

ASIAN'04,
Chiang Mai 2004

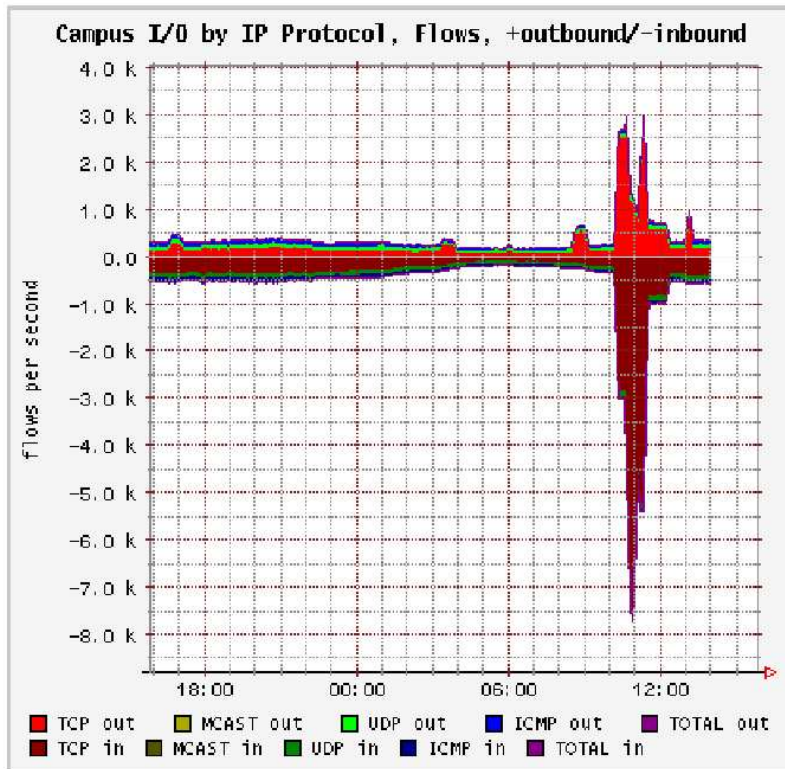


Counting by Coin Tossings

Philippe Flajolet, INRIA, France

<http://algo.inria.fr/flajolet>

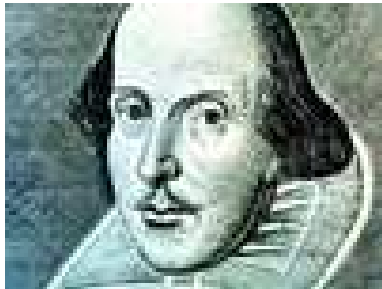
From Estan-Varghese-Fisk: traces of attacks
Need number of active connections in time slices.



Incoming/Outgoing flows at 40Gbits/second.
Code Red Worm: 0.5GBytes of compressed data per hour (2001).
CISCO: in 11 minutes, a worm infected 500,000,000 machines.

The situation is like listening to a play of Shakespeare and at the end *estimate the number of different words*.

Rules: Very little computation per element scanned, very little auxiliary memory.



From Durand-Flajolet, LogLog Counting (ESA-2003):

Whole of Shakespeare, $m = 256$ small “bytes” of 4 bits each = 128bytes

```
ghfffghfghgghggggghghheehfhfhhgghghghhfgffffhhhiigfhhffgfiihfhhh  
igigighfgihfffghigihghigfhhgeegeghgghhhgghhfhidiigihighihehhffgg  
hfgighigffghdieghhhggghhfgghfiieffghghihifgggffihgihfggighgiiif  
fjgfgjhhjiihfjhgehghfhfhjhiggghghihigghhihihgiighgfhlgjfgjjmfl
```

Estimate $n^\circ \approx 30,897$ vs $n = 28,239$ distinct words. Error: +9.4% w/ 128 bytes!

Uses:

- **Routers:** intrusion, flow monitoring & control
- **Databases:** Query optimization, cf $\mathfrak{M} \cup \mathfrak{M}'$ for *multisets*; Estimating the size of queries & “sketches”.
- **Statistics gathering:** on the fly, fast and with little memory even on “unclean” data \simeq layer 0 of “*data mining*”.

This talk:

- Estimating characteristics of large data streams
 - sampling; size & cardinality & nonuniformity index (F_1, F_0, F_2)
 - ↷ power of randomization via hashing
 - ◇ Gains by a factor of >400 (Palmer *et al.*)
- Analysis of algorithms
 - generating functions, complex asymptotics, Mellin transforms
 - ◇ Nice problems for theoreticians.
- Theory and Practice
 - Interplay of analysis and design ↷ super-optimized algorithms.

1 **PROB. ALG. ON STREAMS**

Given: S = a large stream $S = (r_1, r_2, \dots, r_\ell)$ with duplicates

— $\|S\|$ = **length** or **size**: total # of records (ℓ)

— $|S|$ = **cardinality**: # of *distinct* records (c)

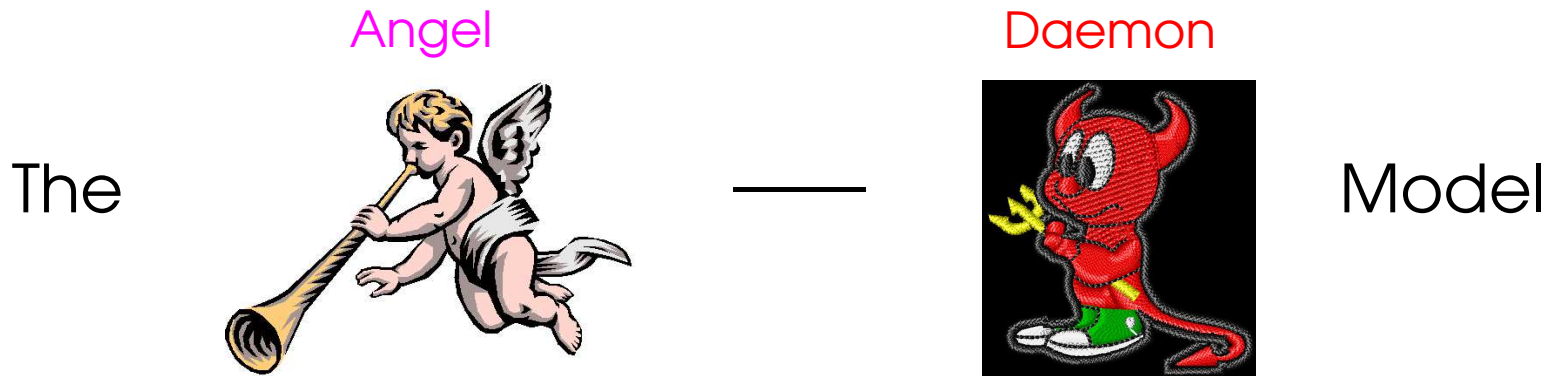
◇ How to estimate size, cardinality, etc?

More generally, if f_v is frequency of value v : $F_p := \sum_{v \in \mathbb{D}} (f_v)^p$.

Cardinality is F_0 ; size is F_1 ; F_2 is indicator of nonuniformity of distribution;
“ F_∞ ” is most frequent element (Alon, Matias, Szegedy, STOC96)

◇ How to sample?

— with or without multiplicity



Pragmatic assumptions/ Engineer's point of view:
 Can get random bits from data: **Works fine!**

(A1) There exists a "good" hash function

$$h : \mathcal{D} \mapsto \mathcal{B} \equiv \{0, 1\}^L$$

Data domain \mapsto Bits

Typically: $L = 30\text{--}32$ (more or less, maybe).

$$h(x) := \lambda \cdot \langle x \text{ in base } B \rangle \bmod p$$

Sometimes, also: (A2) There exists a "good" pseudo-random number gen. $T : \mathcal{B} \mapsto \mathcal{B}$, s.t. iterates $Ty_0, T^{(2)}y_0, T^{(3)}y_0, \dots$ look random. ($T(y) := (a \cdot y \bmod p)$)

Two preparatory examples.

Let a flow of people enter a room.

— *Birthday Paradox*: It takes on average 23 to get a **birthday collision**

— *Coupon Collector*: After 365 persons have entered, expect a **partial collection** of ~ 231 different days in the year; it would take more than 2364 to reach a full collection.

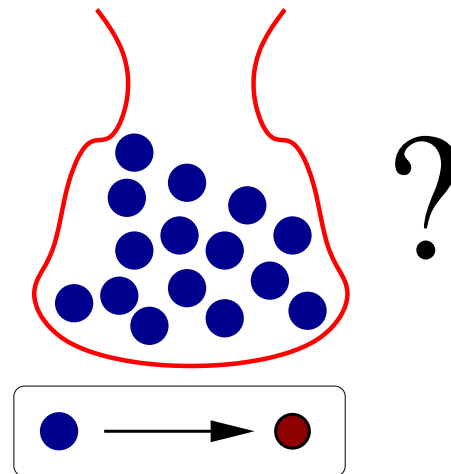
B 1st birthday coll.	n	C complete coll.
$\mathbb{E}_n(B) \sim \sqrt{\frac{\pi n}{2}}$	$\approx ne^{-1}$	$\mathbb{E}_n(C) = nH_n \sim n \log n$

Suppose we *didn't know* the number N of days in the year but could identify people with the same birthday. Could we *estimate* N ?

1.1 Birthday paradox counting

- A warm-up “abstract” example due to Brassard-Bratley (Book 1996) = a *Gedanken* experiment.

How to weigh an urn by shaking it?



Urn contains unknown number N of balls.

♠ Deterministic: Empty it one by one: cost is $O(N)$.

♥ Probabilistic $O(\sqrt{N})$: (shake, draw, paint)*; stop!

ALG: Birthday Paradox Counting

Shake, pull out a ball, mark it with paint;

repeat until draw an already marked ball.

Infer N from T = number of steps.

We have $\mathbb{E}(T) \sim \sqrt{\pi N/2}$ by Birthday Paradox.

• Invert and try $X := \frac{2}{\pi}T^2$. Estimate is biased

•• Analyse 2nd moment of BP, find $\mathbb{E}(T^2) \sim 2N$ and propose $X := T^2/2$. Estimate is now (asymptotically) unbiased.

••• Wonder about accuracy: Standard Error := $\frac{\text{Std Deviation of estimate } (X)}{\text{Exact value } (N)}$.

\leadsto Need to analyse fourth moment $\mathbb{E}(T^4)$. Do maths:

$$\mathbb{E}_N(T^{2r}) = 2^r r! N^r, \quad \mathbb{E}_N(T^{2r+1}) = (1 \cdot 3 \cdots (2r-1)) \sqrt{\frac{\pi}{2}} N^{r+\frac{1}{2}}.$$

$\implies E(T^4) \sim 8N^2$. Standard error \implies Estimate $\in (0, 3N)$. ($N = 10^6$): 384k; 3,187k; 635k; 29k; 2,678k; 796k; 981k, ...

•••• Improve algorithm. Repeat m times and average.

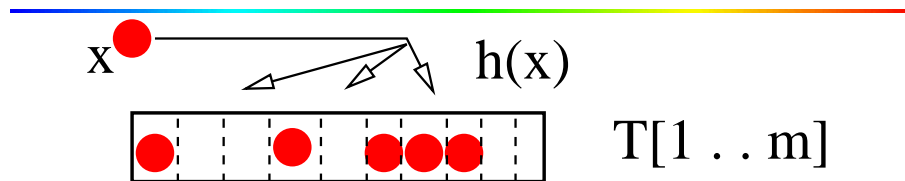
\leadsto Time cost: $O(m\sqrt{N})$ for accuracy $O\left(\frac{1}{\sqrt{m}}\right)$.

Shows usefulness of maths: Ramanujan's $Q(n)$ function, Laplace's method for sums or integrals (cf Knuth, Vol 1); singularity analysis. . .

1.2 Coupon Collector Counting

First Counting Algorithm: Estimate cardinalities \equiv # of **distinct** elements.

This is real CS, motivated by *query optimization in data bases*. (Whang et al, ACM TODS 1990)



ALG: Coupon Collector Counting

Given multiset $S = (s_1, \dots, s_\ell)$; Estimate $\text{card}(S)$?

Set up a table $T[1..m]$ of m bit-cells.

— for x in S do mark cell $T[h(x)]$;

Return $-m \log V$, where $V :=$ fraction of empty cells.

Simulate hashing table; Alg. is indep. of replications.

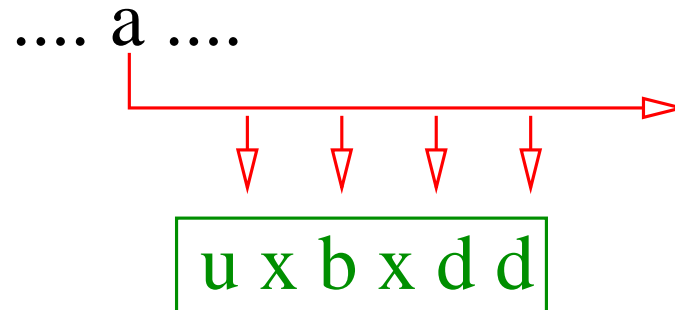
Let n be sought cardinality. Then $\alpha := n/m$ is *filling ratio*. Expect $V \approx e^{-\alpha}$ empty cells by classical analysis of occupancy. Distribution is concentrated. Invert!

Count cardinalities till N_{\max} using $\frac{1}{10} N_{\max}$ bits, for accuracy (standard error) = 2%.

Generating functions for occupancy; Stirling numbers; basic depositions.

2 SAMPLING

A very classical problem (Vitter, ACM TOMS 1985)



ALG: Reservoir Sampling (*with multiplicities*)

Sample m elements from $S = (s_1, \dots, s_N)$; (N unknown a priori)

Maintain a cache (reservoir) of size m ;

— for each coming s_{t+1} :

place it in cache with probability $m/(t+1)$; drop random element;

Math: Need analysis of skipping probabilities.

Complexity of Vitter's best alg. is $O(m \log N)$.

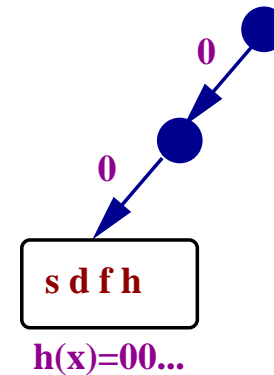
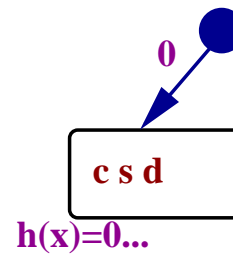
Useful for building "sketches", order-preserving H-fns & DS.

Can we sample *values* (i.e., without multiplicity)?

Algorithm due to (Wegman, ca 1984, unpub.), analysed by (F.1990).

Sample of size $\leq b$:
depth $d = 0, 1, 2, \dots$

c x a s d



ALG: Adaptive Sampling (*without multiplicities*)

Get a sample of size m from S 's values.

Set $b := 4m$ (bucket capacity);

— Oversample by adaptive method;

– Get sample of m elements from the ($b \equiv 4m$) bucket.

Analysis.

View collection of records as a set of bitstrings.

Digital tree aka trie, paged version:

$$\left[\begin{array}{l} \text{Trie}(\omega) \equiv \omega \text{ if } \text{card}(\omega) \leq b \\ \text{Trie}(\omega) = \overbrace{\text{Trie}(\omega \setminus 0) \quad \text{Trie}(\omega \setminus 1)}^{\bullet} \text{ if } \text{card}(\omega) > b \end{array} \right.$$

(Underlies dynamic and extendible hashing, paged DS, etc)

Refs: (Knuth Vol 3), (Sedgewick, Algorithms), Books by Mahmoud, Szpankowski. General analysis by (Clément-F-Vallée, Alg. 2001), etc.

Depth in Adaptive Sampling is length of leftmost branch;
Bucket size is # of elements in leftmost page.

For recursively defined parameters: $\alpha[\omega] = \beta[\omega \setminus 0]$:

$$\mathbb{E}_n(\alpha) := \sum_{k=0}^n \frac{1}{2^n} \binom{n}{k} \mathbb{E}_k(\beta).$$

Introduce *exponential generating functions (EGF)*:

$A(z) := \sum_n \mathbb{E}_n(\alpha) \frac{z^n}{n!}$ &c. Then $A(z) = e^{z/2} B\left(\frac{z}{2}\right)$.

For recursive parameter ϕ : $\Phi(z) = e^{z/2} \Phi\left(\frac{z}{2}\right) + \text{Init}(z)$

Solve by iteration, extract coefficients; Mellin-ize \leadsto later!

Bonus: Second Counting Algorithm for *cardinalities*.

Let $d :=$ sampling *depth*; $\xi :=$ sample size.

Theorem [F90]: $X := 2^d \xi$ estimates the cardinality of S using b words of memory, in a way that is unbiased and with standard error $\approx 1.20/\sqrt{b}$.

- $1.20 \doteq 1/\sqrt{\log 2}$: with $b = 1,000W$, get 4% accuracy.
- Distributional analysis by (Louchard RSA 1997).
- Related to folk algorithm for **leader election** on channel: “*Talk, flip coin if noisy; sleep if Tails; repeat!*”
- Related to “**tree protocols with counting**”
» Ethernet. Cf (Greenberg-F-Ladner JACM 1987).

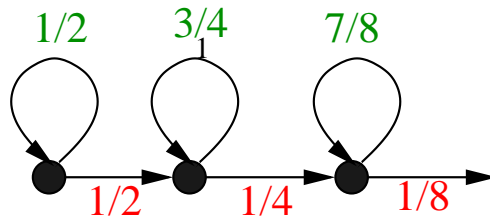
3 APPROXIMATE COUNTING

The oldest algorithm (Morris CACM:1977), analysis (F, 1985).

Maintain F_1 , i.e., counter subject to $C := C + 1$.

Theorem: Count till n probabilistically using $\log_2 \log n + \delta$ bits, with accuracy about $0.59 \cdot 2^{-\delta/2}$.

Beats information theory(!?): 8 bits for counts $\leq 2^{16}$ w/ accuracy $\approx 15\%$.



ALG: Approximate Counting

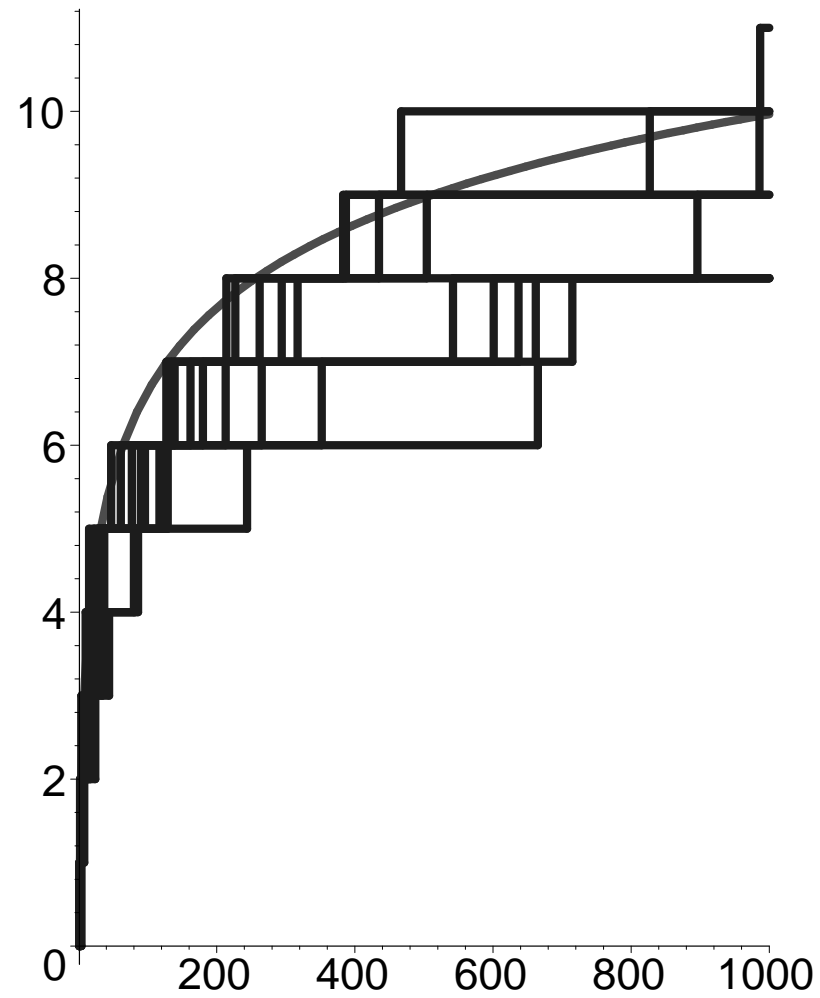
Initialize: $X := 1$;

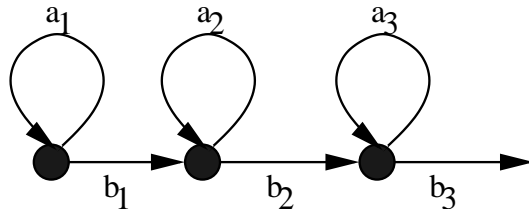
Increment: do $X := X + 1$ with probability 2^{-X} ;

Output: $2^X - 2$.

In base $q < 1$: increment with probability q^X ; output $(q^{-x} - q^{-1}) / (q^{-1} - 1)$; use $q = 2^{-2^{-\delta}} \approx 1$.

10 runs of of APCO: value of X ($n = 10^3$)





Methodology:

♡ Paths in graphs \mapsto Generating Functions:

$$(f_n) \mapsto f(z) := \sum_n f_n z^n.$$

Here: Symbolically describe all paths:

$$\frac{(a_1)^* b_1 (a_2)^* b_2 (a_3)^*}{\frac{1}{1-a_1} b_1 \frac{1}{1-a_2} b_2 \frac{1}{1-a_3}} \quad \text{since} \quad \frac{1}{1-f} = 1 + f + f^2 + \dots \simeq (f)^*.$$

Perform probabilistic valuation $a_j \mapsto q^j$; $b_j \mapsto 1 - q^j$:

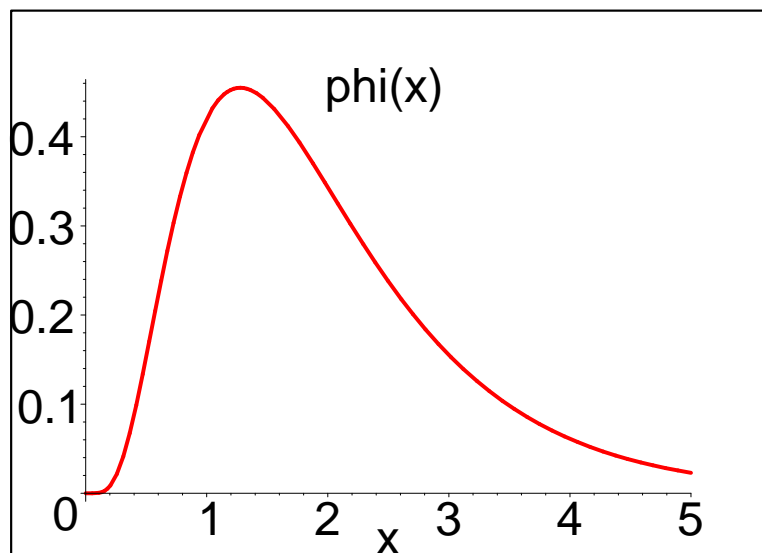
$$H_3(z) = \frac{q^{1+2} z^2}{(1 - (1 - q)z)(1 - (1 - q^2)z)(1 - (1 - q^3)z)}.$$

♡ (Prodinger'94) Euler transform $\xi := z/(1-z)$: $zH_k(z) = \frac{q^{\binom{k}{2}} \xi^{k-1}}{(1 - \xi q) \cdots (1 - \xi q^k)}.$

Exact moments of X and estimate q^X via Heine's transformation of q -calculus:
 mean is unbiased, variance $\rightsquigarrow 0.59$.

♥ Partial fraction expansions \leadsto asymptotic distribution
 = quantify typical behaviour + risk! (Exponential tails \gg Chebyshev ineq.)

We have $\mathbb{P}_n(X = \ell) \sim \phi(n/2^\ell)$, where $((q)_j := (1 - q) \cdots (1 - q_j).$)



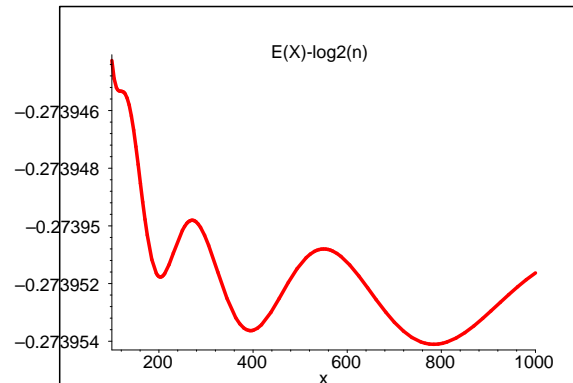
$$\phi(x) := \sum_{j \geq 0} \frac{(-1)^j q \binom{j}{2} e^{-xq^{-j}}}{(q)_\infty (q)_j}$$

♣ Fluctuations: $\dots, \frac{n}{2^L}, \dots, \frac{n}{4}, \dots$ depend on $L = \lfloor \log_2 n \rfloor$.
 cf. Szpankowski, Mahmoud, Fill, Prodinger, ...

Analyse storage utilization via Mellin transform

Approximate Counting

Mean(X) $- \log_2 n$:



The Mellin transform (F. Régnier Sedgewick 1985); (FIGoDu 1995)

$$f^*(s) := \int_0^{\infty} f(x)x^{s-1} dx.$$

(P_1) Mapping properties (complex analysis):

Asympt(f) \longleftrightarrow Singularities(f^*). Pole at $\sigma \longrightarrow \boxed{\approx x^\sigma}$

(P_2) Harmonic sums (superposition of models)

$$\left[\sum_k \lambda_k f(\mu_k x) \right] \mapsto \left[\sum_k \lambda_k (\mu_k)^{-s} \right] f^*(s).$$

♣ EXAMPLE: dyadic sum, $F(x) = \sum \phi\left(\frac{x}{2^\ell}\right) \rightsquigarrow F^*(s) = \frac{f^*(s)}{1 - 2^s}$.

Standard asymptotic terms + $x^{i\chi} \equiv \exp(i\chi \log x)$.

Cultural flashes

- Morris (1977): Counting a large number of events in small memory.
- The power of probabilistic machines & approximation (Freivalds IFIP 1977)
- The FTP protocol: Additive Increase Multiplicative Decrease (AIMD) leads to similar functions (Robert et al, 2001)
- Probability theory: Exponentials of Poisson processes (Yor et al, 2001)

4 **CARDINALITY ESTIMATORS**

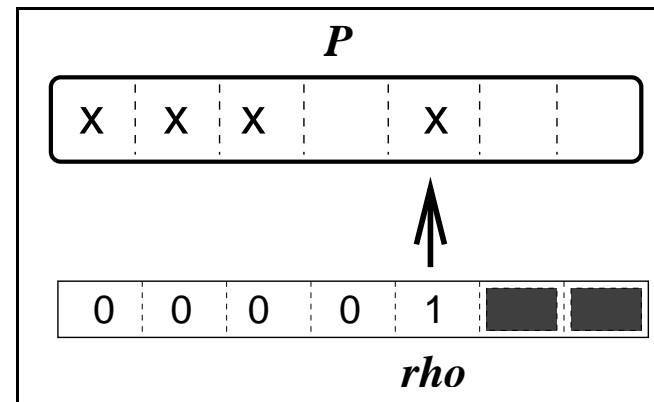
F_0 = Number of different values

- 1983–1985: (F-Martin, FOCS+JCSS) Probabilistic Count.
- 1987–1990: (Whang et al) Coupon Coll. Counting
- 1984–1990: (Wegner) (F90 COMP) Adaptive Sampling
- 1996: (Alon et al, STOC) F_p statistics \rightsquigarrow later
- 2000: (Indyk FOCS) Stable Law Counting \rightsquigarrow later
- 2001: (Estan-Varghese SIGCOMM) Multiresolution Bitmap
- 2003: (Durand-F ESA) Loglog Counting

4.1 Probabilistic Counting

Third Counting Algorithm
for cardinalities

:



ALG: Probabilistic Counting

Input: a stream S ; Output: cardinality $|S|$

For each $x \in S$ do /* $\rho \equiv$ position of leftmost 1-bit */

Set $\text{BITMAP}[\rho(\text{hash}(w))] := 1$; od;

Return P where P is position of first 0.

- P estimates $\log_2(\varphi n)$ for $\varphi \doteq 0.77351$
- Average over m trials $A = \frac{1}{m}[A_1 + \cdots + A_m]$; return $\frac{1}{\varphi} 2^A$.
- In fact, use stochastic averaging, which needs only *one* hash function: $S \mapsto (S_{000}, \dots, S_{111})$.
- Analysis provides

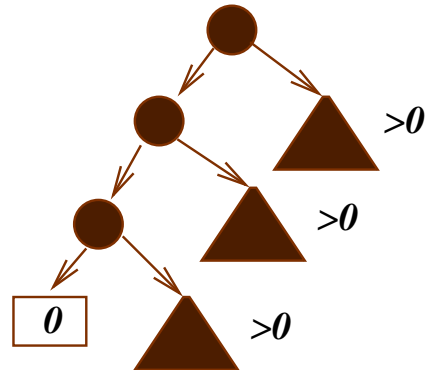
$$\varphi = \frac{e^\gamma}{\sqrt{2}} \prod_{m \geq 2}^* m^{\epsilon(m)}, \quad \epsilon(m) := (-1)^{\sum \text{bits}(m)}.$$

$\epsilon(19) = \epsilon(\langle 10011 \rangle_2) = (-1)^3 = -1$. Standard error is $0.78/\sqrt{m}$ for m Words of $\log_2 N$ bits. + Exponential Tails \gg Chebyshev.

(AMS96) and subsequent literature claim wrongly that *several* hash functions are needed!

Theorem [FM85]: Prob. Count. is asymptotically unbiased. Accuracy is $\frac{0.78}{\sqrt{m}}$ for m Words of size $\log_2 N$. E.g. 1,000W = 4kbytes \sim 2.5% accuracy.

Proof:
trie analysis



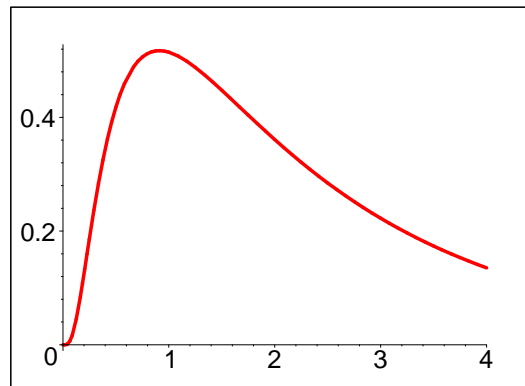
$$1 \cdot (e^{x/8} - 1)(e^{x/4} - 1)(e^{x/2} - 1)$$

$$(1 - q)(1 - q^2)(1 - q^4) \cdots = \sum_n \overbrace{(-1)^{\sum \text{bits}(n)}}^{\epsilon(n)} q^n.$$

Distribution:

$$Q(x) := e^{-x/2} \prod_{j=0}^{\infty} (1 - e^{-x2^j})$$

$$\mathbb{P}_n(X = \ell) \sim Q\left(\frac{n}{2^\ell}\right)$$



+ Mellin requires $N(s) := \sum_{n \geq 1} \frac{\epsilon(n)}{n^s}$. One finds $\log_2 \varphi \equiv -\Gamma'(1) - N'(0) + \frac{1}{2}$, &c.

Data mining of the Internet graph

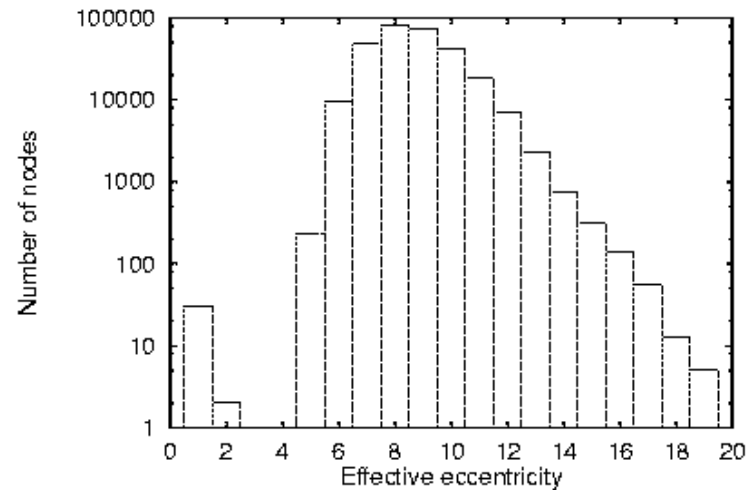
(Palmer, Gibbons, Faloutsos², Siganos 2001)

Internet graph: 285k nodes, 430k edges.

For each vertex v , define ball $B(v; R)$ of radius R .

Want: histograms of $|B(v, R)|$ $R = 1 \dots 20$

Get it in minutes of CPU rather than a day (400× speedup)



b) Histogram of diameters

Update procedure: $(h - 1) \mapsto h$ is

for each edge (u, v) do $B(v, h) := B(v, h) \cup B(u, h - 1)$

Use: *Probabilistic Counting*. Operate *in core*.

$$\text{Card}_{PC}(\mathcal{S} \cup \mathcal{T}) = \text{Card}_{PC}(\mathcal{S}) \vee \text{Card}_{PC}(\mathcal{T}).$$

where Card_{PC} is BITMAP evaluator of cardinalities.

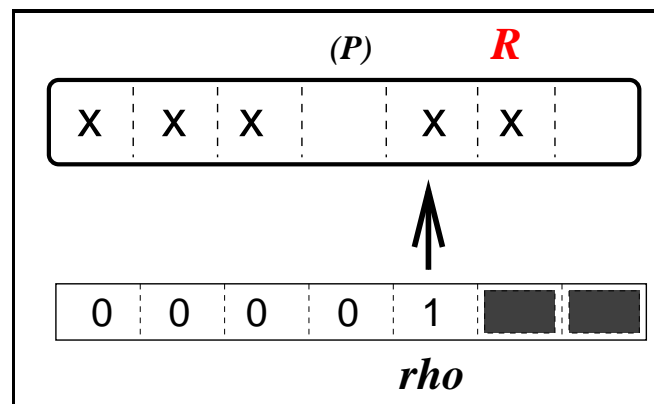
Allows for *fully distributed implementation*.

4.2 LogLog Counting

Fourth Counting Algorithm for cardinalities: (Durand-F, ESA 2003)

Claim: the best algorithm on the market!

- Hash values and get $\rho(h(x)) =$ position of leftmost 1-bit = a geometric RV $G(x)$.
- To set S associate $R(S) := \max_{v \in S} G(v)$.



- Max of geometric RVs are well-known (Prodinger*).
 $R(s)$ estimates $\sim \log(\hat{\varphi} \text{card}(S))$, with $\hat{\varphi} := e^{-\gamma} \sqrt{2}$.

- Do **stochastic averaging** with $m = 2^\ell$:

E.g., $S \cong \langle S_{00}, S_{01}, S_{10}, S_{11} \rangle$: count separately.

Return $\frac{m}{\hat{\varphi}} 2^{\text{Average}}$.

- ++ Switch to **Coupon Collector Counting** for small cardinalities.
- ++ Optimize by pruning discrepant values \rightsquigarrow **superLogLog**.

Theorem. LogLog needs m “bytes”, each of length $\log_2 \log N$.
 Accuracy is: $\frac{1.30}{\sqrt{m}}$ (BASIC) or $\frac{0.95}{\sqrt{m}}$ (SUPER)

PROOF: Generating Functions + Saddle-point depoissonization
 (Jacquet-Szpankowski) + Mellin. $1.30 \doteq \sqrt{\frac{1}{12} \log^2 2 + \frac{1}{6} \pi^2}$.



Whole of Shakespeare:

$m = 256$ small “bytes” of 4 bits each = 128bytes

```
ghfffghfghgghggggghghheehfhhfhggghghghhfgffffhhhiigfhhffgfiihfhhh
igigighfgihffffghigihghigfhhgeegeghgghhhgghhfhidiigihighihehhffgg
hfgighigffghdieghhhggghhfgghfiiheffghghihifgggffihgihfggighgiif
fjgfgjhhjiihfjhgehgghfhhfhjhiggghghihigghhihihgiihgfhlgjfgjjmfl
```

Estimate $n^\circ \approx 30,897$ against $n = 28,239$ distinct words

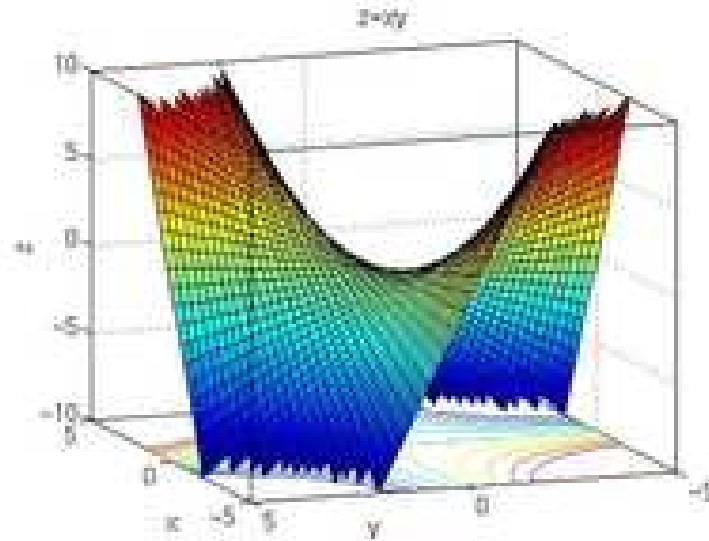
Error is +9.4% for 128 bytes(!!)

An aside: **Analytic depoissonization** (JaSz95⁺)

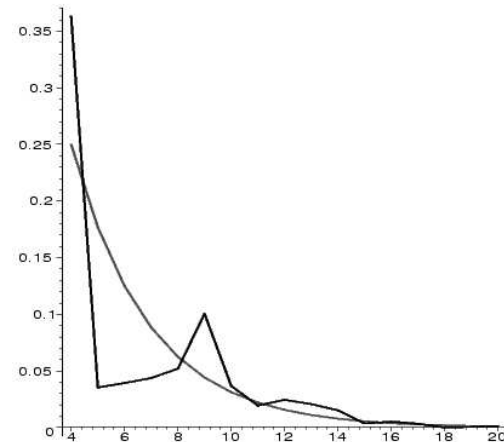
• PROBLEM: Recover asympt. f_n from $f(z) = \sum_n f_n \frac{z^n}{n!}$?

• Intuition: “with luck” $f_n \sim \phi(n)$ where $\phi(z) := e^{-z} f(z)$ is Poisson g.f..
(Here: “Luck” means good lifting of $\phi(z)$ to $\mathbb{C} \equiv$ Poisson flow of complex rate!)

$$f_n = \frac{n!}{2i\pi} \oint f(z) \frac{dz}{z^{n+1}} \approx \phi(n)$$



Features: Errors \approx *Gaussian*, seldom more than $2\times$ standard error. Algorithm *scales down* (for small cardinalities) and Algorithm *scales up* (large memory size): **HYBRIDIZE with Collision Counting**.



Mahābhārata: 8MB, 1M words, 177601 diff.

HTTP server: 400Mb log pages 1.8 M distinct req.

m	2^6 (50by)	2^{10} (0.8kby)	2^{14} (12kb)	2^{18} (200kb)
Obs:	8.9%	2.6%	1.2%	0.32%
σ :	11%	2.8%	0.7%	0.36%

Summary

Analytic results ($\lg \equiv \log_2$): Alg/Mem/Accuracy

CouponCC	AdSamp	ProbC	LogLog
$\approx \frac{N}{10}$ bits	$m \cdot \lg N$ Words	$m \cdot \lg \frac{N}{m}$ Words	$m \cdot \lg \lg \frac{N}{m}$ Bytes
$\approx 2\%$	$\frac{1.20}{\sqrt{m}}$ W	$\frac{0.78}{\sqrt{m}}$ W	$\approx \frac{1.30-0.95}{\sqrt{m}}$ By

F_0 statistics, $N = 10^8$ & 2% error

- Coupon Collector Counting = 1 Mbyte
- Adaptive Sampling = 16 kbytes
- Probabilistic Counting: = 8 kbytes
- Multiresolution bitmap (analysis?) = 5 kbytes?
- Loglog Counting = 2 kbytes

(NB: LogLog counting + compression $\rightsquigarrow \lg \lg N + O(m)$ bits !?)

5 FREQUENCY MOMENTS

5.1 AMS's F_2 algorithm

Recall: Alon, Matias, Szegedy (STOC 1996)^{***}

$$F_2 := \sum_v (f_v)^2,$$

where f_v is frequency of value v .

A beautifully simple idea: $\text{flip}(x) \equiv \epsilon(x) = \pm 1$ based on $\text{hash}(x)$.

ALG: F2;

Initialize $Z := 0$;

For each x in S do $Z := Z + \text{flip}(x)$.

Return Z^2 .

Collect m Z -values and average, with T -transform.

$$\mathbb{E}(Z^2) = \mathbb{E} \left(\sum_{x \in S} \epsilon(v) \right)^2 = \mathbb{E} \left(\sum_j f_j \cdot \epsilon(j) \right)^2 = \sum_j (f_j)^2.$$

(Actually, they prove stronger complexity result by complicated (impractical?) algorithm.) (What about stochastic averaging?)

5.2 Indyk's F_p algorithm

A beautiful idea of Piotr Indyk (FOCS 2000)^{***} for $F_p, p \in (0, 2)$.

- Stable law of parameter $p \in (0, 2)$: $\mathbb{E}(e^{itX}) = e^{-|t|^p}$.

No second moment; no 1st moment if $p \in (0, 1)$.

$$c_1 X_1 + c_2 X_2 \stackrel{\mathcal{L}}{\cong} \mu X, \text{ with } \mu := (c_1^p + c_2^p)^{1/p}.$$

ALG: F_p ;

Initialize $Z:=0$;

For each x in S do $Z := Z + \text{Stable}_\alpha(x)$.

Return Z .

Estimate F_p parameter from m copies of Z -values.

Remark: Use of $\log(|Z|)$ to estimate seems better than median(?)

6 CONCLUSIONS

For streams, using **practically $O(1)$ storage**, one can:

- **Sample** positions and even distinct values;
- **Estimate F_0, F_1, F_2, F_p** ($0 < p \leq 2$) even for *huge* data sets;
- Need **no assumption on nature of data**.



The algorithms are based on **randomization** \mapsto **Analysis fully applies**

- They **work exactly** as predicted on real-life data;
- They often have **a wonderfully elegant structure**;
- Their analysis involves **beautiful methods** for AofA: “Symbolic modelling by generating functions, Singularity analysis, Saddle Point and analytic depositions, Mellin transforms, stable laws and Mittag-Leffler functions, etc.”

That's All, Folks!

