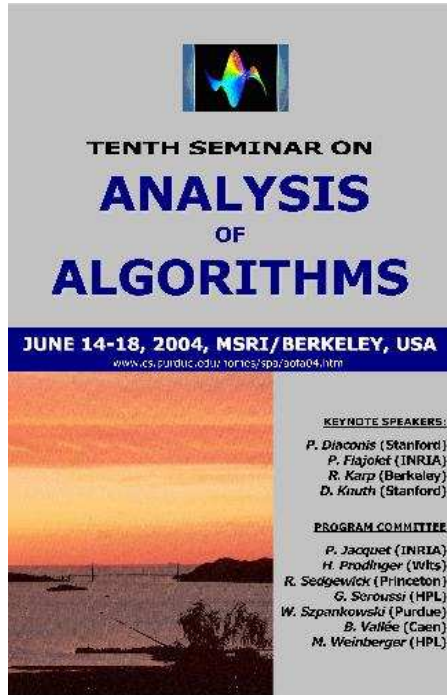
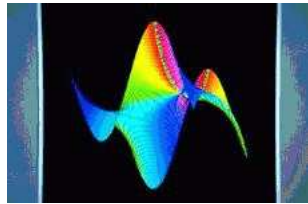


MSRI, Berkeley, June 2004



**TENTH SEMINAR ON  
ANALYSIS  
OF  
ALGORITHMS**

**JUNE 14-18, 2004, MSRI/BERKELEY, USA**  
[www.cs.purdue.edu/homes/sps/aota04.htm](http://www.cs.purdue.edu/homes/sps/aota04.htm)

**KEYNOTE SPEAKERS:**  
*P. Diaconis* (Stanford)  
*P. Flajolet* (INRIA)  
*R. Karp* (Berkeley)  
*D. Knuth* (Stanford)

**PROGRAM COMMITTEE**  
*P. Jacquet* (INRIA)  
*H. Prodinger* (Wits)  
*R. Sedgewick* (Princeton)  
*G. Serrousi* (HPL)  
*W. Szpankowski* (Purdue)  
*B. Vallée* (Caen)  
*M. Weinberger* (HPL)



## Singularity Analysis: A Perspective

PHILIPPE FLAJOLET

(INRIA, France)



Analysis of Algorithms



Average-Case, Probabilistic



Properties of Random Structures?

- Counting and asymptotics

$$n! \sim n^n e^{-n} \sqrt{2\pi n}$$

- Asymptotic laws

$$\Omega_n \xrightarrow{\mathcal{D}} \frac{1}{\sqrt{2\pi}} \int_{-\infty}^x e^{-t^2/2} dt. \quad (\text{e.g., Monkey and typewriter!})$$

— Probabilistic, stochastic

— Analytic Combinatorics: **Generating Functions**

# 1. Introduction

## “Symbolic” Methods

Rota-Stanley; Foata-Schutzenberger; Joyal and UQAM group; Jackson-Goulden, &c; F.; ca 1980 $\pm$ . F-Salvy-Zimmermann 1991  $\leadsto$  *Computer Algebra*.

Basic combinatorial constructions admit of direct translations as operators over generating functions (GF's).

$\mathcal{C}$  : class of comb. structures;

$C_n$  : # objects of size  $n$

$$\begin{array}{l} \text{(counting)} \\ \text{(params)} \end{array} \left\{ \begin{array}{l} C(z) := \sum C_n z^n \\ \widehat{C}(z) := \sum C_n \frac{z^n}{n!} \\ C(z, u) := \sum C_{n,k} z^n u^k \\ \widehat{C}(z, u) := \sum C_{n,k} u^k \frac{z^n}{n!} \end{array} \right.$$

Ordinary GF's for unlabelled structures. Exponential GF's for labelled structures.

## “Dictionaries”

= Constructions viewed as Operators over GF's.

Constr.	Operations	
Union	+	+
Product	×	×
Sequence	$(1 - f)^{-1}$	$(1 - f)^{-1}$
MultiSet	Pólya Exp.	$e^f$
Cycle	Pólya Log.	$\log(1 - f)^{-1}$
	(unlab.)	(lab.)

$$\text{Exp}(f) := \exp\left(f(z) + \frac{1}{2}f(z^2) + \dots\right)$$

$$\text{Log}(f) := \log \frac{1}{1-f(z)} + \dots$$

Books: Goulden-Jackson, Bergeron-LL, Stanley, F-Sedgewick

⇒ How to extract coeff., especially, asymptotically??

## “Complex–analytic Structures”

*Interpret:*

- ♡ Counting GF as analytic transformation of  $\mathbb{C}$ ;
- ♡ Comb. Construction as analytic functional.

**Singularities** are crucial to asymptotic prop's!

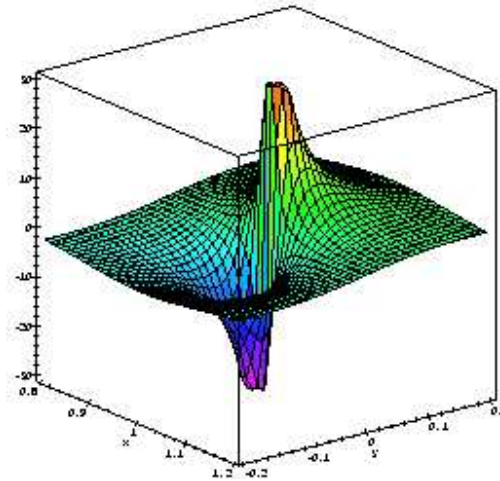
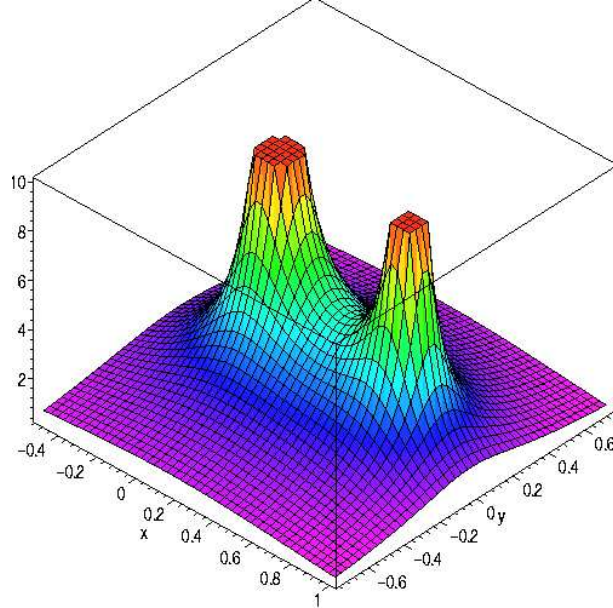
(cf. analytic number theory, complex analysis, etc)

Asymptotic counting via Singularity Analysis (S.A.)

Asymptotic laws via Perturbation + S.A.

$$\frac{1}{2i\pi} \int \frac{1}{1-z-z^2} \frac{dz}{z^{n+1}}$$

$$\Im f(z), \quad f(z) = (1-z-z^2)^{-1}.$$



Refs: F–Odlyzko, SIAM A&DM, 1990  $\ll$  FO82 on tree height; Odlyzko’s 1995 survey in *Handbook of Combinatorics*

+ Banderier, Fill, J. Gao, Gonnet, Gourdon, Kapur, G. Labelle, Lafortest, T. Lafforgue, Noy, Odlyzko, Panario, Poblete, Pouyanne, Prodinger, Puech, Richmond, Robson, Salvy, Schaeffer, Sipala, Soria, Steyaert, Szpankowski, B. Vallée, Viola .

♠ Location of singularity at  $z = \rho$ :  $\text{coeff. } [z^n]f(z) = \rho^{-n} \cdot \text{coeff. } [z^n]f(\rho z)$

♠ Nature of singularity at  $z = 1$ :

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$\frac{1}{(1-z)^2}$	$\longrightarrow$	$n+1$	$\sim$	$n$
$\frac{1}{1-z} \log \frac{1}{1-z}$	$\longrightarrow$	$H_n \equiv \frac{1}{1} + \dots + \frac{1}{n}$	$\sim$	$\log n$
$\frac{1}{1-z}$	$\longrightarrow$	$1$	$\sim$	$1$
$\frac{1}{\sqrt{1-z}}$	$\longrightarrow$	$\frac{1}{2^{2n}} \binom{2n}{n}$	$\sim$	$\frac{1}{\sqrt{\pi n}}$

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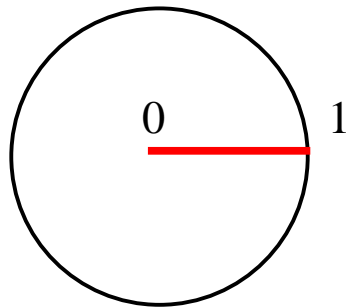
{	Location of sing's :	Exponential factor	$\rho^{-n}$
	Nature of sing's :	“Polynomial” factor	$\vartheta(n)$



Generating Function  $\rightsquigarrow$  Coefficients

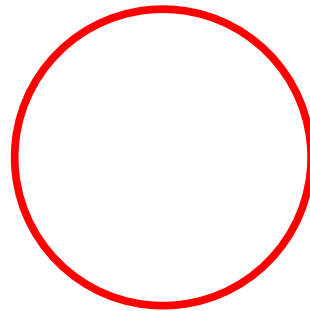
Solving a “Tauberian” problem

Real-Tauberian



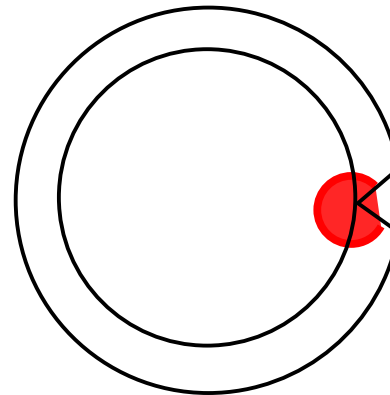
(large  $\implies$  large)

Darboux-Pólya



(smooth  $\implies$  small)

Singularity An.



(Full mappings)

Combinatorial constructions  $\rightsquigarrow$  Analytic Functionals

$\implies$  Analytic continuation prevails for comb. GF's

## 2. Basic Singularity Analysis

**Theorem 1.** *Basic scale translates:*

$$\begin{aligned} \sigma_{\alpha,\beta}(z) &:= (1-z)^{-\alpha} \left( \frac{1}{z} \log \frac{1}{1-z} \right)^\beta \\ \implies [z^n] \sigma_{\alpha,\beta} &\underset{n \rightarrow \infty}{\sim} \frac{n^{\alpha-1}}{\Gamma(\alpha)} (\log n)^\beta. \end{aligned}$$

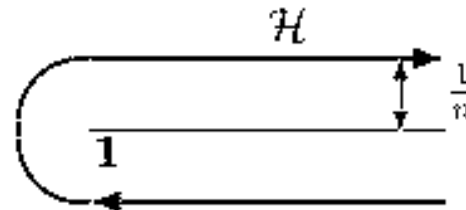
PROOF. Cauchy's coefficient integral,  $f(z) = (1-z)^{-\alpha}$

$$[z^n] f(z) = \frac{1}{2i\pi} \int_{\gamma} f(z) \frac{dz}{z^{n+1}}$$

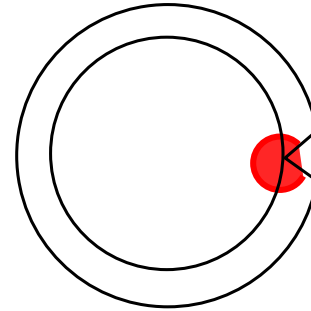
$$\Downarrow \quad \left( z = 1 + \frac{t}{n} \right) \quad \Downarrow$$

$$\frac{1}{2i\pi} \int_{\mathcal{H}} \left( -\frac{t}{n} \right)^{-\alpha} e^{-t} \frac{dt}{n}$$

$$n^{\alpha-1} \times \frac{1}{\Gamma(\alpha)}.$$



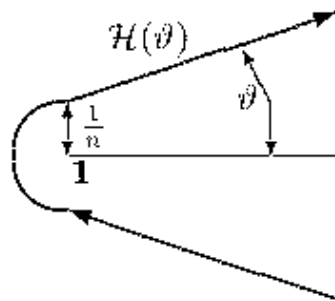
“Camembert”



**Theorem 2.**  *$\mathcal{O}$ -transfers: Under continuation in a  $\Delta$ -domain,*

$$f(z) = O(\sigma_{\alpha,\beta}(z)) \implies [z^n]f(z) = O([z^n]\sigma_{\alpha,\beta}(z)).$$

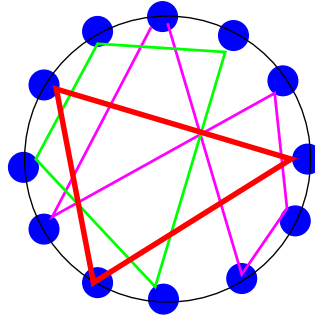
PROOF:



Usage:  $\left\{ \begin{array}{l} f(z) = \lambda\sigma(z) + \mu\tau(z) + \dots + O(\omega(z)) \\ \implies \\ f_n = \lambda\sigma_n + \mu\tau_n + \dots + O(\omega_n). \end{array} \right.$

Similarly: *o*-transfer.

- Dominant singularity at  $\rho$  gives factor  $\rho^{-n}$ .
- Finitely many singularities work fine



EXAMPLE 1. *2-regular graphs* [Comtet] (Originally by Darboux-Pólya.)

$$\mathcal{G} = \mathfrak{M} \left( \frac{1}{2} \mathfrak{e}_{\geq 3}(z) \right)$$

$$\widehat{G}(z) = \exp \left( \frac{1}{2} \log \frac{1}{1-z} - \frac{z}{2} - \frac{z^2}{4} \right)$$

$$\widehat{G}(z) \underset{z \rightarrow 1}{\sim} \frac{e^{-3/4}}{\sqrt{1-z}}$$

$$\frac{G_n}{n!} \underset{n \rightarrow \infty}{\sim} \frac{e^{-3/4}}{\sqrt{\pi n}}$$

□

> equivalent(exp(-z/2-z^2/4)/sqrt(1-z),z,n,4); # By SALVY

$$\frac{\exp(-3/4) (1/n)^{1/2}}{\pi^{1/2}} - \frac{5}{8} \frac{\exp(-3/4) (1/n)^{3/2}}{\pi^{1/2}} + \frac{1}{128} \frac{\exp(-3/4) (1/n)^{5/2}}{\pi^{1/2}}$$

**EXAMPLE 2.** *Richness index of trees* [F-Sipala-Steyaert,90]

= Number of different terminal subtrees. Catalan case:

$$K(z) = \frac{1}{2z} \sum_{k \geq 0} \frac{1}{k+1} \binom{2k}{k} \left( \sqrt{1-4z-4z^{k+1}} - \sqrt{1-4z} \right)$$

$$K(z) \underset{z \rightarrow 1/4}{\approx} \frac{1}{\sqrt{Z \log Z}}, \quad Z := 1 - 4z$$

$$\text{Mean index} \underset{n \rightarrow \infty}{\sim} C \frac{n}{\sqrt{\log n}}, \quad C \equiv \sqrt{\frac{8 \log 2}{\pi}}.$$

= Compact tree representations as DAGs = Common Subexpression Pb. □

## Extensions

- ♡ Slowly varying  $\implies$  slowly varying: Log-log  $\implies$  Log-Log, ...
- ♡ Full asymptotic expansions
- ♡ Uniformity of coefficient extraction  $[z^n]\{F_u(z)\}_{u \in \Omega} = \sim$  later!.
- ♡ Some cases with natural boundary [Fl-Gourdon-Panario-Pouyanne]

**EXAMPLE 3.** *Distinct Degree Factorization* [DDF] in Polynomial Fact  $\leadsto$  Greene–Knuth:

$$[z^n] \prod_{k=1}^{\infty} \left(1 + \frac{z^k}{k}\right).$$

Hybrid w/ Darboux:  $e^{-\gamma} + \frac{e^{-\gamma}}{n} + \dots + \star \frac{(-1)^n}{n^3} + \star \frac{\omega^n}{n^3} + \dots$  □

Cf. Hardy-Ramanujan's partition analysis "without contrast".

### 3. Closure Properties

Function of S.A.-type = amenable to singularity analysis

- is continuable in a  $\Delta$ -domain,
- admits singular expansion in scale  $\{\sigma_{\alpha,\beta}\}$ .

**Theorem 3.** *Generalized polylogarithms*

$$\text{Li}_{\alpha,k} := \sum (\log n)^k n^{-\alpha} z^n$$

*are of S.A.-type.*

PROOF. Cauchy-Lindelöf representations

$$\sum \varphi(n)(-z)^n = -\frac{1}{2i\pi} \int_{1/2-i\infty}^{1/2+i\infty} \varphi(s) z^s \frac{\pi}{\sin \pi s} ds.$$

+ Mellin transform techniques (Ford, Wong, F.).



EXAMPLE 4. *Entropy of Bernoulli distribution*

$$H_n := - \sum_k \pi_{n,k} \log \pi_{n,k}, \quad \pi_{n,k} \equiv \binom{n}{k} p^k (1-p)^{n-k}$$

$$\text{involves } \sum \log(k!) z^k = (1-z)^{-1} \text{Li}_{0,1}(z)$$

$$\frac{1}{2} \log n + \frac{1}{2} + \log \sqrt{2\pi p(1-p)} + \dots$$

Redundancy, coding, information th.; [Jacquet-Szpankowski](#) via Analytic dePoissonization. □

- Elements like  $\log n, \sqrt{n}$  in combinatorial sums

**Theorem 4.** *Functions of S.A.-type are closed under integration and differentiation.*

PROOF. Adapt from Olver, Henrici, etc.

**Theorem 5.** *Functions of S.A.-type are closed under Hadamard product*

$$f(z) \odot g(z) := \sum_n (f_n g_n) z^n.$$

PROOF. Start from Hadamard's formula

$$f(z) \odot g(z) = \frac{1}{2i\pi} \int_{\gamma} f(t) g\left(\frac{w}{t}\right) \frac{dt}{t}.$$

+ adapt Hankel contours [H., Jungen, R. Wilson  $\rightsquigarrow$  **Fill-F-Kapur**]

EXAMPLE 5. *Divide-and-conquer recurrences*

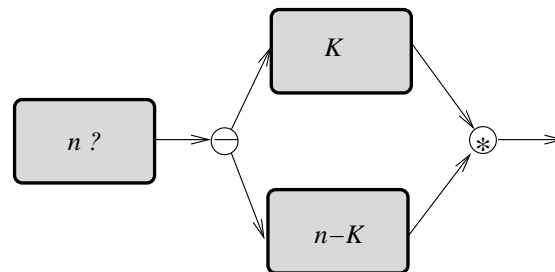
$$f_n = t_n + \sum \pi_{n,k}(f_k + f_{n-k})$$

$$\text{Sing}(f(z)) = \Phi(\text{Sing}(t(z)))$$

$$\text{Asympt}[f_n] = \Psi(\text{Sing}(t)).$$

E.g., Catalan statistics: need  $\sum \binom{2n}{n} \log n \cdot z^n$ .

Useful in random tree applications [Fill-F-Kapur, 2004<sup>+</sup>, Fill-Kapur] //  
Neininger-Hwang *et al.*  $\ll$  Knuth-Pittel. Moments  $\leftrightarrow$  contraction method  
[Rösler-Rüschendorf-Neininger] □



## 4. Functional Equations

- Rational functions. Linear system  $\mathbb{Q}_{\geq 0}[z]$  implies **polar singularities**:

$$[z^n]f(z) \approx \sum \omega^n n^k, \quad \omega \in \overline{\mathbb{Q}}, \quad k \in \mathbb{Z}_{\geq 0}.$$

+ irreducibility: Perron-Frobenius  $\implies$  *simple dom. pole*.

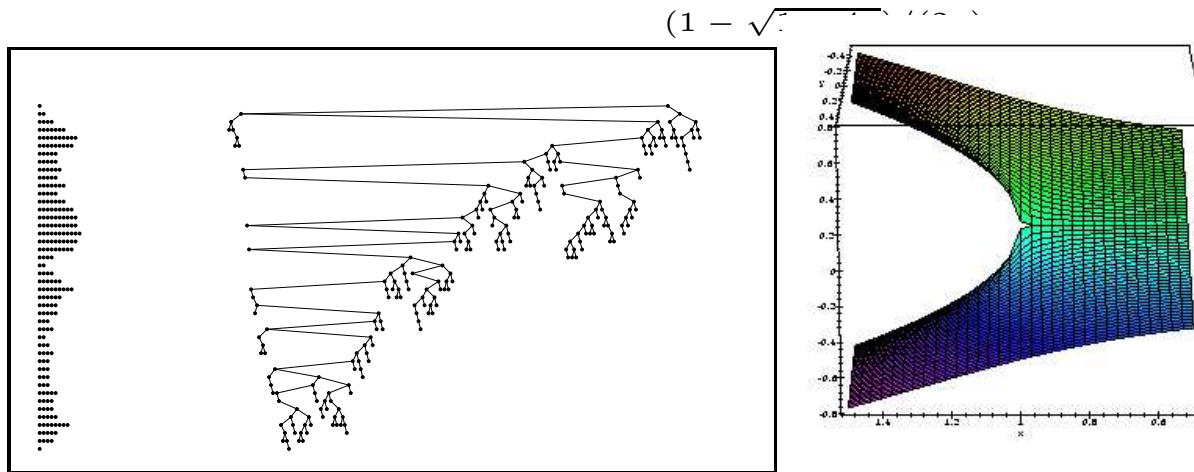
- Word problems from regular language models;
  - Transfer matrices [Bender-Richmond]: dimer in strip, knights, etc.
- $\rightsquigarrow$  Vallée's generalization to dynamical sources via transfer operators.

- Algebraic functions, by Puiseux expansions  $(Z^{p/q}) \ll$  S.A. or Darboux!

$$[z^n]f(z) \approx \sum \sum \omega^n n^{p/q}, \quad \omega \in \overline{\mathbb{Q}}, \quad p/q \in \mathbb{Q},$$

*Asymptotics of coeff. is decidable [Chabaud-F-Salvy].*

- Word problems from context-free models;
  - Trees; Geom. configurations (non-crossing graphs, polygonal triang.);
- Planar Maps [Tutte...]; Walks [Banderier Bousquet-M., Schaeffer], ...



Square-root singularity is “**universal**” for many recursive classes = controlled “failure” of Implicit Function Theorem  $Z \propto Y^2$   
 Entails coeff. asymptotic  $\approx \omega^n n^{-3/2}$  with **critical exponent**  $-3/2$ .

E.g., unbalanced 2–3 trees (Meir-Moon):  $f = z\phi(f)$ ,  $\phi(u) = 1 + u^2 + u^3$ .  
 Pólya’s combinatorial chemistry programme:

$$f(z) = z \text{Exp}(f(z)) \equiv z e^{f(z) + \frac{1}{2}f(z^2) + \frac{1}{3}f(z^3) + \dots}$$

Starting with Pólya 1937; Otter 1949; Harary-Robinson et al. 1970’s;  
 Meir-Moon 1978; Bender/Meir-Moon; **Drmotá-Lalley-Woods thm.** 1990<sup>+</sup>

- “Holonomic” functions. Defined as solutions of linear ODE’s with coeffs in  $\mathbb{C}(z)$  [Zeilberger]  $\equiv \mathcal{D}$ -finite.

$$\mathcal{L}[f(z)] = 0, \quad \mathcal{L} \in \mathbb{C}(z)[\partial_z].$$

- Stanley, Zeilberger, Gessel: Young tableaux and permutation statistics; regular graphs, constrained matrices, etc.

Fuchsian case (or “regular” singularity) ( $Z^\beta \log^k Z$ ):

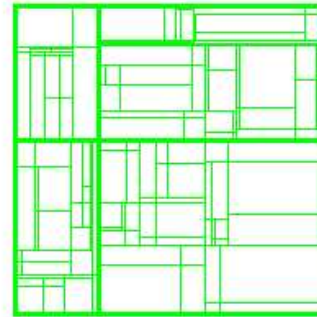
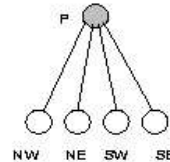
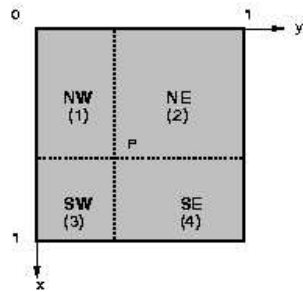
$$[z^n]f(z) \approx \sum \omega^n n^\beta (\log n)^k, \quad \omega, \beta \in \overline{\mathbb{Q}}, \quad k \in \mathbb{Z}_{\geq 0}.$$

S.A. applies automatically to classical classification.

Asymptotics of coeff is decidable

- general case: modulo oracle for connection problem;
- strictly positive case: “usually” OKay.

QTrees:



**EXAMPLE 6.** *Quadtrees—Partial Match* [FGPR'92]

Divide-and-conquer recurrence with coeff. in  $\mathbb{Q}(n)$

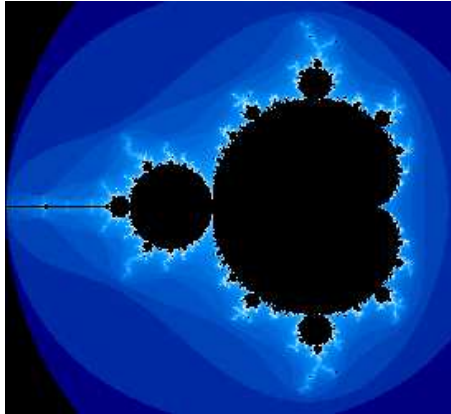
Fuchsian equation of order  $d$  (dimension) for GF

$$Q_n^{(d=2)} \asymp n^{(\sqrt{17}-3)/2}.$$

E.g.,  $d = 2$ : Hypergeom  ${}_2F_1$  with algebraic arguments. □

Extended by Hwang et al. Cf also Hwang's *Cauchy ODE* cases.

Panholzer-Prodinger+Martinez, ...



- Functional Equations and Substitution.

- Early example of *balanced 2-3 trees* by Odlyzko, 1979.

$$T(z) = z + T(\tau(z)), \quad \tau(z) := z^2 + z^3.$$

Infinitely many exponents with common real part implies periodicities:

$$T_n \sim \frac{\phi^n}{n} \Omega(\log n).$$



- Singular iteration for *height of trees* (binary and other simple varieties; F-Gao-Odlyzko-Richmond; cf Renyi-Szekeres):

$$y_h = z + y_{h-1}^2, \quad y_0 = z.$$

— Moments and convergence in law; Local limit law of  $\vartheta$ -type.

Applies to branching processes conditioned on total progeny.

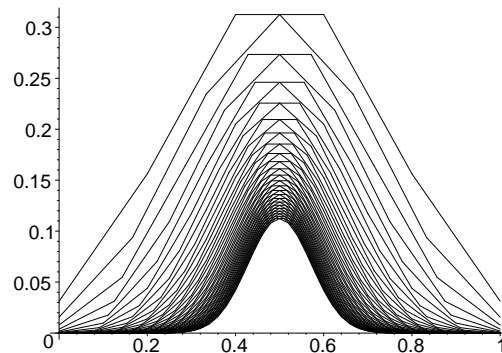
Cf Chassaing-Marckert for // probabilistic approaches  $\rightsquigarrow$  **width**

- *Digital search trees* via  $q$ -hypergeometrics: singularities accumulate geometrically  $\rightsquigarrow$  **periodicities** [F-Richmond]:

$$\partial_z^k f(z) = t(z) + 2e^{z/2} f\left(\frac{z}{2}\right).$$

- *Order of binary trees* (Horton-Strahler, Register function; F-Prodinger) via **Mellin** tr. of GF and & singularities.

## 5. Limit Laws



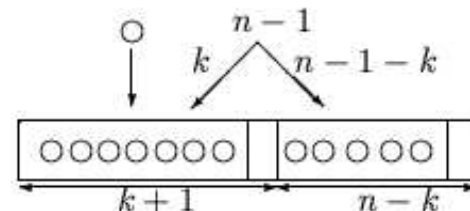
- Moment pumping from bivariate GF

Early theories by Kirschenhofer-Prodinger-Tichy (1987)

Factorial moment of order  $k$ :  $[z^n] \left( \frac{\partial}{\partial k} F(z, u) \right)_{u=1}$

**EXAMPLE 7.** *Airy distribution of areas* shows up in area below paths, path length in trees, Linear Probing Hashing, inversions in increasing trees, connectivity of graphs.

$$\frac{\partial}{\partial z} F(z, q) = F(z, q) \cdot \frac{F(z, q) - qF(qz, q)}{1 - q}$$



Louchard-Takács<sup>[Darboux]</sup>; Knuth; F-Poblete-Viola // Chassaing-Marckert □

Classical probability theory: sums of Random Variables  $\leadsto$  powers of fixed function (PGF, Fourier tr.)  $\leadsto$  Normal Law.

For problems expressed by Bivariate GF (BGF): field founded by E. Bender *et al.* + developments by F, Soria, Hwang, ...

Idea: BGF  $F(z, u) = \sum f_n(u)z^n$ , where  $f_n(u)$  describes parameter on objects of size  $n$ . If (for  $u$  near 1)

$$f_n(u) \approx \omega(u)^{\kappa_n}, \quad \kappa_n \rightarrow \infty,$$

then speak of Quasi-Powers approximation. Recycle continuity theorem, Berry-Esseen, Chernov, etc.  $\implies$  Normal law and many goodies...

(speed of convergence, large deviation fn, local limits)

Two important cases:

- Movable singularity:

$$F(z, u) \approx \left(1 - \frac{z}{\rho(u)}\right)^{-\alpha} \implies \frac{f_n(u)}{f_n(1)} \approx \left(\frac{\rho(1)}{\rho(u)}\right)^n.$$

- Variable exponent:

$$F(z, u) \approx \left(1 - \frac{z}{\rho}\right)^{-\alpha(u)} \implies \frac{f_n(u)}{f_n(1)} \approx \begin{cases} n^{\alpha(u) - \alpha(1)} \\ \left(e^{\alpha(u) - \alpha(1)}\right)^{\log n} \end{cases}.$$

Requires *uniformity* afforded by *Singularity Analysis*  
( $\neq$  Tauber or Darboux).

Singularity Perturbation analysis (smoothness)



Uniform Quasi-Powers for coeffs



*Normal limit law*

EXAMPLE 8. *Polynomials over finite fields.*

> Factor( $x^7+x+1$ ) mod 29;

$$(x^3 + x^2 + 3x + 15) (x^2 + 25x + 25) (x^2 + 3x + 14)$$

- Polynomial is a *Sequence* of coeffs:  $\mathcal{P}$  has *Polar singularity*.
- By unique factorization,  $\mathcal{P}$  is also *Multiset of Irreducibles*:  
 $\mathcal{I}$  has *log singularity*.

$\implies$  *Prime Number Theorem for Polynomials*  $I_n \sim \frac{q^n}{n}$ .

- Marking number of  $\mathcal{I}$ -factors is approx *u*th power:

$$P(z, u) \approx \left( e^{I(z)} \right)^u.$$

*Variable Exponent*  $\implies$  *Normality of # of irred. factors.*

(cf Erdős-Kac for integers.)

□

(Analysis of polynomial fact. algorithms, [F-Gourdon-Panario])

ACCGATCATTAGCAGATTATCATTTACTGAGAGTACTTAACATGCCA

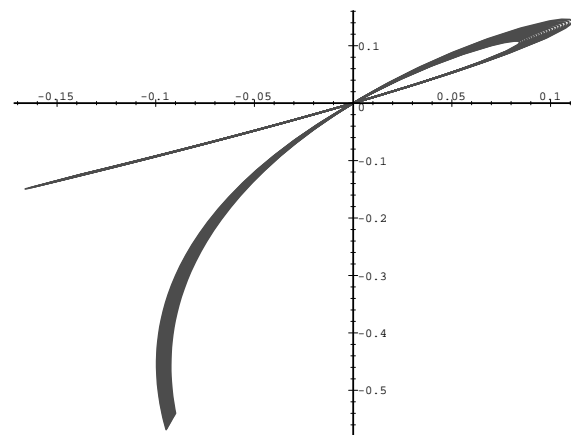
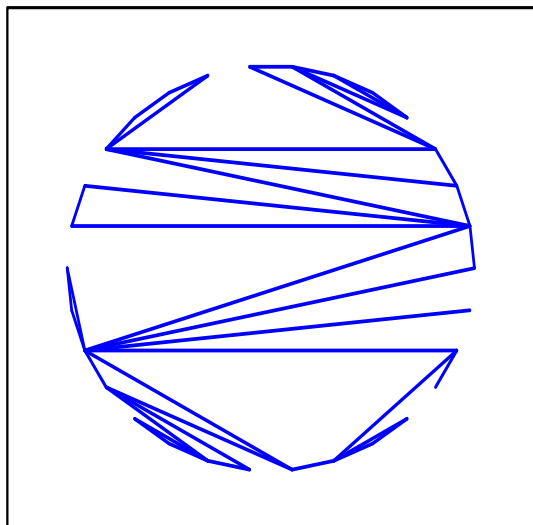
**EXAMPLE 9.** *Patterns in Random Strings* = Perturbation of linear system of eqns. (& many problems with finite automata, paths in graphs)

Linear system  $X = X_0 + \mathbf{T}X$  w/ Perron-Frobenius. Auxiliary mark  $u$  induces smooth singularity displacement. For “natural” problems:

*Normal limit law.* cf [Régnier & Szpankowski], ... □

Also sets of patterns; similarly for patterns in increasing labelled trees, in permutations, in *binary search trees* [F-Gourdon-Martinez].

Generalized patterns and/or sources by Szpankowski, Vallée, ...



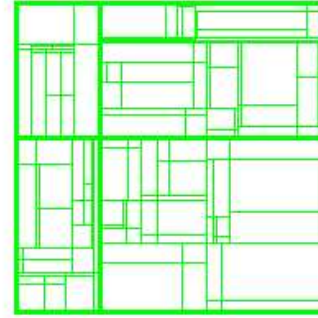
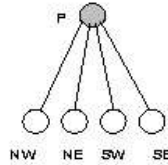
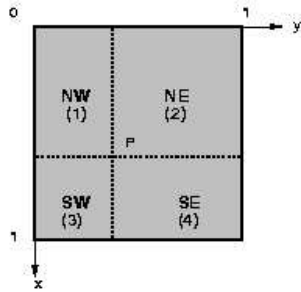
**EXAMPLE 10.** *Non crossing graphs.* [F-Noy]  
 = Perturbation of algebraic equation.

$$G^3 + (2z^2 - 3z - 2)G^2 + (3z + 1)G = 0$$

$$G^3 + (2u^3z^2 - 3u^2z + u - 3)G^2 + (3u^2 - 2u + 3)G + u - 1 = 0$$

Movable singularity scheme applies: **Normality.**

+ Patterns in context-free languages, in combinatorial tree models, in functional graphs: Drmota's version of Drmota-Lalley-Woods. □



EXAMPLE 11. *Profile of Quadtrees.*

$$F(z, u) = 1 + 2^3 u \int_0^z \frac{dx_1}{x_1(1-x_1)} \int_0^{x_1} \frac{dx_2}{1-x_2} \int_0^{x_2} F(x_3, u) \frac{dx_3}{1-x_3}.$$

Solution is of the form  $(1-z)^{-\alpha(u)}$  for algebraic branch  $\alpha(u)$ ;

Variable Exponent  $\implies$  Normality of search costs. □

Applies to many linear differential models that behave like *cycles-in-perms*.





**EXAMPLE 12.** *Urn models.*  $2 \times 2$ -balanced.

$$(u^5 z - u) \frac{\partial G}{\partial z} + (1 - u^6) \frac{\partial G}{\partial u} + u^5 G = 0$$

[FGP'03]  $\leftrightarrow$  Janson, Mahmoud, Puyhaubert,  
Panholzer-Proding, ...  $\square$

- *Coalescence of singularities and/or exponents*: e.g. Maps  
= Airy Law  $\equiv$  Stable( $\frac{3}{2}$ ) [BFSS'01]. Cf Pemantle, Wilson, Lladser,  
.....

# Conclusions

For combinatorial **counting** and **limit laws**:

Modest technical apparatus & *generic technology*.

High-level for applications, esp., *analysis of algorithms*.

Plug-in on *Symbolic Combinatorics & Symbolic Computation*.

Discussion of *Schemas & “Universality”* in metric aspects of random discrete structures.

E.g. Borges’ theorem for words, trees, labelled trees, mappings, permutations, increasing trees, maps, etc.

THANK YOU!

