u)}\right)^n.$$

- Variable exponent:

$$F(z, u) \approx \left(1 - \frac{z}{\rho}\right)^{-\alpha(u)} \implies \frac{f_n(u)}{f_n(1)} \approx \begin{cases} n^{\alpha(u) - \alpha(1)} \\ \left(e^{\alpha(u) - \alpha(1)}\right)^{\log n} \end{cases}.$$

Requires *uniformity* afforded by *Singularity Analysis*
(\neq Tauber or Darboux).

Singularity Perturbation analysis (smoothness)



Uniform Quasi-Powers for coeffs



Normal limit law

EXAMPLE 6. *Polynomials over finite fields.*

```
> Factor(x^7+x+1) mod 29;  
      3      2      2      2  
(x  + x  + 3 x + 15) (x  + 25 x + 25) (x  + 3 x + 14)
```

- \mathcal{P} olynomial is a *Sequence* of coeffs: \mathcal{P} has *Polar singularity*.
- By unique factorization, \mathcal{P} is also *Multiset of Irreducibles*:
 \mathcal{I} has *log singularity*.

\Rightarrow *Prime Number Theorem for Polynomials* $I_n \sim \frac{q^n}{n}$.

- Marking number of \mathcal{I} -factors is approx u th power:

$$P(z, u) \approx \left(e^{I(z)} \right)^u.$$

Variable Exponent \Rightarrow *Normality of # of irred. factors.*

(cf Erdős-Kac for integers.)

□

(Analysis of algorithms, [F-Gourdon-Panario])

EXAMPLE 7. *Patterns in Random Strings*

= Perturbation of linear system of eqns.

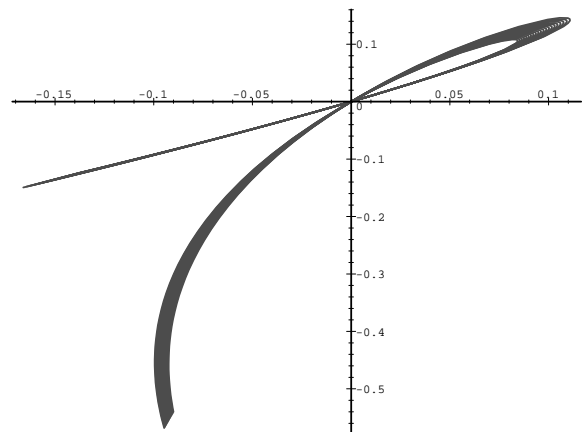
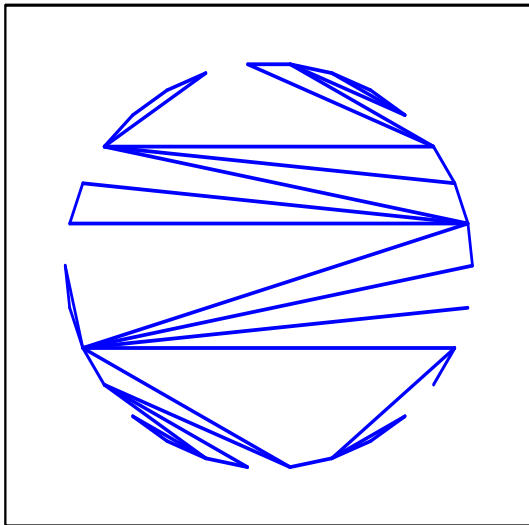
(& many problems with finite automata, paths in graphs)

Linear system $X = X_0 + TX$ w/ Perron-Frobenius.

Auxiliary mark u induces smooth singularity displacement.

For “natural” problems: **Normal limit law.**

Also sets of patterns; similarly for patterns in increasing labelled trees, in permutations, etc. □

**EXAMPLE 8.** *Non crossing graphs.* [F-Noy]

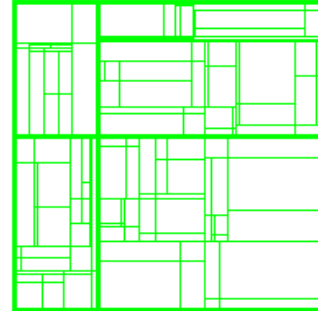
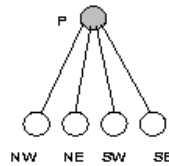
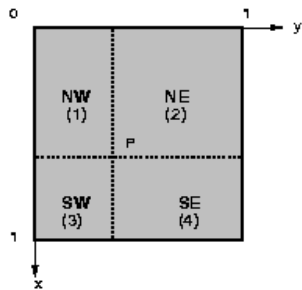
= Perturbation of algebraic equation.

$$G^3 + (2z^2 - 3z - 2)G^2 + (3z + 1)G = 0$$

$$G^3 + (2u^3z^2 - 3u^2z + u - 3)G^2 + (3u^2 - 2u + 3)G + u - 1 = 0$$

Movable singularity scheme applies: **Normality.**

+ Patterns in context-free languages, in combinatorial tree models, in functional graphs [Cf. Drmota]. □



EXAMPLE 9. *Profile of Quadtrees.*

$$F(z, u) = 1 + 2^3 u \int_0^z \frac{dx_1}{x_1(1-x_1)} \int_0^{x_1} \frac{dx_2}{1-x_2} \int_0^{x_2} F(x_3, u) \frac{dx_3}{1-x_3}.$$

Solution is of the form $(1-z)^{-\alpha(u)}$ for algebraic branch $\alpha(u)$;
Variable Exponent \Rightarrow **Normality of search costs.** \square

Applies to many linear differential models that then behave like *cycles-in-perms*.

- **Coalescence** of singularities and/or exponents: e.g. **Airy Law** \equiv **Stable** $(\frac{3}{2})$ [BFSS'01; cf also Pemantle].

Conclusion

For combinatorial *counting* and *limit laws*:

Modest technical apparatus & *generic technology*.

High-level for applications, esp., *analysis of algorithms*.

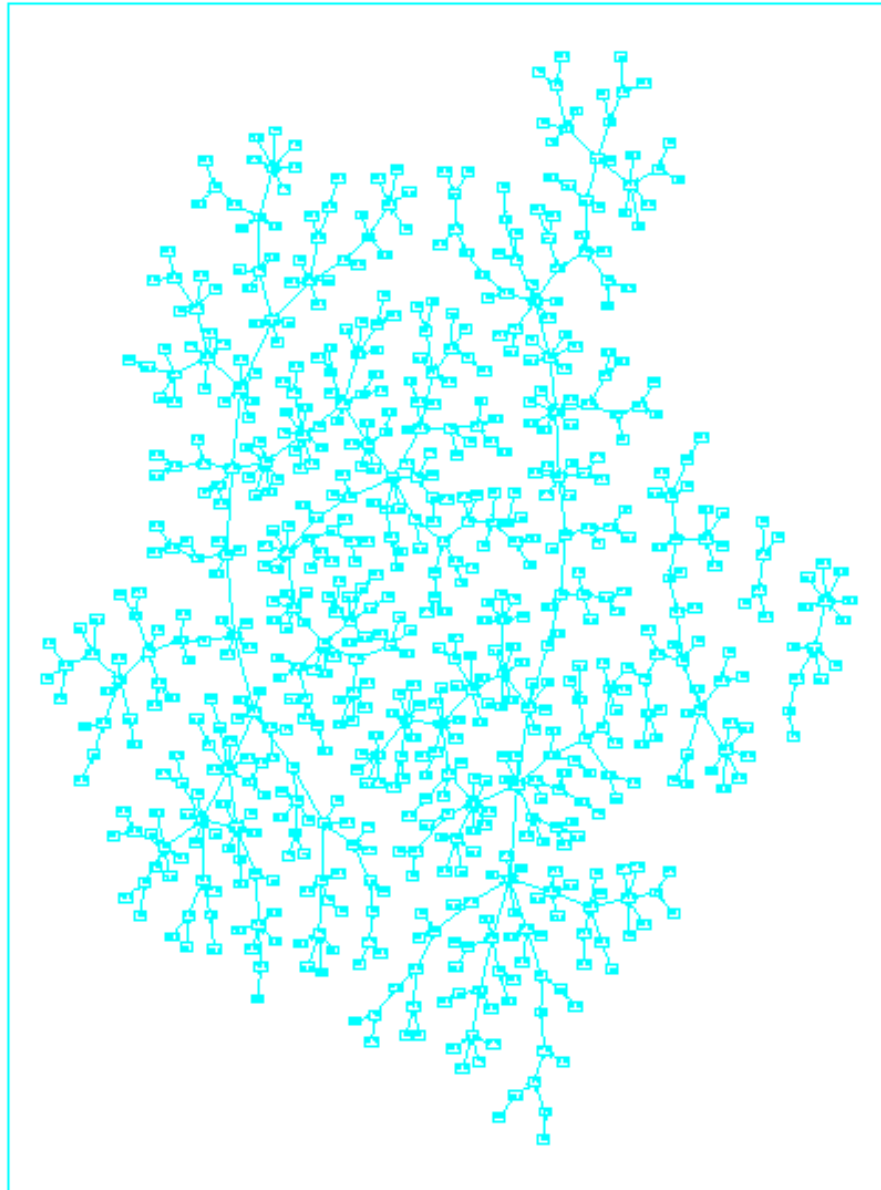
Plug-in on *Symbolic Combinatorics* & Symbolic Computation.

Discussion of *Schemas & Universality* in metric aspects of random discrete structures.

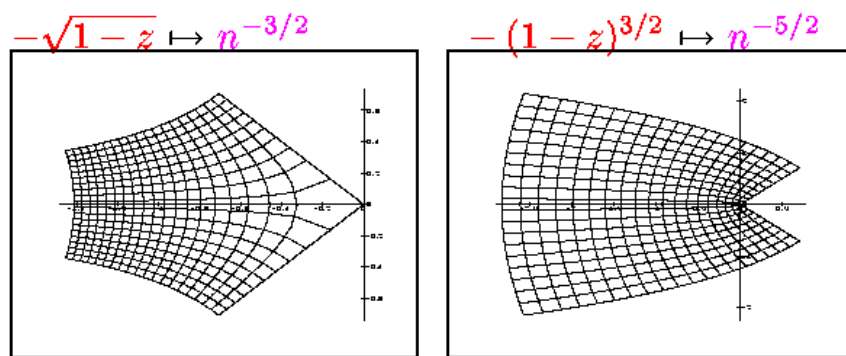
E.g. Borges' theorem for words, trees, labelled trees, mappings, permutations, increasing trees, etc.

THANK YOU!

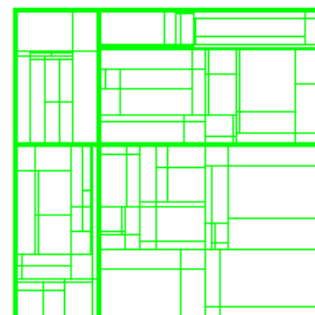
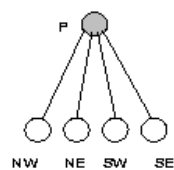
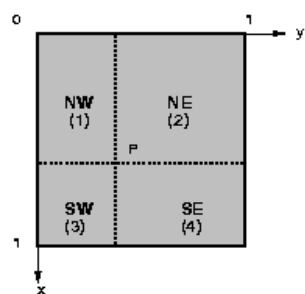




A random functional graph (mapping)



QTrees:



$$f^{\star}(s) := \int_0^{\infty} f^{\star}(x) x^{s-1} dx.$$